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SERIES Q: SWITCHING AND SIGNALLING

Digital subscriber Signalling System No. 1 – Data link
layer

**ISDN user-network interface – Data link layer
specification**

ITU-T Recommendation Q.921

(Previously CCITT Recommendation)

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ITU-T RECOMMENDATION Q.921

ISDN USER-NETWORK INTERFACE – DATA LINK LAYER SPECIFICATION

Summary

This Recommendation specifies the Link Access Procedures on the D-channel (LAPD) of an ISDN customer access. Implementations of this Recommendation are in use in existing networks.

This Recommendation has been revised for the sake of clarity, to cover new functionality, and to reflect the requirements of Q.921 in Protocol Implementation Conformance Statement (PICS) proformas.

Clarifications are:

- Definition of terms used in Q.921 (new subclause 1.3);
- Amendment to 5.2.1 for clarification on signalling arrangement and circumstances under which TEI management and UI frame need not be supported;
- Amendment to 5.3.4 for clarification on TEI procedures;
- Amendment to Annex A of Q.921 for clarification on signalling arrangement and circumstances under which TEI management need not be supported.

New functions were specified as follows:

- New SAPI value 12 for teleaction communication (revised Table 2/Q.921);
- New Annex E "Provision of Multi-Selective Reject option".

Protocol Implementation Conformance Statement (PICS) proformas were amended, or generated, as appropriate, as follows:

- Amendment to Annex F (**former Annex E**) of Q.921 "PICS Basic Rate User-side" in order to reflect modifications to Q.921, as indicated above;
- New Annex G to Q.921 "PICS Basic Rate Network-side";
- New Annex H to Q.921 "PICS Rate User-side";
- New Annex I to Q.921 "PICS Primary Rate Network-side".

Source

ITU-T Recommendation Q.921 was revised by ITU-T Study Group 11 (1997-2000) and was approved under the WTSC Resolution No. 1 procedure on the 12th of September 1997.

FOREWORD

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NOTE

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Recommendation Q.921¹

ISDN USER-NETWORK INTERFACE – DATA LINK LAYER SPECIFICATION

(revised in 1997)

1 General

1.1 Introduction

This Recommendation specifies the frame structure, elements of procedure, format of fields and procedures for the proper operation of the Link Access Procedure on the D-channel, LAPD.

The concepts, terminology, overview description of LAPD functions and procedures, and the relationship with other Recommendations are described in general terms in Recommendation Q.920 [1].

NOTE 1 – As stated in Recommendation Q.920 [1], the term "data link layer" is used in the main text of this Recommendation. However, mainly in figures and tables, the terms "layer 2" and "L2" are used as abbreviations. Furthermore, in accordance with Recommendations Q.930 [2] and Q.931 [3], the term "layer 3" is used to indicate the layer above the data link layer.

NOTE 2 – All references within this Recommendation to "layer management entity" and/or "connection management entity" refer to those entities at the data link layer.

The abstract test suites for testing conformance to this Recommendation are contained in Recommendation Q.921 *bis* [4].

1.2 References

The following ITU-T Recommendations and other references contain provisions which, through reference in this text constitute provisions of this Recommendation. At the time of publication, the editions indicated were valid. All Recommendations and other references are subject to revision; all users of this Recommendation are therefore encouraged to investigate the possibility of applying the most recent edition of the Recommendations and other references listed below. A list of the currently valid ITU-T Recommendations is regularly published.

- [1] ITU-T Recommendation Q.920 (1993), *ISDN user-network interface data link layer – General aspects*.
- [2] ITU-T Recommendation Q.930 (1993), *ISDN user-network interface layer 3 – General aspects for basic call control*.
- [3] ITU-T Recommendation Q.931 (1993), *ISDN user-network interface layer 3 specification*.
- [4] ITU-T Recommendation Q.921 *bis* (1993), *Abstract test suite for LAPD conformance testing*.
- [5] ITU-T Recommendation I.430 (1995), *Basic user-network interface – layer 1 specification*.

¹ This Recommendation is published in the Series-I Recommendations under the number Recommendation I.441.

- [6] ITU-T Recommendation I.431 (1993), *Primary rate user-network interface – layer 1 specification*.
- [7] ITU-T Recommendation X.25 (1996), *Interface between Data Terminal Equipment (DTE) and Data Circuit-Terminating Equipment for terminals operating in the packet mode and connected to public data networks by dedicated circuit*.

1.3 Definitions

For the purposes of this Recommendation the following definitions apply, together with those given in Recommendations referenced.

1.3.1 Assignment Source Point (ASP): Layer management entity at the network side performing TEI management.

1.3.2 automatic TEI assignment: Layer management procedure between user side and network side (ASP) which associates within one interface a unique numeric value for a layer 2 terminal identity (TEI value) to a variable called TEI of a specific terminal equipment. The TEI, which is part of the DLCI, is selected by the ASP.

1.3.3 broadcast data link connection; broadcast connection: A connection with the capability to support more than two connection-endpoints [see (5.3.1.4/X.200) multi-endpoint-connection].

1.3.4 confirm (primitive): (see 3.2.7/X.210) A primitive issued by a service-provider to complete, at a particular service-access-point, some procedure previously invoked by a request at that service-access-point.

1.3.5 connection: (see 5.3.1.2/X.200) An association established by the "service provider" layer between two or more "service user" entities for the transfer of data.

1.3.6 connection-endpoint: (see 5.3.1.3/X.200) A terminator at one end of a connection within a service-access-point.

1.3.7 Connection Endpoint Identifier (CEI): (see 5.4.1.5/X.200) An identifier of a connection-endpoint which can be used to identify the corresponding connection at a service-access-point.

1.3.8 Connection Endpoint Suffix (CES): (see 5.4.1.6/X.200) A part of a connection-endpoint-identifier which is unique within the scope of a service-access-point.

1.3.9 Connection Management Entity (CME): An entity for the purpose of management of resources that have impact on an individual data link connection.

1.3.10 D-channel: (see Recommendation I.412) The D-channel represents the portion of the information-carrying capacity of the ISDN user-network interface primarily intended to carry access signalling information. In addition, a D-channel may also be used to carry other information such as packet-switched data, teleaction information, etc.

1.3.11 data link connection: (see Recommendation X.212) An association established by a Data Link Layer between two or more Data Link Service users for the transfer of data, which provides explicit identification of a set of Data Link data transmissions and agreement concerning the Data Link transmission services to be provided for the set.

NOTE – This definition clarifies the definition given in Recommendation X.200.

1.3.12 Data Link Connection Identifier (DLCI): An address conveyed in a PDU which indicates the source and destination of an intended instance of communication at the data link layer.

1.3.13 function: (see 5.2.1.7/X.200) A part of the activity of entities.

1.3.14 indication (primitive): (see 3.2.5/X.210) A primitive issued by a service-provider is either:

- i) to invoke some procedure; or
- ii) to indicate that a procedure has been invoked by the service-user at the peer service-access-point.

1.3.15 Integrated Services Digital Network (ISDN): (see Recommendation I.112, Nos. 307, 308) A network that provides or supports a range of different telecommunication services and provides digital connections between user-network interfaces.

1.3.16 layer: (see 5.2.1.2/X.200) A subdivision of the system architecture, constituted by subsystems of the same rank.

1.3.17 layer management: (see 8.1.6/X.200) Functions related to the management of the layer partly performed in the layer itself according to the protocol of the layer (activities such as activation and error control) and partly performed as a subset of systems management.

1.3.18 Layer Management Entity (LME): An entity for the purpose of management of resources that have layer-wide impact.

1.3.19 Link Access Procedure (LAP): Class of a procedure based on HDLC elements of procedures for use on the link layer.

1.3.20 network side: Location in relation to the user-network interface indicating that the context to which this term refers is at the network side of the user-network interface.

1.3.21 Network side system management entity: An entity for the purpose of management communications at the network side of the user-network interface.

1.3.22 non-automatic TEI assignment: Layer management local interaction between layer management entity and data link layer entity at the user side which associates within one interface a numeric value for a layer 2 terminal identity (TEI value) to a variable called TEI of a specific terminal equipment. The TEI, which is part of the DLCI, is selected by the user.

1.3.23 persistent deactivation: The term "persistent layer 1 deactivation" defines condition which shall be satisfied before the data link layer assumes layer 1 deactivation and takes the actions according to the protocol specification. Persistency is achieved if:

- i) the deactivation is an intended action within layer 1 caused by the functional block responsible for deactivation of the layer 1; or
- ii) layer 1 lost connectivity during a time interval, the value of which is outside the scope of this Recommendation, but which should be defined for each specific transmission facility.

1.3.24 point-to-point data link connection; point-to-point connection: A connection with two connection-endpoints.

1.3.25 protocol: (see 5.2.1.9/X.200) A set of rules and formats (semantic and syntactic) which determines the communication behaviour of entities in the performance of functions.

1.3.26 Protocol Data Unit (PDU): (see 5.6.1.3/X.200) A unit of data specified in a protocol and consisting of protocol-control-information and possibly user-data.

1.3.27 reference point: (see 2.3/I.411) Conceptual point dividing set of functions which form functional groups. In a specific access arrangement, a reference point may correspond to a physical interface between pieces of equipment, or there may not be any physical interface corresponding to the reference point. Physical interfaces that do not correspond to a reference point (e.g. transmission link interfaces) will not be the subject of ISDN user-network interface Recommendations.

1.3.28 request (primitive): (see 3.2.4/X.270) A primitive issued by a service-user to invoke some procedure.

1.3.29 response (primitive): (see 3.2.6/X.210) A primitive issued by a service-user to complete, at a particular service-access-point, some procedure previously invoked by an indication at that service-access-point.

1.3.30 service ("layer" service): (see 5.2.1.5/X.200) A capability of the providing layer and the layers beneath it, which is provided to "service user" entities at the boundary between the "service provider" layer and the "service user" layer.

1.3.31 Service Access Point (SAP): (see 5.2.1.8/X.200) The point at which services are provided by a "service provider" entity to a "service user" entity.

1.3.32 Service Data Unit (SDU): (see 5.6.1.4/X.200) An amount of interface-data whose identity is preserved from one end of a connection to the other.

1.3.33 service-primitive; primitive: (see 3.2.3/X.210) An abstract, implementation independent interaction between a service-user and the service-provider.

1.3.34 service-provider: (see 3.2.2/X.210) An abstract machine which models the behaviour of the totality of the entities providing the service, as viewed by the user.

1.3.35 service-user: (see 3.2.1/X.210) An abstract representation of the totality of those entities in a single system that make use of a service through a single access point.

1.3.36 system management: (see 8.1.4/X.200) Function in the Application Layer related to the management of various system resources and their status across all layers of the system architecture.

1.3.37 system management entity: (see 8.1.5/X.200) an entity for the purpose of systems-management communications.

1.3.38 Terminal Endpoint Identifier (TEI): Portion of a DLCI associated with one (point-to-point data link) or more than one (broadcast data link) terminal equipment.

1.3.39 user-data: (see 5.6.1.2/X.200) The data transferred between "service provider" entities on behalf of the "service user" entities for whom "service provider" entities are providing services.

1.3.40 user side: Location in relation to the user-network interface indicating that the context to which this term refers is at the user side of the user-network interface.

1.3.41 user side system management entity: An entity for the purpose of management communications at the user side of the user-network interface.

1.4 Abbreviations and acronyms used in this Recommendation

This Recommendation uses the following abbreviations

ACK	Acknowledgment
Ai	Action indicator
ASP	Assignment Source Point
CEI	Connection Endpoint Identifier
CES	Connection Endpoint suffix
C/R	Command/response field bit
DISC	Disconnect
DL-	Communication between Layer 3 and data link layer
DLCI	Data Link Connection Identifier
DM	Disconnected mode

EA	Extended address field bit
ERR	Error
ET	Exchange Termination
FCS	Frame Check Sequence
FRMR	Frame Reject
HDLC	High-level Data Link Control procedures
I	Information
ID	Identity
IND	Indication
ISDN	Integrated Services Digital Network
k	Maximum number of outstanding frames (window size)
L1	Layer 1
L2	Layer 2
L3	Layer 3
LAP	Link Access Procedure
LAPB	Link Access Procedure - Balanced
LAPD	Link Access Procedure on the D-channel
M	Modifier function bit
MDL-	Communication between layer management entity and data link layer
MPH-	Communication between system management and physical layer
N(R)	Receive sequence number
N(S)	Send sequence number
PDU	Protocol Data Unit
P/F	Poll/Final bit
PH-	Communication between data link layer and physical layer
PI	Parameter Identifier
PL	Parameter Length
PV	Parameter Value
RC	Retransmission Counter
REC	Receiver
REJ	Reject
REQ	Request
Ri	Reference number
RNR	Receive Not Ready
RR	Receive Ready
S	Supervisory

S ²	Supervisory function bit
SABME	Set Asynchronous Balanced Mode Extended
SAP	Service Access Point
SAPI	Service Access Point Identifier
SDL	Specification Description Language
SDU	Service Data Unit
TE	Terminal Equipment
TEI	Terminal Endpoint Identifier
TX	Transmit
U	Unnumbered
UA	Unnumbered Acknowledgement
UI	Unnumbered Information
V(A)	Acknowledge state variable
V(M)	Recovery state variable
V(R)	Receive state variable
V(S)	Send state variable
XID	Exchange identification

2 Frame structure for peer-to-peer communication

2.1 General

All data link layer peer-to-peer exchanges are in frames conforming to one of the formats shown in Figure 1. Two format types are shown in the figure: format A for frames where there is no information field and format B for frames containing an information field.

2.2 Flag sequence

All frames shall start and end with the flag sequence consisting of one 0 bit followed by six contiguous 1 bits and one 0 bit. The flag preceding the address field is defined as the opening flag. The flag following the Frame Check Sequence (FCS) field is defined as the closing flag. The closing flag may also serve as the opening flag of the next frame, in some applications. However, all receivers must be able to accommodate receipt of one or more consecutive flags. See ISDN User-Network Interfaces: Layer 1 Recommendations I.430 [5] and I.431 [6] for applicability.

2.3 Address field

The address field shall consist of two octets as illustrated in Figure 1. The format of the address field is defined in 3.2.

² A different acronym has to be found for Supervisory function bit.

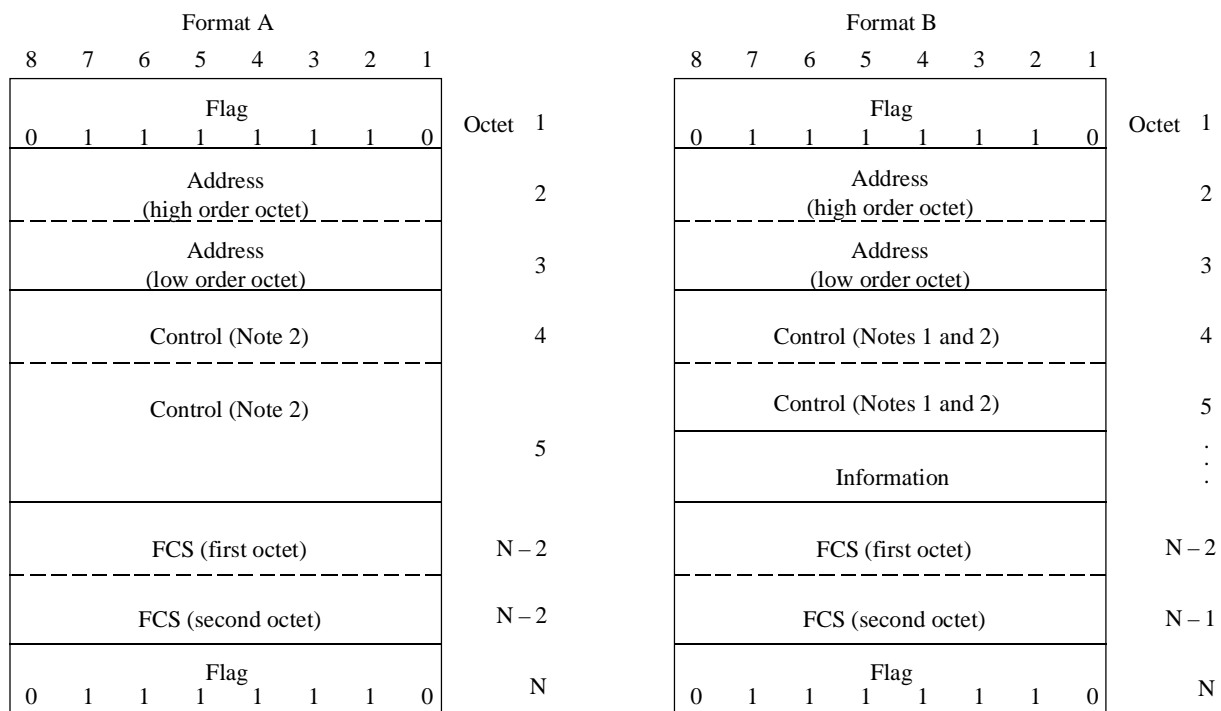
A single octet address field is reserved for LAPB (Link Access Procedure – Balanced) operation in order to allow a single LAPB [7] data link connection to be multiplexed along with LAPD data link connections.

NOTE – The support of a LAPB data link connection within the D-channel is optional at both the network and user side.

2.4 Control field

The control field shall consist of one or two octets. Figure 1 illustrates the two frame formats (A and B), each with a control field of one or two octets, depending upon the type of frame.

The format of the control field is defined in 3.4.



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NOTE 1 – For unacknowledged operation, format B applies and one octet control field is used.

NOTE 2 – For multiple frame operation, frames with sequence numbers contain a two-octet control field and frames without sequence numbers contain a one-octet control field. Connection management information transfer frames contain a one-octet control field.

Figure 1/Q.921 – Frame formats

2.5 Information field

The information field of a frame, when present, follows the control field (see 2.4 above) and precedes the frame check sequence (see 2.7 below). The contents of the information field shall consist of an integer number of octets.

The maximum number of octets in the information field is defined in 5.9.3.

2.6 Transparency

A transmitting data link layer entity shall examine the frame content between the opening and closing flag sequences, (address, control, information and FCS fields) and shall insert a 0 bit after all sequences of five contiguous 1 bits (including the last five bits of the FCS) to ensure that a flag or an

abort sequence is not simulated within the frame. A receiving data link layer entity shall examine the frame contents between the opening and closing flag sequences and shall discard any 0 bit which directly follows five contiguous 1 bits.

2.7 Frame Check Sequence (FCS) field

The FCS field shall be a 16-bit sequence. It shall be the ones complement of the sum (modulo 2) of:

- a) the remainder of $x^k (x^{15} + x^{14} + x^{13} + x^{12} + x^{11} + x^{10} + x^9 + x^8 + x^7 + x^6 + x^5 + x^4 + x^3 + x^2 + x + 1)$ divided (modulo 2) by the generator polynomial $x^{16} + x^{12} + x^5 + 1$, where k is the number of bits in the frame existing between, but not including, the final bit of the opening flag and the first bit of the FCS, excluding bits inserted for transparency; and
- b) the remainder of the division (modulo 2) by the generator polynomial $x^{16} + x^{12} + x^5 + 1$, of the product of x^{16} by the content of the frame existing between, but not including, the final bit of the opening flag and the first bit of the FCS, excluding bits inserted for transparency.

As a typical implementation at the transmitter, the initial content of the register of the device computing the remainder of the division is preset to all 1s and is then modified by division by the generator polynomial (as described above) on the address, control and information fields; the ones complement of the resulting remainder is transmitted as the 16-bit FCS.

As a typical implementation at the receiver, the initial content of the register of the device computing the remainder is preset to all 1s. The final remainder, after multiplication by x^{16} and then division (modulo 2) by the generator polynomial $x^{16} + x^{12} + x^5 + 1$ of the serial incoming protected bits and the FCS, will be 0001110100001111 (x^{15} through x^0 , respectively) in the absence of transmission errors.

2.8 Format convention

2.8.1 Numbering convention

The basic convention used in this Recommendation is illustrated in Figure 2. The bits are grouped into octets. The bits of an octet are shown horizontally and are numbered from 1 to 8. Multiple octets are shown vertically and are numbered from 1 to n .

2.8.2 Order of bit transmission

The octets are transmitted in ascending numerical order; inside an octet bit 1 is the first bit to be transmitted.

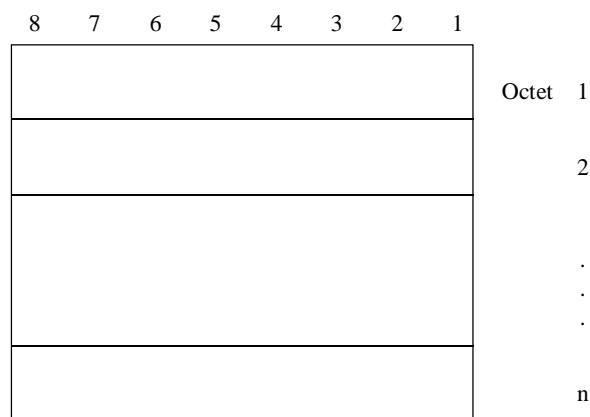
2.8.3 Field mapping convention

When a field is contained within a single octet, the lowest bit number of the field represents the lowest order value.

When a field spans more than one octet, the order of bit values within each octet progressively decreases as the octet number increases. The lowest bit number associated with the field represents the lowest order value.

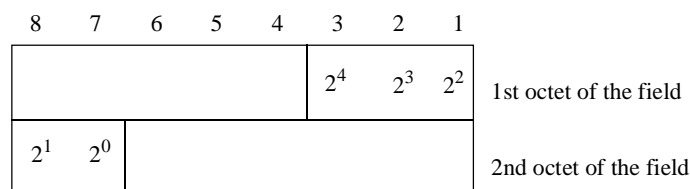
For example, a bit number can be identified as a couple (o, b) where o is the octet number and b is the relative bit number within the octet. Figure 3 illustrates a field that spans from bit $(1, 3)$ to bit $(2, 7)$. The high order bit of the field is mapped on bit $(1, 3)$ and the low order bit is mapped on bit $(2, 7)$.

An exception to the preceding field mapping convention is the data link layer FCS field, which spans two octets. In this case, bit 1 of the first octet is the high order bit and bit 8 of the second octet is the low order bit (see Figure 4).



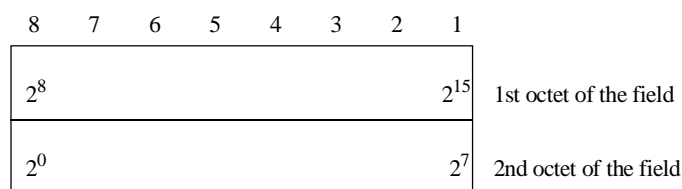
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Figure 2/Q.921 – Format convention



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Figure 3/Q.921 – Field mapping convention



T1161610-94

Figure 4/Q.921 – FCS mapping convention

2.9 Invalid frames

An invalid frame is a frame which:

- a) is not properly bounded by two flags; or
- b) has fewer than six octets between flags of frames that contain sequence numbers and fewer than five octets between flags of frames that do not contain sequence numbers; or
- c) does not consist of an integral number of octets prior to zero bit insertion or following zero bit extraction; or
- d) contains a frame check sequence error; or
- e) contains a single octet address field; or
- f) contains a service access point identifier (see 3.3.3) which is not supported by the receiver.

Invalid frames shall be discarded without notification to the sender. No action is taken as the result of that frame.

2.10 Frame abort

Receipt of seven or more contiguous 1 bits shall be interpreted as an abort and the data link layer shall ignore the frame currently being received.

3 Elements of procedures and formats of fields for data link layer peer-to-peer communication

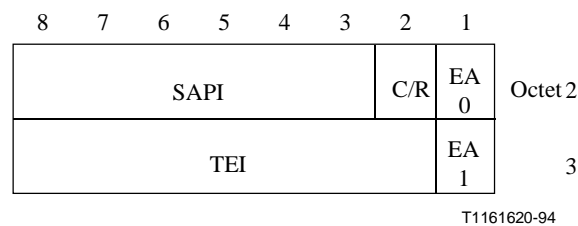
3.1 General

The elements of procedures define the commands and responses that are used on the data link connections carried on the D-channel.

Procedures are derived from these elements of procedures and are described in clause 5.

3.2 Address field format

The address field format shown in Figure 5 contains the address field extension bits, a command/response indication bit, a data link layer Service Access Point Identifier (SAPI) subfield, and a Terminal Endpoint Identifier (TEI) subfield.



- EA Address field extension bit
- C/R Command/response field bit
- SAPI Service Access Point Identifier
- TEI Terminal Endpoint Identifier

Figure 5/Q.921 – Address field format

3.3 Address field variables

3.3.1 Address field extension bit (EA)

The address field range is extended by reserving the first transmitted bit of the address field octets to indicate the final octet of the address field. The presence of a 1 in the first bit of an address field octet signals that it is the final octet of the address field. The double octet address field for LAPD operation shall have bit 1 of the first octet set to a 0 and bit 1 of the second octet set to 1, otherwise the frame shall be ignored.

3.3.2 Command/response field bit (C/R)

The C/R bit identifies a frame as either a command or a response. The user side shall send commands with the C/R bit set to 0, and responses with the C/R bit set to 1. The network side shall

do the opposite; that is, commands are sent with C/R set to 1, and responses are sent with C/R set to 0. The combinations for the network side and user side are shown in Table 1.

Table 1/Q.921 – C/R field bit usage

Command/Response	Direction	C/R value
Command	Network side → user side	1
	User side → network side	0
Response	Network side → user side	0
	User side → network side	1

In conformance with HDLC (High-level Data Link Control procedures) rules, both peer entities on a point-to-point data link connection use the same Data Link Connection Identifier (DLCI) composed of a SAPI-TEI where SAPI and TEI conform to the definitions contained in 3.3.3 and 3.3.4 and define the data link connection as described in 3.4.1/Q.920 [1].

3.3.3 Service Access Point Identifier (SAPI)

The SAPI identifies a point at which data link layer services are provided by a data link layer entity type to a layer 3 or management entity. Consequently, the SAPI specifies a data link layer entity type that should process a data link layer frame and also a layer 3 or management entity which is to receive information carried by the data link layer frame. The SAPI allows 64 service access points to be specified, where bit 3 of the address field octet containing the SAPI is the least significant binary digit and bit 8 is the most significant. The SAPI values are allocated as shown in Table 2.

Table 2/Q.921

SAPI value	Related layer 3 or management entity
0	Call control procedures
1 – 11	Reserved for future standardization
12	Teleaction communication
13 – 15	Reserved for future standardization
16	Packet communication conforming to X.25 level 3 procedures
17 – 31	Reserved for future standardization
63	Layer 2 management procedures
All others	Not available for Q.921 procedures
NOTE – The reservation of SAPI values for experimental purposes is for further study.	

3.3.4 Terminal Endpoint Identifier (TEI)

It is possible to associate a TEI with a single Terminal Equipment (TE) for a point-to-point data link connection. If a TEI is not the group TEI (see 3.3.4.1) and is not associated with any TE, that TEI is unassigned. A TE may contain one or more TEIs used for point-to-point data transfer. The TEI for a broadcast data link connection is associated with all user side data link layer entities containing the same SAPI. The TEI subfield allows 128 values where bit 2 of the address field octet containing the TEI is the least significant binary digit and bit 8 is the most significant binary digit. The following conventions shall apply in the assignment of these values.

3.3.4.1 TEI for broadcast data link connection

The TEI subfield bit pattern 111 1111 (= 127) is defined as the group TEI. The group TEI is assigned permanently to the broadcast data link connection associated with the addressed Service Access Point (SAP).

3.3.4.2 TEI for point-to-point data link connection

TEI, values other than 127 are used for the point-to-point data link connections associated with the addressed SAP. The range of TEI values shall be allocated as shown in Table 3.

Table 3/Q.921

TEI Value	User Type
0-63	Non-automatic TEI assignment user equipment
64-126	Automatic TEI assignment user equipment

Non-automatic TEI values are selected by the user, and their allocation is the responsibility of the user.

Automatic TEI values are selected by the network, and their allocation is the responsibility of the network.

For further information regarding point-to-point situations, see Annex A.

3.4 Control field formats

The control field identifies the type of frame which will be either a command or response. The control field will contain sequence numbers, where applicable.

Three types of control field formats are specified: numbered information transfer (I format), supervisory functions (S format), and unnumbered information transfers and control functions (U format). The control field formats are shown in Table 4.

Table 4/Q.921 – Control field formats

Control field bits (modulo 128)	8	7	6	5	4	3	2	1	
I format	N(S)							0	Octet 4
	N(R)							P	5
S format	X	X	X	X	S	S	0	1	Octet 4
	N(R)							P/F	5
U format	M	M	M	P/F	M	M	1	1	Octet 4
N(S)	Transmitter send sequence number		M	Modifier function bit					
N(R)	Transmitter receive sequence number		P/F	Poll bit when issued as a command, final bit when issued as a response					
S	Supervisory function bit		X	Reserved and set to 0					

3.4.1 Information transfer (I) format

The I format shall be used to perform an information transfer between layer 3 entities. The functions of N(S), N(R) and P (defined in 3.5) are independent; that is, each I frame has an N(S) sequence

number, an N(R) sequence number which may or may not acknowledge additional I frames received by the data link layer entity, and a P bit that may be set to 0 or 1.

The use of N(S), N(R) and P is defined in clause 5.

3.4.2 Supervisory (S) format

The S format shall be used to perform data link supervisory control functions such as: acknowledge I frames, request retransmission of I frames, and request a temporary suspension of transmission of I frames. The functions of N(R) and P/F are independent, that is, each supervisory frame has an N(R) sequence number which may or may not acknowledge additional I frames received by the data link layer entity, and a P/F bit that may be set to 0 or 1.

3.4.3 Unnumbered (U) format

The U format shall be used to provide additional data link control functions and unnumbered information transfers for unacknowledged information transfer. This format does not contain sequence numbers. It does include a P/F bit that may be set to 0 or 1.

3.5 Control field parameters and associated state variables

The various parameters associated with the control field formats are described in this subclause. The coding of the bits within these parameters is such that the lowest numbered bit within the parameter field is the least significant bit.

3.5.1 Poll/Final (P/F) bit

All frames contain the Poll/Final (P/F) bit. The P/F bit serves a function in both command frames and response frames. In command frames the P/F bit is referred to as the P bit. In response frames it is referred to as the F bit. The P bit set to 1 is used by a data link layer entity to solicit (poll) a response frame from the peer data link layer entity. The F bit set to 1 is used by a data link layer entity to indicate the response frame transmitted as a result of a soliciting (poll) command.

The use of the P/F bit is described in clause 5.

3.5.2 Multiple frame operation – variables and sequence numbers

3.5.2.1 Modulus

Each I frame is sequentially numbered and may have the value 0 through n minus 1 (where n is the modulus of the sequence numbers). The modulus equals 128 and the sequence numbers cycle through the entire range, 0 through 127.

NOTE – All arithmetic operations on state variables and sequence numbers contained in this Recommendation are affected by the modulus operation.

3.5.2.2 Send state variable V(S)

Each point-to-point data link connection endpoint shall have an associated V(S) when using I frame commands. V(S) denotes the sequence number of the next I frame to be transmitted. The V(S) can take on the value 0 through n minus 1. The value of V(S) shall be incremented by 1 with each successive I frame transmission, and shall not exceed V(A) by more than the maximum number of outstanding I frames k . The value of k may be in the range of $1 \leq k \leq 127$.

3.5.2.3 Acknowledge state variable V(A)

Each point-to-point data link connection endpoint shall have an associated V(A) when using I frame commands and supervisory frame commands/responses. V(A) identifies the last I frame that has been

acknowledged by its peer [$V(A) - 1$ equals the $N(S)$ of the last acknowledged I frame]. $V(A)$ can take on the value 0 through n minus 1. The value of $V(A)$ shall be updated by the valid $N(R)$ values received from its peer (see 3.5.2.6). A valid $N(R)$ value is one that is in the range $V(A) \leq N(R) \leq V(S)$.

3.5.2.4 Send sequence number $N(S)$

Only I frames contain $N(S)$, the send sequence number of transmitted I frames. At the time that an in-sequence I frame is designated for transmission, the value of $N(S)$ is set equal to $V(S)$.

3.5.2.5 Receive state variable $V(R)$

Each point-to-point data link connection endpoint shall have an associated $V(R)$ when using I frame commands and supervisory frame commands/responses. $V(R)$ denotes the sequence number of the next in-sequence I frame expected to be received. $V(R)$ can take on the value 0 through n minus 1. The value of $V(R)$ shall be incremented by one with the receipt of an error-free, in-sequence I frame whose $N(S)$ equals $V(R)$.

3.5.2.6 Receive sequence number $N(R)$

All I frames and supervisory frames contain $N(R)$, the expected send sequence number of the next received I frame. At the time that a frame of the above types is designated for transmission, the value of $N(R)$ is set equal to $V(R)$. $N(R)$ indicates that the data link layer entity transmitting the $N(R)$ has correctly received all I frames numbered up to and including $N(R) - 1$.

3.5.3 Unacknowledged operation – variables and parameters

No variables are defined. One parameter is defined, $N201$ (see 5.9.3).

3.6 Frame types

3.6.1 Commands and responses

The following commands and responses are used by either the user or the network data link layer entities and are represented in Table 5. Each data link connection shall support the full set of commands and responses for each application implemented. The frame types associated with each of the two applications are identified in Table 5.

Frame types associated with an application not implemented shall be discarded and no action shall be taken as a result of that frame.

For purposes of the LAPD procedures in each application, those encodings not identified in Table 5 are identified as undefined command and response control fields. The actions to be taken are specified in 5.8.5.

The commands and responses in Table 5 are defined in 3.6.2 to 3.6.12.

Table 5/Q.921 – Commands and responses (modulo 128)

Application	Format	Commands	Responses	Encoding								Octet	
				8	7	6	5	4	3	2	1		
Unacknowledged and Multiple Frame Acknowledged Information Transfer	Information transfer	I (information)		N(S)							0	4	
				N(R)							P	5	
	Supervisory	RR (receive ready)	RR (receive ready)	0	0	0	0	0	0	0	1	4	
				N(R)							P/F	5	
		RNR (receive not ready)	RNR (receive not ready)	0	0	0	0	0	1	0	1	4	
				N(R)							P/F	5	
		REJ (reject)	REJ (reject)	0	0	0	0	1	0	0	1	4	
				N(R)							P/F	5	
	Unnumbered	SABME (set asynchronous balanced mode extended)		0	1	1	P	1	1	1	1	4	
				0	0	0	F	1	1	1	1	4	
		DISC (disconnect)	UI (unnumbered information)		0	0	0	P	0	0	1	1	4
					0	1	0	P	0	0	1	1	4
					0	1	1	F	0	0	1	1	4
					1	0	0	F	0	1	1	1	4
XID (Exchange Identification) (Note)	XID (Exchange Identification) (Note)		1	0	1	P/F	1	1	1	1	4		

NOTE – Use of the XID frame other than for parameter negotiation procedures (see 5.4) is for further study.

3.6.2 Information (I) command

The function of the information (I) command is to transfer, across a data link connection, sequentially numbered frames containing information fields provided by layer 3. This command is used in the multiple frame operation on point-to-point data link connections.

3.6.3 Set Asynchronous Balanced Mode Extended (SABME) command

The SABME unnumbered command is used to place the addressed user side or network side into modulo 128 multiple frame acknowledged operation.

No information field is permitted with the SABME command. A data link layer entity confirms acceptance of an SABME command by the transmission at the first opportunity of a UA response. Upon acceptance of this command, the data link layer entity's V(S), V(A) and V(R) are set to 0. The transmission of an SABME command indicates the clearance of all exception conditions.

Previously transmitted I frames that are unacknowledged when this command is processed remain unacknowledged and are discarded. It is the responsibility of a higher level (for example, layer 3) or the management entity to recover from the possible loss of the contents of such I frames.

3.6.4 Disconnect (DISC) command

The DISC unnumbered command is used to terminate the multiple frame operation.

No information field is permitted with the DISC command. The data link layer entity receiving the DISC command confirms the acceptance of a DISC command by the transmission of a UA response.

The data link layer entity sending the DISC command terminates the multiple frame operation when it receives the acknowledging UA or DM response.

Previously transmitted I frames that are unacknowledged when this command is processed remain unacknowledged and are discarded. It is the responsibility of a higher level (for example, layer 3) or the management entity to recover from the possible loss of the contents of such I frames.

3.6.5 Unnumbered information (UI) command

When a layer 3 or management entity requests unacknowledged information transfer, the UI unnumbered command is used to send information to its peer without affecting data link layer variables. UI command frames do not carry a sequence number and therefore, the UI frame may be lost without notification.

3.6.6 Receive ready (RR) command/response

The RR supervisory frame is used by a data link layer entity to:

- a) indicate it is ready to receive an I frame;
- b) acknowledge previously received I frames numbered up to and including $N(R) - 1$ (as defined in clause 5); and
- c) clear a busy condition that was indicated by the earlier transmission of an RNR frame by that same data link layer entity.

In addition to indicating the status of a data link layer entity, the RR command with the P bit set to 1 may be used by the data link layer entity to ask for the status of its peer data link layer entity.

3.6.7 Reject (REJ) command/response

The REJ supervisory frame is used by a data link layer entity to request retransmission of I frames starting with the frame numbered $N(R)$. The value of $N(R)$ in the REJ frame acknowledges I frames numbered up to and including $N(R) - 1$. New I frames pending initial transmission shall be transmitted following the retransmitted I frame(s).

Only one REJ exception condition for a given direction of information transfer is established at a time. The REJ exception condition is cleared (reset) upon the receipt of an I frame with an $N(S)$ equal to the $N(R)$ of the REJ frame. An optional procedure for the retransmission of an REJ response frame is described in Appendix I.

The transmission of an REJ frame shall also indicate the clearance of any busy condition within the sending data link layer entity that was reported by the earlier transmission of an RNR frame by that same data link layer entity.

In addition to indicating the status of a data link layer entity, the REJ command with the P bit set to 1 may be used by the data link layer entity to ask for the status of its peer data link layer entity.

3.6.8 Receive not ready (RNR) command/response

The RNR supervisory frame is used by a data link layer entity to indicate a busy condition; that is, a temporary inability to accept additional incoming I frames. The value of $N(R)$ in the RNR frame acknowledges I frames numbered up to and including $N(R) - 1$.

In addition to indicating the status of a data link layer entity, the RNR command with the P bit set to 1 may be used by the data link layer entity to ask for the status of its peer data link layer entity.

3.6.9 Unnumbered acknowledgement (UA) response

The UA unnumbered response is used by a data link layer entity to acknowledge the receipt and acceptance of the mode-setting commands (SABME or DISC). Received mode-setting commands

are not processed until the UA response is transmitted. No information field is permitted with the UA response. The transmission of the UA response indicates the clearance of any busy condition that was reported by the earlier transmission of an RNR frame by that same data link layer entity.

3.6.10 Disconnected mode (DM) response

The DM unnumbered response is used by a data link layer entity to report to its peer that the data link layer is in a state such that multiple frame operation cannot be performed. No information field is permitted with the DM response.

3.6.11 Frame reject (FRMR) response

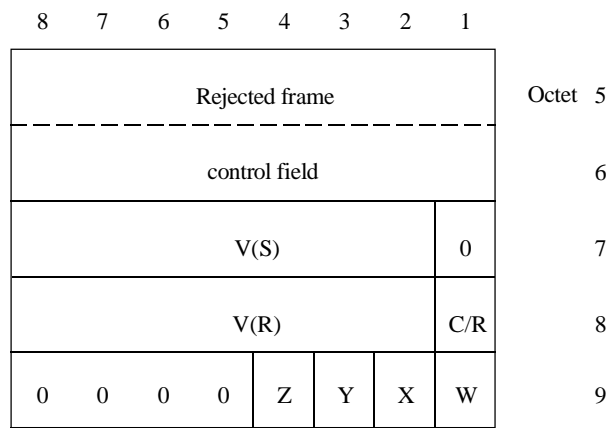
The FRMR unnumbered response may be received by a data link layer entity as a report of an error condition not recoverable by retransmission of the identical frame, i.e. at least one of the following error conditions resulting from the receipt of a valid frame:

- a) the receipt of a command or response control field that is undefined;
- b) the receipt of a supervisory or unnumbered frame with incorrect length;
- c) the receipt of an invalid N(R); or
- d) the receipt of a frame with an information field which exceeds the maximum established length.

An undefined control field is any of the control field encodings that are not identified in Table 5.

A valid N(R) value is one that is in the range $V(A) \leq N(R) \leq V(S)$.

An information field which immediately follows the control field and consists of five octets (modulo 128 operation) is returned with this response and provides the reason for the FRMR response. This information field format is given in Figure 6.



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NOTE 1 – Rejected frame control field is the control field of the received frame which caused the frame reject. When the rejected frame is an unnumbered frame, the control field of the rejected frame is positioned in octet 5, with octet 6 set to 0000 0000.

NOTE 2 – V(S) is the current send state variable value on the user side or network side reporting the rejection condition.

NOTE 3 – C/R is set to 1 if the frame rejected was a response and is set to 0 if the frame rejected was a command.

NOTE 4 – V(R) is the current receive state variable value on the user side or network side reporting the rejection condition.

NOTE 5 – W set to 1 indicates that the control field received and returned in octets 5 and 6 was undefined.

NOTE 6 – X set to 1 indicates that the control field received and returned in octets 5 and 6 was considered invalid because the frame contained an information field which is not permitted with this frame or is a supervisory or unnumbered frame with incorrect length. Bit W must be set to 1 in conjunction with this bit.

NOTE 7 – Y set to 1 indicates that the information field received exceeded the maximum established information field length (N201) of the user side or network side reporting the rejection condition.

NOTE 8 – Z set to 1 indicates that the control field received and returned in octets 5 and 6 contained an invalid N(R).

NOTE 9 – Octet 7 bit 1 and octet 9 bits 5 through 8 shall be set to 0.

Figure 6/Q.921 – FRMR information field format – extended (modulo 128) operation

3.6.12 Exchange identification (XID) command/response

The XID frame may contain an information field in which the identification information is conveyed. The exchange of XID frames is a compelled arrangement used in connection management (i.e. when a peer connection management entity receives an XID command, it shall respond with an XID response at the earliest time possible). No sequence numbers are contained within the control field.

The information field is not mandatory. However, if a valid XID command contains an information field and the receiver can interpret its contents, the receiver should then respond with an XID response also containing an information field. If the information field cannot be interpreted by the receiving entity, or a zero length information field has been received, an XID response frame shall be issued containing a zero length information field. The maximum length of the information field must conform to the value N201.

Sending or receiving an XID frame shall have no effect on the operational mode or state variables associated with the data link layer entities.

4 Elements for layer-to-layer communication

4.1 General

Communications between layers and, for this Recommendation, between the data link layer and the layer management are accomplished by means of primitives.

Primitives represent, in an abstract way, the logical exchange of information and control between the data link and adjacent layers. They do not specify or constrain implementations.

Primitives consist of commands and their respective responses associated with the services requested of a lower layer. The general syntax of a primitive is:

XX – Generic name – Type: Parameters

where XX designates the interface across which the primitive flows. For this Recommendation, XX is:

- DL for communication between layer 3 and the data link layer;
- PH for communication between the data link layer and the physical layer;
- MDL for communication between the layer management and the data link layer; or
- MPH for communication between the management entity and the physical layer.

4.1.1 Generic names

The generic name specifies the activity that should be performed. Table 6 illustrates the primitives defined in this Recommendation. Note that not all primitives have associated parameters.

The primitive generic names that are defined in this Recommendation are:

4.1.1.1 DL-ESTABLISH

The DL-ESTABLISH primitives are used to request, indicate and confirm the outcome of the procedures for establishing multiple frame operation.

4.1.1.2 DL-RELEASE

The DL-RELEASE primitives are used to request, indicate and confirm the outcome of the procedures for terminating a previously established multiple frame operation, or for reporting an unsuccessful establishment attempt.

4.1.1.3 DL-DATA

The DL-DATA primitives are used to request and indicate SDUs containing layer 3 PDUs which are to be transmitted, or have been received, by the data link layer using the acknowledged information transfer service.

4.1.1.4 DL-UNIT DATA

The DL-UNIT DATA primitives are used to request and indicate SDUs containing layer 3 PDUs which are to be transmitted, or have been received, by the data link layer using the unacknowledged information transfer service.

4.1.1.5 MDL-ASSIGN

The MDL-ASSIGN primitives are used by the layer management entity to request that the data link layer associate the TEI value contained within the parameter data of the primitive with the specified Connection Endpoint Suffix (CES), across all SAPIs which support point-to-point data links. The MDL-ASSIGN primitive is used by the data link layer to indicate to the layer management entity the need for a TEI value to be associated with the CES specified in the primitive parameter data.

4.1.1.6 MDL-REMOVE

The MDL-REMOVE primitives are used by the layer management entity to request that the data link layer remove the association of the specified TEI value with the specified CES, across all SAPIs which support point-to-point data links. The TEI and CES are specified by the MDL-REMOVE primitive parameter data.

4.1.1.7 MDL-ERROR

The MDL-ERROR primitives are used to indicate to the connection management entity that an error has occurred, associated with a previous management function request or detected as a result of communication with the data link layer peer entity. The layer management entity may respond with an MDL-ERROR primitive if the layer management entity cannot obtain a TEI value.

Table 6/Q.921 – Primitives associated with this Recommendation

Generic name	Type				Parameters		Parameter data contents
	Request	Indication	Response	Confirm	Priority indicator	Parameter data	
L3 ↔ L2							(Note 1)
DL-ESTABLISH	X	X	–	X	–	–	–
DL-RELEASE	X	X	–	X	–	–	–
DL-DATA	X	X	–	–	–	X	Layer 3 PDU (peer-to-peer message)
DL-UNIT DATA	X	X	–	–	–	X	Layer 3 PDU (peer-to-peer message)
M ↔ L2							
MDL-ASSIGN	X	X	–	–	–	X	TEI value, CES (Note 2)
MDL-REMOVE	X	–	–	–	–	X	TEI value, CES
MDL-ERROR	–	X	X	–	–	X	Reason for error message
MDL-UNIT DATA	X	X	–	–	–	X	Layer management PDU (peer-to-peer message)
MDL-XID	X	X	X	X	–	X	Connection management PDU (peer-to-peer XID frame)
L2 ↔ L1							
PH-DATA	X	X	–	–	X	X	Data link layer PDU (peer-to-peer frame)
PH-ACTIVATE	X	X	–	–	–	–	–
PH-DEACTIVATE	–	X	–	–	–	–	–
M ↔ L1							
MPH-ACTIVATE	–	X	–	–	–	–	–

Table 6/Q.921 – Primitives associated with this Recommendation (concluded)

Generic name	Type				Parameters		Parameter data contents
	Request	Indication	Response	Confirm	Priority indicator	Parameter data	
MPH-DEACTIVATE	X	X	–	–	–	–	–
MPH-INFORMATION	–	X	–	–	–	X	Connected/disconnected
X Exists – Does not exist L3 ↔ L2 Layer 3/data link layer boundary L2 ↔ L1 Data link layer/physical layer boundary M ↔ L2 Management entity/data link layer boundary M ↔ L1 Management entity/physical layer boundary NOTE 1 – Although not shown below, the CES is implicitly associated with each L3-L2 primitive, indicating the applicable connection endpoint. NOTE 2 – TEI value is included only in the MDL-ASSIGN request.							

4.1.1.8 MDL-UNIT DATA

The MDL-UNIT DATA primitives are used to request and indicate SDUs containing layer management PDUs which are to be transmitted, or have been received, by the data link layer using the unacknowledged information transfer service.

4.1.1.9 MDL-XID

The MDL-XID primitives are used by the connection management entity to request and respond to the data link layer and by the data link layer to indicate and confirm to the connection management entity service data units which are to be transmitted, or have been received, by the data link layer using the XID procedures.

4.1.1.10 PH-DATA

The PH-DATA primitives are used to request and indicate SDUs containing frames used for data link layer peer-to-peer communications passed to and from the physical layer.

4.1.1.11 PH-ACTIVATE

The PH-ACTIVATE primitives are used to request activation of the physical layer connection or to indicate that the physical layer connection has been activated.

4.1.1.12 PH-DEACTIVATE

The PH-DEACTIVATE primitive is used to indicate that the physical layer connection has been deactivated.

4.1.1.13 MPH-ACTIVATE (see Appendix III)

The MPH-ACTIVATE primitive is used to indicate that the physical layer connection has been activated.

4.1.1.14 MPH-DEACTIVATE (see Appendix III)

The MPH-DEACTIVATE primitives are used to request deactivation of the physical layer connection or to indicate that the physical layer connection has been deactivated. The MPH-DEACTIVATE request primitive is for use by the network side system management entity.

4.1.1.15 MPH-INFORMATION

The MPH-INFORMATION primitive is for use by the user side layer management entity, and provides an indication as to whether the terminal is:

- connected; or
- disconnected or unable to provide sufficient power to support the TEI management procedures.

4.1.2 Primitive types

The primitive types defined in this Recommendation are:

4.1.2.1 request

The request primitive type is used when a higher layer or layer management is requesting a service from the lower layer.

4.1.2.2 indication

The indication primitive type is used by a layer providing a service to inform the higher layer or layer management.

4.1.2.3 response

The response primitive type is used by layer management as a consequence of the indication primitive type.

4.1.2.4 confirm

The confirm primitive type is used by the layer providing the requested service to confirm that the activity has been completed.

Figure 7 illustrates the relationship of the primitive types to layer 3 and the data link layer.

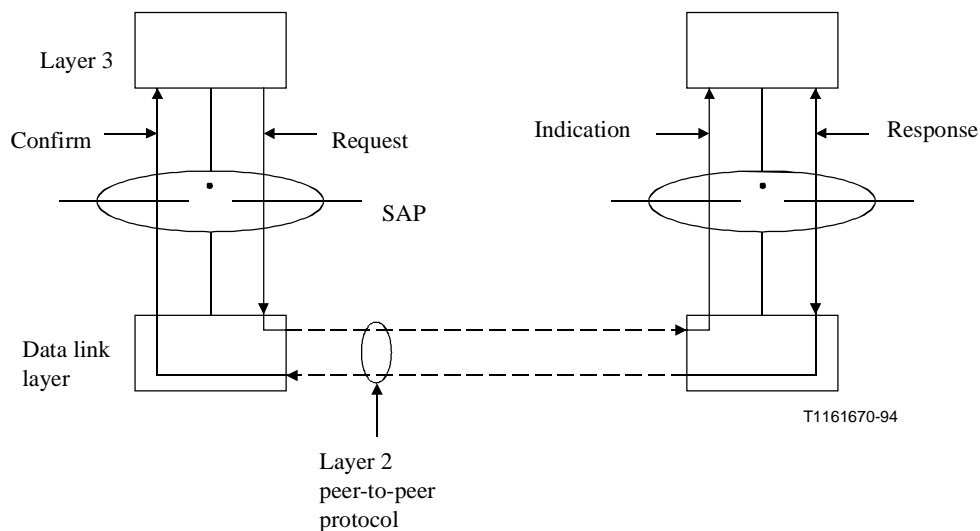


Figure 7/Q.921 – Relationship of the primitive types to layer 3 and the data link layer

4.1.3 Parameter definition

A parameter consists of two parts, the priority indicator and parameter data such as: service user data, reasons or TEI.

4.1.3.1 Priority indicator

Since several SAPs may exist on the network side or user side, SDUs sent across one SAP may contend with those sent across other SAPs for the physical resources available for information transfer. The priority indicator is used to determine which SDU will have greater priority when contention exists. The priority indicator is only needed at the user side for distinguishing SDUs sent across the SAP with a SAPI value of 0 from all other SDUs.

4.1.3.2 Parameter data

The parameter data is associated with a primitive and contains information related to the service. In the case of the DATA primitives, the parameter data contains the SDU which allows the service user to transmit its PDU to the peer service user entity. For example, the DL-DATA parameter data contains layer 3 information. The PH-DATA parameter data contains the data link layer frame.

NOTE – The operations across the data link layer/layer 3 boundary shall be such that the layer sending a primitive can assume a temporal order of the bits within the parameter data and that the layer receiving the primitive can reconstruct the information with its assumed temporal order.

4.2 Primitive procedures

4.2.1 General

Primitive procedures specify the interactions between adjacent layers to invoke and provide a service. The service primitives represent the elements of the procedures.

In the scope of this Recommendation the interactions between layer 3 and the data link layer are specified.

4.2.2 Layer 3 – Data link layer interactions

The states of a data link connection endpoint may be derived from the internal states of the data link layer entity supporting this type of a data link connection.

Data link connection endpoint states are defined as follows:

- a) Broadcast data link connection endpoint:
 - *information transfer state.*
- b) Point-to-point data link connection endpoint:
 - *link connection released state;*
 - *awaiting establish state;*
 - *awaiting release state;*
 - *link connection established state.*

The primitives provide the procedural means to specify conceptually how a data link service user can invoke a service.

This subclause defines the constraints on the sequences in which the primitives may occur. The sequences are related to the states at one point-to-point data link connection endpoint.

The possible overall sequences of primitives at a point-to-point data link connection endpoint are defined in the state transition diagram, Figure 8. The *link connection released* and *link connection established* states are stable states whilst the *awaiting establish* and *awaiting release* states are transition states.

The model illustrates the behaviour of layer 2 as seen by layer 3. This model assumes that the primitives passed between layers are implemented using a first in first out queue. In this model,

"collisions" of request and indication primitives can occur thereby illustrating actions that seem to be in conflict with the actual layer 2 protocol description. In some implementations, these collisions could occur.

4.3 Block interaction diagram of the data link layer

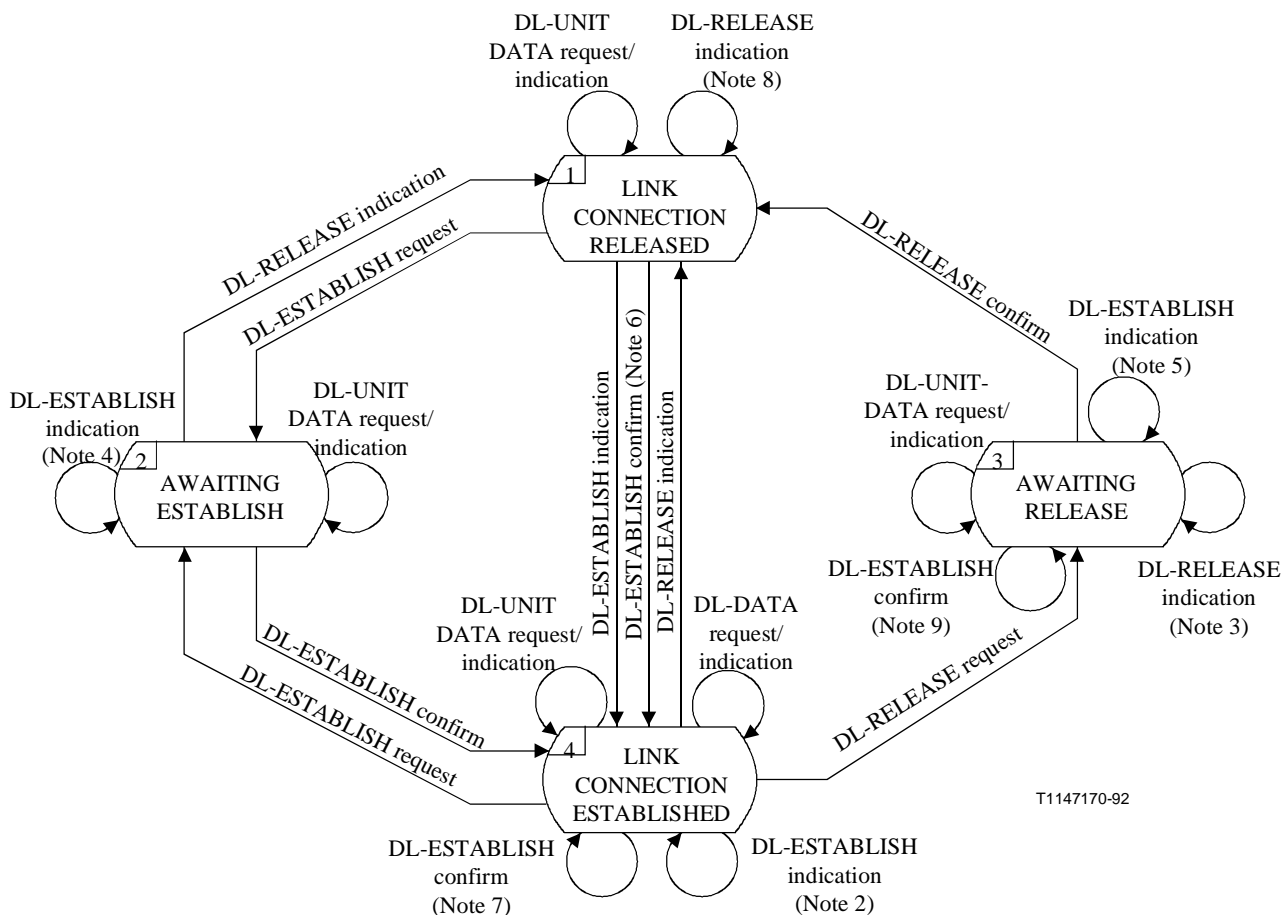
Subclause 4.1 defines the primitives associated with this Recommendation and 4.2 defines the primitive procedures between layer 3 and the data link layer.

Subclause 5.4/Q.920 [1] provides a functional block diagram which includes the functional blocks relevant to the data link layer.

This subclause clarifies how the primitives defined in this Recommendation apply to the various functional blocks.

A block interaction diagram relates the service primitives to these functional blocks which have to interact, see Figure 9. Additional signals are needed for the internal use within the data link layer for the communication between point-to-point link procedures or broadcast link procedures, respectively, and the multiplex procedure.

The Figure 9 is an aid to illustrate the relationship between various functional blocks. It is not intended to constrain implementation. The primitives contained in Figure 9 are those defined in 4.1. Other additional primitives may be defined in other Recommendations, e.g. dealing with maintenance requirements.



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NOTE 1 – If the data link layer entity issues a DL-ESTABLISH indication (this applies to the case of data link layer initiated or peer system initiated re-establishment), DL-RELEASE confirm or DL-RELEASE indication, this indicates the discard of all the data link service data units representing DL-DATA requests.

NOTE 2 – This primitive notifies layer 3 of link re-establishment.

NOTE 3 – This primitive will occur if a DL-RELEASE request collides with a DL-RELEASE indication.

NOTE 4 – This primitive will occur if a DL-ESTABLISH request collides with a DL-ESTABLISH indication.

NOTE 5 – This primitive will occur if a DL-RELEASE request collides with a DL-ESTABLISH indication.

NOTE 6 – This primitive will occur if a DL-ESTABLISH request (this applies to the case of layer 3 initiated re-establishment) collides with a DL-RELEASE indication. Since this DL-RELEASE indication is not related to the DL-ESTABLISH request, the data link layer will establish the link and issue a DL-ESTABLISH confirm. It may also occur if establishment was initiated upon receipt of an unsolicited DM response with the F bit set to 0.

NOTE 7 – This primitive will occur as a result of multiple collisions of primitives. If a first DL-ESTABLISH request collides with a DL-RELEASE indication, the data link layer will establish the link and issue a DL-ESTABLISH confirm (see Note 6). This DL-ESTABLISH confirm (it is related to the first DL-ESTABLISH request) would collide with a subsequent DL-ESTABLISH request which may be issued since layer 3 is not aware that the DL-RELEASE indication was not related to the first DL-ESTABLISH request. Since layer 3 relates this DL-ESTABLISH confirm to the subsequent DL-ESTABLISH request it assumes that the data link layer is in the link connection established state, but the data link layer will re-establish the link and issue again a DL-ESTABLISH confirm.

NOTE 8 – This primitive will occur if a DL-ESTABLISH request (this applies to the case of layer 3 initiated re-establishment) collides with a DL-RELEASE indication. Since this DL-RELEASE indication is not related to the DL-ESTABLISH request, the data link layer will try to establish the link and if this is not possible, it issues a DL-RELEASE indication.

NOTE 9 – This primitive will occur as a result of multiple collisions of primitives. If a first DL-ESTABLISH request collides with a DL-RELEASE indication, the data link layer will establish the link and issue a DL-ESTABLISH confirm (see Note 6). This DL-ESTABLISH confirm may collide with a subsequent DL-ESTABLISH request and the data link layer will re-establish the link and issue again a DL-ESTABLISH confirm (see Note 7). This second DL-ESTABLISH confirm (it is related to the second DL-ESTABLISH request) may collide with a subsequent DL-RELEASE request which may be issued since layer 3 is not aware that the DL-RELEASE indication was not related to the first DL-ESTABLISH request. Since layer 3 relates this first DL-ESTABLISH confirm to the subsequent DL-ESTABLISH request it assumes the data link layer in the link connection established state, but the data link layer will re-establish the link and issue again a DL-ESTABLISH confirm (see Note 7).

Figure 8/Q.921 – State Transition Diagram for sequences of Primitives at a point-to-point Data Link Connection as seen by Layer 3 (Note 1)

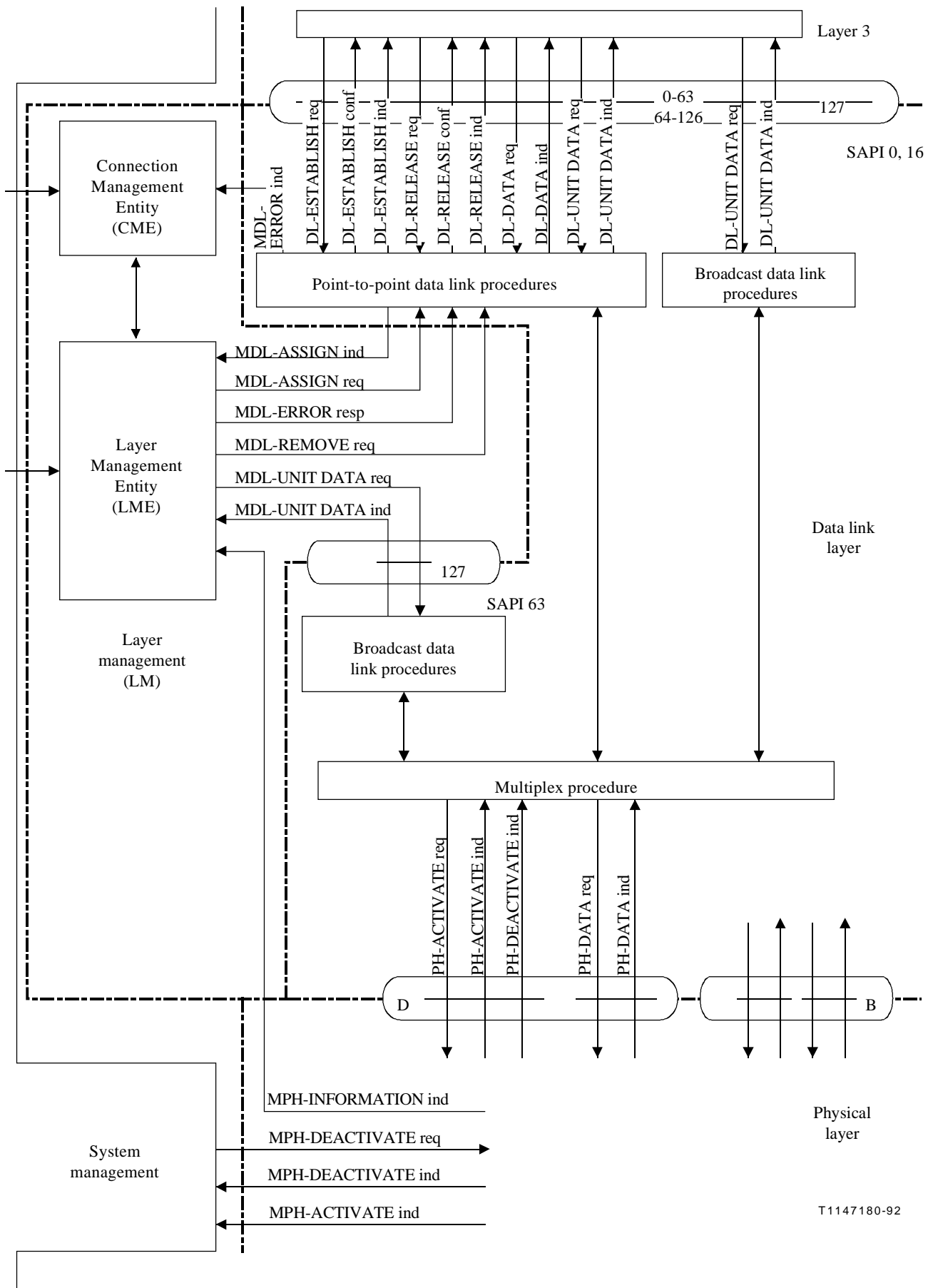


Figure 9/Q.921 – Block interaction diagram

5 Definition of the peer-to-peer procedures of the data link layer

The procedures for use by the data link layer are specified in the following subclauses.

The elements of procedure (frame types) which apply are:

- a) for unacknowledged information transfer (see 5.2):
UI-command;
- b) for multiple frame acknowledged information transfer (see 5.5 to 5.8):
SABME-command;
UA-response;
DM-response;
DISC-command;
RR-command/response;
RNR-command/response;
REJ-command/response;
I-command;
FRMR-response (see Note);

NOTE – An FRMR-response shall not be generated by a data link layer entity; however, on receipt of this frame actions according to 5.8.6 shall be taken.

- c) for connection management entity information transfer:
XID-command/response.

5.1 Procedure for the use of the P/F bit

5.1.1 Unacknowledged information transfer

For unacknowledged information transfer the P/F bit is not used and shall be set to 0.

5.1.2 Acknowledged multiple frame information transfer

A data link layer entity receiving an SABME, DISC, RR, RNR, REJ or I frame, with the P bit set to 1, shall set the F bit to 1 in the next response frame it transmits, as defined in Table 7.

Table 7/Q.921 – Immediate response operation of P/F bit

Command received with P bit = 1	Response transmitted with F bit = 1
SABME, DISC	UA, DM
I, RR, RNR, REJ	RR, RNR, REJ (Note)
NOTE – A LAPB data link layer entity may transmit an FRMR or DM response with the F bit set to 1 in response to an I frame or supervisory command with the P bit set to 1.	

5.2 Procedures for unacknowledged information transfer

5.2.1 General

The procedures which apply to the transmission of information in unacknowledged operation are defined in 5.2.2 and 5.2.3.

No data link layer error recovery procedures are defined for unacknowledged operation.

If an implementation conforming to this Recommendation is used only in conjunction with a defined set of applications which do not use the unacknowledged information transfer service, and an open interface is not provided to user applications, e.g. for later inclusion of a new application, then, as an implementation option:

- a) the layer 2 protocol entity supporting any particular user application which does not use unacknowledged information transfer service, does not need to implement the generation of the UI frame, and may treat received UI frames as frames associated with an application not implemented according to 3.6.1 of;
- b) the layer 2 protocol entity supporting layer 2 management (broadcast data link procedures) does not need to be implemented, if no peer-to-peer layer management is provided (see Annex A).

5.2.2 Transmission of unacknowledged information

NOTE 1 – The term "transmission of a UI frame" refers to the delivery of a UI frame by the data link layer to the physical layer.

SDUs to be conveyed by means of unacknowledged information transfer are passed to the data link layer by layer 3 or management entities using the primitives DL-UNIT DATA request or MDL-UNIT DATA request, respectively. The SDUs passed by layer 3 or layer management shall be transmitted in a UI command frame.

For broadcast operation, the TEI value in the UI command address field shall be set to 127 (binary 111 1111, the group value).

For point-to-point operation, the appropriate TEI value shall be used.

The P bit shall be set to 0.

In the case of persistent layer 1 deactivation, the data link layer will be informed by an appropriate indication. Upon receipt of this indication, all UI transmission queues shall be discarded. At the network side, the system management entity provides that the PH-DEACTIVATE indication primitive will be issued only if persistent deactivation has occurred. However, at the user side, the conditions to issue a PH-DEACTIVATE indication primitive depend on the implementation of the physical layer.

NOTE 2 – The network side system management deactivation procedures should ensure that layer 1 is not deactivated before all UI data transfer is completed.

5.2.3 Receipt of unacknowledged information

On receipt of a UI command frame with a SAPI which is supported by the receiver and TEI which has been assigned to the receiver, the contents of the information field shall be passed to the layer 3 or management entity using the data link layer to layer 3 primitive DL-UNIT DATA indication or the data link layer to management primitive MDL-UNIT DATA indication, respectively. Otherwise, the UI command frame shall be discarded.

5.3 Terminal Endpoint Identifier (TEI) management procedures

5.3.1 General

This subclause defines the TEI management protocols for TEI values to be used for point-to-point data link connections (TEI value is in the range from 0 through 126). In particular, this subclause is not applicable to the management of the group TEI (TEI = 127).

TEI management is based on the following procedural means:

- TEI assignment procedures (see 5.3.2);

- TEI check procedures (see 5.3.3);
- TEI removal procedures (see 5.3.4);
- optional user equipment initiated TEI Identity verify procedures (see 5.3.5).

A user equipment in the *TEI-unassigned* state shall use the TEI assignment procedures to enter the *TEI-assigned* state. Conceptually, these procedures exist in the layer management entity. The layer management entity on the network side is referred to as the Assignment Source Point (ASP) in this Recommendation.

The purpose of these procedures is to:

- a) allow automatic TEI equipment to request the network to assign a TEI value that the data link layer entities within the requesting user equipment will use in their subsequent communications;
- b) allow a network to remove a previously assigned TEI value from specific or all user equipment;
- c) allow a network to check:
 - whether or not a TEI value is in use; or
 - whether duplicate TEI assignment has occurred;
- d) allow user equipment the option to request that the network invoke TEI check procedures.

The user side layer management entity shall instruct the user data link layer entities to remove all TEI values when it is notified that the terminal is disconnected at the interface (as defined in Recommendation I.430).

Additionally, the user side layer management entity should instruct the user data link layer entity to remove a TEI value for its own internal reasons (for example, losing the ability to communicate with the network). The layer management entity shall use the MDL-REMOVE request primitive for these purposes.

Subclause 5.3.4.1 includes the actions taken by a data link layer entity receiving an MDL-REMOVE request primitive.

Typically, one TEI value would be used by the user equipment (for example, a data link layer entity which has been assigned a TEI value could use that value for all SAPs which it supports). If required, a number of TEI values may be requested by multiple use of the procedures defined in 5.3.2. It shall be the responsibility of the user to maintain the association between TEI and SAPI values.

The initiation of TEI assignment procedures occurs on the receipt of a request for establishment or unacknowledged information transfer while in the TEI-unassigned state. The data link layer entity shall inform the layer management entity using the MDL-ASSIGN indication primitive. Alternatively, the user side layer management entity may initiate the TEI assignment procedures for its own reasons.

NOTE – In the case of initialization from a no power condition, the user equipment should postpone the start of the TEI assignment procedure until a layer 2 service that needs a TEI is to be provided.

All layer management entity PDUs used for these TEI management procedures are transmitted to, or received from, the data link layer entity in the form of SDUs using the MDL-UNIT DATA request primitive, or the MDL-UNIT DATA indication primitive, respectively. The data link layer entity shall transmit SDUs for the support of management procedures in UI command frames. The SAPI value shall be 63. The TEI value shall be 127.

5.3.2 TEI assignment procedure

If the user equipment is of the non-automatic TEI assignment category, the user side layer management entity shall deliver the TEI value to be used to the data link layer entity(s) via the MDL-ASSIGN request primitive.

If the user equipment is of the automatic TEI assignment category, upon initiation of the automatic TEI assignment procedure, the user side layer management entity shall transmit to its peer a message containing the following elements:

- a) message type = Identity request;
- b) Reference number (Ri); and
- c) Action indicator (Ai).

The Reference number, Ri, shall be used to differentiate between a number of user equipment which may simultaneously request assignment of a TEI value. The Ri shall be 2 octets in length and shall be randomly generated for each request message by the user equipment.

All values in the range 0 to 65535 shall be available from the random number generator.

NOTE – The design of the random number generator should minimize the probability of identical reference numbers being generated by terminals which initiate their TEI assignment procedures simultaneously. However, there exists a small probability that double assignment will occur. Possible procedures to resolve this problem are listed in 5.3.3 to 5.3.5.

The single-octet Action indicator, Ai, shall be used to indicate a request to the ASP for the assignment of any TEI value available.

The coding of the Ai shall be Ai = 127. This Ai value requests the ASP to assign any TEI value.

A timer T202 shall be started.

The ASP, on receipt of the Identity request message, shall either:

- select a TEI value;
- deny Identity requests with Ai values in the range 64-126, or ignore Identity requests with the Ai value in the range 0-63; or
- ignore the Identity request message if a previous Identity request message that contains an identical Ri has been received and no response has been issued. In this case, the ASP shall not assign a TEI value to either request.

Selection of a TEI value shall be on the basis of information stored at the ASP. This may consist of:

- a map of the full range of automatic TEI values; or
- an updated list of all automatic TEI values available for assignment, or a smaller subset.

The ASP, after having selected the TEI value, shall inform the network data link layer entities by means of the MDL-ASSIGN request primitive and transmit to its peer a message containing the following elements:

- a) message type = Identity assigned;
- b) Reference number (Ri); and
- c) the assigned TEI value in the Ai field.

If the available TEI information/resources are exhausted, a TEI check procedure should be initiated.

A user side layer management entity receiving this Identity assigned message shall compare the TEI value in the Ai field with its own TEI value(s) (if any) to see if it is already allocated if an Identity

request message is outstanding. Additionally, the TEI value in the Ai field may be compared with its TEI(s) on the receipt of all Identity assigned messages.

If there is a match, the management entity shall either:

- initiate TEI removal; or
- initiate the TEI identity verify procedures.

If there is no match, the user side layer management entity shall:

- compare the Ri value with any outstanding Identity request message and if it matches, consider the TEI value assigned to the user equipment, discard the value of Ri, inform the user side data link layer entities by means of the MDL-ASSIGN request primitive and stop timer T202;
- compare the Ri value with any outstanding Identity request message and if there is no match, do nothing;
- if there is no outstanding Identity request message, do nothing.

When the data link layer receives the MDL-ASSIGN request primitive from the layer management entity, the data link layer entity shall:

- enter the TEI-assigned state; and
- proceed with data link establishment procedures if a DL-ESTABLISH request primitive is outstanding, or proceed with the transmission of a UI command frame if a DL-UNIT DATA request primitive is outstanding.

To deny an Identity request message, the ASP shall transmit to its peer a message containing the following elements:

- a) message type = Identity denied;
- b) Reference number (Ri); and
- c) the value of TEI which is denied in the Ai field (a value of 127 indicates that no TEI values are available).

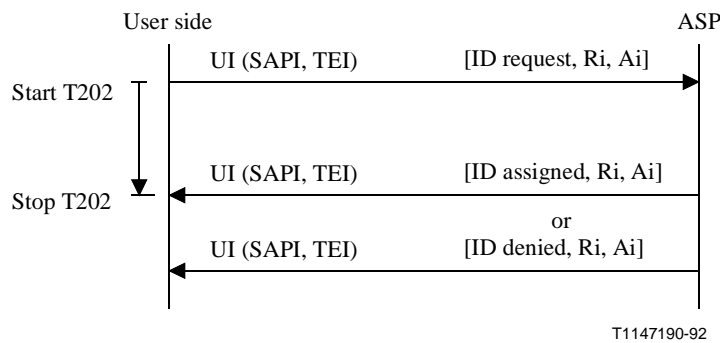
5.3.2.1 Expiry of timer T202

If the user receives either no response or an Identity denied message to its Identity request message, then on expiry of timer T202, the timer shall be restarted and the Identity request message shall be retransmitted with a new value of Ri.

After N202 unsuccessful attempts to acquire a TEI value, the layer management entity shall inform the data link layer entity using the MDL-ERROR response primitive. The data link layer entity receiving the MDL-ERROR response primitive shall respond with the DL-RELEASE indication primitive if a request for establishment had previously occurred, and shall discard all unserved DL-UNIT DATA request primitives.

The values of T202 and N202 are specified in 5.9.

The TEI assignment procedure is illustrated in Figure 10.



ID Identity
 SAPI Service Access Point Identifier = 63
 TEI Group TEI = 127
 Ai Action indicator (see Table 8)
 Ri Reference number
 () Contents of the data link layer command address field
 [] Contents of the data link layer command information field

Figure 10/Q.921 – TEI assignment procedure

5.3.3 TEI check procedure

5.3.3.1 Use of the TEI check procedure

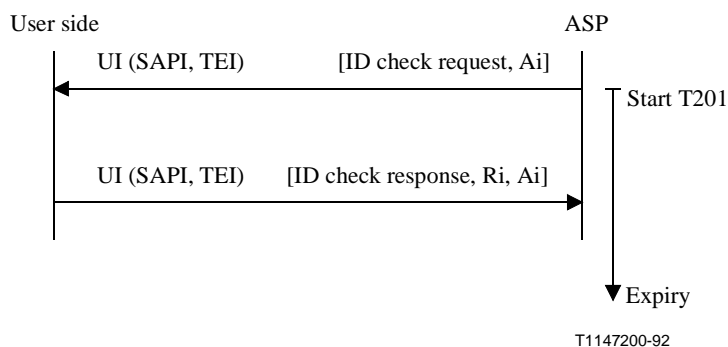
The TEI check procedure shall be used in the TEI audit and recovery procedures. The TEI check procedure allows the network side layer management entity either:

- to establish that a TEI value is in use; or
- to verify duplicate TEI assignment.

The TEI check procedure for verifying duplicate TEI assignment may also optionally be invoked as a response to an Identity verify request message from the user equipment.

5.3.3.2 Operation of the TEI check procedure

The TEI check procedure is illustrated in Figure 11.



ID Identity
 SAPI SAPI = 63
 TEI TEI = 127
 Ai Action indicator (see Table 8)
 Ri Reference number
 () Contents of the data link layer command address field
 [] Contents of the data link layer command information field

Figure 11/Q.921 – TEI check procedure

The ASP shall transmit a message containing the following elements:

- a) message type = Identity check request; and
- b) Ai field which contains the TEI value to be checked or the value 127 when all TEI values are to be checked.

Timer T201 shall be started.

If any user equipment has been assigned the TEI value specified in the identity check request message, it shall respond by transmitting a message containing the following elements:

- a) message type = Identity check response;
- b) the TEI value in the Ai field; and
- c) Reference number (Ri).

NOTE – The randomly-generated Ri is present in the Identity check response message to ensure that in the case where more than one user equipment happens to commence transmission of the Identity check response message at precisely the same time (i.e. the first "0" bit of the opening flag coincides) due to different Ri values a collision at layer 1 (see ISDN user-network interfaces; Recommendation I.430 [5] for clarification) occurs. The resolution of this collision results in multiple Identity check response messages.

When the TEI check procedure is used to verify duplicate TEI assignment:

- if more than one identity check response message with the Ai field indicating identical TEI values is received within the T201 time period, then duplicate TEI assignment shall be considered present; otherwise the request shall be repeated once and timer T201 restarted;
- if more than one Identity check response message with the Ai field indicating identical TEI values is received within the second T201 time period, duplicate TEI assignment shall be considered present;
- if no Identity check response message is received after both T201 periods, the TEI value shall be assumed to be free and available for (re)assignment;
- if one Identity check response message is received in one or both T201 periods, the TEI value shall be assumed to be in use.

When the TEI check procedure is used to test whether a TEI value is in use, it is completed upon the receipt of the first TEI Identity check response message, and the TEI value is assumed to be in use. Otherwise:

- if no Identity check response message is received within T201, the identity check request message shall be repeated once and timer T201 restarted;
- if no Identity check response message is received after the second Identity check request message, the TEI value shall be assumed to be free and available for reassignment.

If the Ai value in the Identity check request message is equal to 127, it is preferred that the receiving user side layer management entity respond with a single Identity check response message that contains all of the TEI values in use within that user equipment (see 5.3.6.5). If an Identity check request message with Ai equal to 127 is transmitted and an Identity check response message is received making use of the extension facility, each Ai variable in the Ai field shall be processed as if received in separate Identity check response messages for parallel Identity check request messages.

5.3.4 TEI removal procedure

When the network side layer management entity determines that the removal of a TEI value (see 5.3.4.2) is necessary, the ASP shall transmit a message containing the following elements and issue an MDL-REMOVE request primitive:

- a) message type = Identity remove; and

- b) TEI value which is to be removed, as indicated in the Ai field (the value 127 indicates that all user equipments are requested to remove their TEI values; otherwise, the specific TEI value is requested to be removed); and
- c) Layer management entity identifier.

The Identity remove message shall be sent twice in succession, to overcome possible message loss.

When the user side layer management entity determines that the removal of a TEI value is necessary (see 5.3.4.2), it shall instruct the data link layer entity to enter the TEI-unassigned state, using the MDL-REMOVE request primitive.

Further action to be taken shall be:

- a) if automatic TEI values apply, initiation of automatic TEI assignment for a new TEI value; or
- b) if non-automatic TEI values apply, notification to the equipment user for the need for corrective action.

NOTE – In point-to-point configurations, unless there exists a pre-arrangement between the user side and the network side for the TEI value to be used, the user side shall initiate the link layer establishment procedures in order to allow the network side to offer calls.

5.3.4.1 Action taken by the data link layer entity receiving the MDL-REMOVE request primitive

A data link layer entity receiving an MDL-REMOVE request primitive shall:

- a) if no DL-RELEASE request primitive is outstanding and the user equipment is not in the TEI-assigned state, issue a DL-RELEASE indication primitive; or
- b) if a DL-RELEASE request primitive is outstanding, issue a DL-RELEASE confirm primitive.

The data link layer entity shall then enter the *TEI-unassigned* state after discarding the contents of both UI and I queues.

5.3.4.2 Conditions for TEI removal

At the user equipment, automatic TEI values shall be removed under the following conditions:

- a) on request from the ASP by an Identity remove message;
- b) on receipt of an MPH-INFORMATION indication (disconnected) primitive; or
- c) on receipt of an Identity assigned message containing a TEI value in the Ai field, depending upon whether or not an Identity request message is outstanding. If an Identity request message is outstanding and the TEI received in the Ai field is already in use within the user equipment (see 5.3.2), then the user equipment shall either remove the TEI value or invoke the TEI Identity verify procedures. If an Identity request message is not outstanding, then the user equipment shall either do nothing or check if the TEI received in the Ai field is already in use within the user equipment. If the TEI is in use, the user equipment shall either remove the TEI value or invoke the TEI Identity verify procedure.

In addition to the conditions identified above, on receipt of an MDL-ERROR indication primitive indicating that the data link layer entity has assumed possible duplicate assignment of a TEI value, the user side equipment shall:

- i) either remove the concerned TEI value, as the preferred action; or
- ii) request a TEI check procedure by the transmission of an Identity verify request message and proceed according to 5.3.5.

At the user equipment, non-automatic TEI values can be removed on request from the ASP by an Identity remove message. On receipt of an MDL-ERROR indication primitive indicating that the data link layer entity has assumed possible duplicate assignment of a TEI value, the concerned TEI can be removed in preference to requesting a TEI check procedure by the transmission of an Identity verify request message. If the Identity verify option is chosen, the user equipment shall proceed according to 5.3.5. When the TEI value is removed, an appropriate indication shall be made to the user.

At the network side, automatic TEI values should be removed under the following conditions:

- a) following a TEI audit procedure showing that a TEI value is no longer in use;
- b) following a TEI audit procedure showing that duplicate TEI assignment has occurred; or
- c) on receipt of an MDL-ERROR indication primitive indicating a possible duplicate TEI assignment, of which confirmation can be requested by the invocation of the TEI check procedures. If the TEI check procedure was invoked, the network side shall proceed according to 5.3.3 in order to evaluate if the TEI value indicated in the MDL-ERROR indication is "free", "single" or "duplicate". The appropriate action for the verdict "free" is remove TEI, for "duplicate" it is remove TEI locally and initiate TEI removal procedures while for "single" no action is taken.

At the network side, non-automatic TEI values should be removed under the following conditions:

- a) following a TEI audit procedure showing that duplicate TEI assignment has occurred; or
- b) on receipt of an MDL-ERROR indication primitive indicating a possible duplicate TEI assignment, of which confirmation can be requested by the invocation of the TEI check procedures. If the TEI check procedure was invoked, the network side shall proceed according to 5.3.3 in order to evaluate if the TEI value indicated in the MDL-ERROR indication is "free", "single" or "duplicate". The appropriate action for the verdict "free" is remove TEI, for "duplicate" it is remove TEI locally and initiate TEI removal procedures while for "single" no action is taken.

5.3.5 TEI identity verify procedure

5.3.5.1 General

The TEI identity verify procedure allows the user side layer management entity to have the capability to request that the network invoke the identity check procedure for verification of duplicate TEI assignment.

The TEI identity verify procedure is optional for both the network and user equipment.

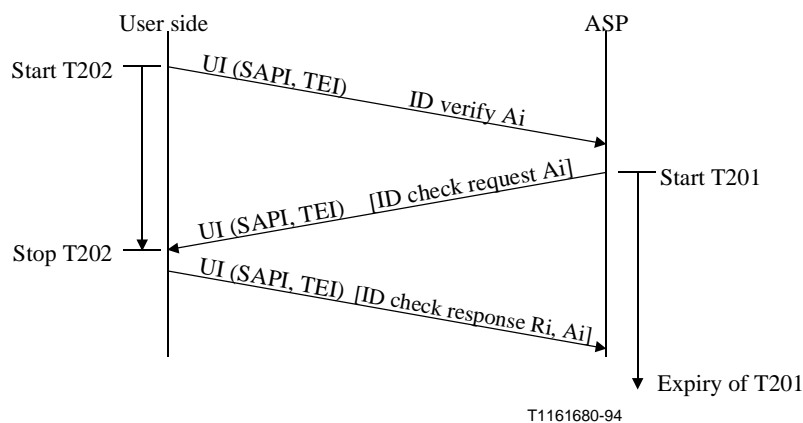
5.3.5.2 Operation of the TEI identity verify procedure

The TEI identity verify procedure is illustrated in Figure 12.

The user equipment shall transmit an Identity verify message containing the following elements:

- a) message type = Identity verify request;
- b) the TEI value to be checked in the Ai field; and
- c) the Ri field, which is not processed by the network and is coded 0.

Timer T202 is started.



SAPI	SAPI = 63
TEI	TEI = 127
ID	Identity
Ai	Action indicator (see Table 8)
Ri	Reference number
()	Contents of the data link layer command address field
[]	Contents of the data link layer command information field

NOTE – The Ai in the ID verify message will be in the range 0 to 126. Ai = 127 is not allowed.

Figure 12/Q.921 – TEI identify verify procedure

The ASP, on receipt of the TEI Identity verify message shall, if implemented, invoke the TEI check procedure as defined in 5.3.3. This will result in the ASP sending an Identity check request message to the user equipment.

The user side layer management entity receives an Identity check request message with the contents of the Ai field equal to its TEI value (for which verification has been requested) or the value 127 (indicating that all TEI values are to be checked), it shall stop timer T202. In any case, it shall respond to an Identity check request message according to the TEI check procedure as defined in 5.3.3

5.3.5.3 Expiry of Timer T202

If the user equipment receives no Identity check request message with an Ai equal to its TEI or an Ai equal to 127 before the expiry of timer T202, the user side layer management entity shall restart the timer and the TEI Identity verify message shall be retransmitted. If no Identity check request message is received from the ASP after the second TEI Identity verify request message, the TEI shall be removed.

5.3.6 Formats and codes

5.3.6.1 General

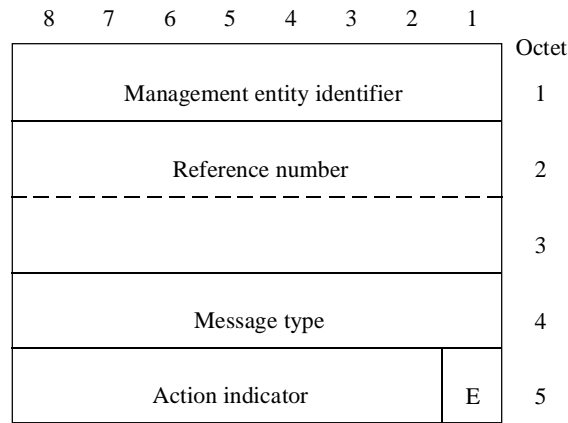
All messages used for TEI management procedures are carried in the information field of UI command frames with a SAPI value set to 63 (binary 11 1111) and TEI value set to 127 (binary 111 1111).

All messages have the structure shown in Figure 13.

Fields that are not used in a specific message are coded all zeros, and are not to be processed by either side.

The coding of each field for the various messages is specified in Table 8.

E is the Action indicator field extension bit (see 5.3.6.5).



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Figure 13/Q.921 – Message used for TEI management procedures

Table 8/Q.921 – Codes for messages concerning TEI management procedures

Message name	Management entity identifier	Reference number Ri	Message type	Action indicator Ai
Identity request (user-to-network)	0000 1111	0-65535	0000 0001	Ai = 127, Any TEI value acceptable
Identity assigned (network-to-user)	0000 1111	0-65535	0000 0010	Ai = 64-126, Assigned TEI value
Identity denied (network-to-user)	0000 1111	0-65535	0000 0011	Ai = 64-126, Denied TEI value
				Ai = 127, No TEI value available
Identity check request (network-to-user)	0000 1111	Not used (coded 0)	0000 0100	Ai = 127, Check all TEI values
				Ai = 0-126, TEI value to be checked
Identity check response (user-to-network)	0000 1111	0-65535	0000 0101	Ai = 0-126, TEI value in use
Identity remove (network-to-user)	0000 1111	Not used (coded 0)	0000 0110	Ai = 127, Request for removal of all TEI values
				Ai = 0-126, TEI value to be removed
Identity verify (user-to-network)	0000 1111	Not used (coded 0)	0000 0111	Ai = 0-126, TEI value to be checked

5.3.6.2 Layer management entity identifier

For TEI administration procedures, the layer management entity identifier octet is 0000 1111. Other values are reserved for further standardization.

5.3.6.3 Reference number (Ri)

Octets 2 and 3 contain Ri. When used, it can assume any value between 0 and 65535.

5.3.6.4 Message type

Octet 4 contains the message type. The purpose of the message type is to identify the function of the message being sent.

5.3.6.5 Action indicator (Ai)

The Ai field is extended by reserving the first transmitted bit of the Ai field octets to indicate the final octet of the Ai field.

Ai variables in the Ai field are coded as follows:

- a) bit 1 is the extension bit and is coded as follows:
 - 0 to indicate an extension (see Note), and
 - 1 to indicate the final octet;
- b) bits 2 to 8 contain the Action indicator.

The purpose of the Action indicator is to identify the concerned TEI value(s).

NOTE – The use of the extension mechanism is confined to the Identity Check Response when all of the TEI values in use within a user equipment are to be reported in a single Identity Check Response upon receipt of an Identity Check Request with an Ai equal to 127 (see 5.3.3.2).

5.4 Initialization of data link layer parameters

5.4.1 General

Each data link layer entity has an associated data link connection management entity. The data link connection management entity has the responsibility for initializing the link parameters necessary for correct peer-to-peer information transport.

The method of initialization of the parameters follows one of the two methods below:

- initialization to the default values as specified in 5.9; or
- initialization based on the values supplied by its peer entity (automatic negotiation of data link layer parameters).

Typically, after the assignment of a TEI value to the management entity, the data link connection management entity is notified by its layer management entity that parameter initialization is required.

After parameter initialization, the data link connection management entity will notify the layer management entity that parameter initialization has occurred, and the layer management entity will issue the MDL-ASSIGN request primitive.

5.4.2 Parameter initialization

The parameter initialization procedure may invoke either the internal initialization procedure or the automatic negotiation of data link parameter procedure.

5.4.2.1 Internal parameter initialization

When the layer management entity notifies the connection management entity of TEI assignment, the connection management entity shall initialize the link parameters to the default values and notify the layer management of task completion.

5.4.2.2 Automatic negotiation of data link layer parameter values

The procedures for automatic negotiation of data link layer parameters are described in Appendix IV.

5.5 Procedures for establishment and release of multiple frame operation

5.5.1 Establishment of multiple frame operation

5.5.1.1 General

These procedures shall be used to establish multiple frame operation between the network and a designated user entity.

Layer 3 will request establishment of the multiple frame operation by the use of the DL-ESTABLISH request primitive. Re-establishment may be initiated as a result of the data link layer procedures defined in 5.7. All frames other than unnumbered frame formats received during the establishment procedures shall be ignored.

5.5.1.2 Establishment procedures

A data link layer entity shall initiate a request for the multiple frame operation to be set by transmitting the SABME command. All existing exception conditions shall be cleared, the retransmission counter shall be reset, and timer T200 shall then be started (timer T200 is defined in 5.9.1). All mode setting commands shall be transmitted with the P bit set to 1.

Layer 3 initiated establishment procedures imply the discard of all outstanding DL-DATA request primitives and all I frames in queue.

A data link layer entity receiving an SABME command, if it is able to enter the multiple-frame-established state, shall:

- a) respond with a UA response with the F bit set to the same binary value as the P bit in the received SABME command;
- b) set V(S), V(R) and V(A) to 0;
- c) enter the *multiple-frame-established* state and inform layer 3 using the DL-ESTABLISH indication primitive;
- d) clear all existing exception conditions;
- e) clear any existing peer receiver busy condition; and
- f) start timer T203 (timer T203 is defined in 5.9.8), if implemented.

If the data link layer entity is unable to enter the *multiple-frame-established* state, it shall respond to the SABME command with a DM response with the F bit set to the same binary value as the P bit in the received SABME command.

Upon reception of the UA response with the F bit set to 1, the originator of the SABME command shall:

- reset timer T200;
- start timer T203, if implemented;
- set V(S), V(R), and V(A) to 0; and
- enter the *multiple-frame-established* state and inform layer 3 using the DL-ESTABLISH confirm primitive.

Upon reception of a DM response with the F bit set to 1, the originator of the SABME command shall indicate this to layer 3 by means of the DL-RELEASE indication primitive, and reset timer T200. It shall then enter the *TEI-assigned* state. DM responses with the F bit set to 0 shall be ignored in this case.

A DL-RELEASE request primitive received during data link layer initiated re-establishment shall be serviced on completion of the establishment mode-setting operation.

5.5.1.3 Procedure on expiry of timer T200

If timer T200 expires before the UA or DM response with the F bit set to 1 is received, the data link layer entity shall:

- retransmit the SABME command as above;
- restart timer T200; and
- increment the retransmission counter.

After retransmission of the SABME command N200 times, the data link layer entity shall indicate this to layer 3 and the connection management entity by means of the DL-RELEASE indication and MDL-ERROR indication primitives, respectively, and enter the *TEI-assigned* state, after discarding all outstanding DL-DATA request primitives and all I frames in queue.

The value of N200 is defined in 5.9.2.

5.5.2 Information transfer

Having either transmitted the UA response to a received SABME command or received the UA response to a transmitted SABME command, I frames and supervisory frames shall be transmitted and received according to the procedures described in 5.6.

If an SABME command is received while in the *multiple-frame-established* state, the data link layer entity shall conform to the re-establishment procedure described in 5.7.

On receipt of a UI command, the procedures defined in 5.2 shall be followed.

5.5.3 Termination of multiple frame operation

5.5.3.1 General

These procedures shall be used to terminate the multiple frame operation between the network and a designated user entity.

Layer 3 will request termination of the multiple frame operation by use of the DL-RELEASE request primitive.

All frames other than unnumbered frames received during the release procedures shall be ignored.

All outstanding DL-DATA request primitives and all I frames in queue shall be discarded.

In the case of persistent layer 1 deactivation the data link layer entity shall discard all I queues and deliver to layer 3 a DL-RELEASE confirm primitive if a DL-RELEASE request primitive is outstanding, or otherwise a DL-RELEASE indication primitive. At the network side, the system management entity provides that the PH-DEACTIVATE indication primitive will be issued only, if persistent deactivation has occurred. At the user side, however, the conditions to issue a PH-DEACTIVATE indication primitive depend on the implementation of the physical layer.

5.5.3.2 Release procedure

A data link layer entity shall initiate a request for release of the multiple frame operation by transmitting the Disconnect (DISC) command with the P bit set to 1. Timer T200 shall then be started and the retransmission counter reset.

A data link layer entity receiving a DISC command while in the *multiple-frame-established* or *timer recovery* state shall transmit a UA response with the F bit set to the same binary value as the P bit in the received DISC command. A DL-RELEASE indication primitive shall be passed to layer 3, and the *TEI-assigned* state shall be entered.

If the originator of the DISC command receives either:

- a UA response with the F bit set to 1; or
- a DM response with the F bit set to 1, indicating that the peer data link layer entity is already in the *TEI-assigned* state,

it shall enter the *TEI-assigned* state and reset timer T200.

The data link layer entity which issued the DISC command is now in the *TEI-assigned* state and will notify layer 3 by means of the DL-RELEASE confirm primitive. The conditions relating to this state are defined in 5.5.4.

5.5.3.3 Procedure on expiry of timer T200

If timer T200 expires before a UA or DM response with the F bit set to 1 is received, the originator of the DISC command shall:

- retransmit the DISC command as defined in 5.5.3.2;
- restart timer T200; and
- increment the retransmission counter.

If the data link layer entity has not received the correct response as defined in 5.5.3.2, after N200 attempts to recover, the data link layer entity shall indicate this to the connection management entity by means of the MDL-ERROR indication primitive, enter the *TEI-assigned* state and notify layer 3 by means of the DL-RELEASE confirm primitive.

5.5.4 TEI-assigned state

While in the TEI-assigned state:

- the receipt of a DISC command shall result in the transmission of a DM response with the F bit set to the value of the received P bit;
- on receipt of an SABME command, the procedures defined in 5.5.1 shall be followed;
- on receipt of an unsolicited DM response with the F bit set to 0, the data link layer entity shall, if it is able to, initiate the establishment procedures by the transmission of an SABME (see 5.5.1.2). Otherwise, the DM shall be ignored;
- on receipt of UI commands, the procedures defined in 5.2 shall be followed;
- on receipt of any unsolicited UA response an MDL-ERROR indication primitive indicating a possible duplicate assignment of a TEI value shall be issued; and
- all other frame types shall be discarded.

5.5.5 Collision of unnumbered commands and responses

5.5.5.1 Identical transmitted and received commands

If the transmitted and received unnumbered commands (SABME or DISC) are the same, the data link layer entities shall send the UA response at the earliest possible opportunity. The indicated state shall be entered after receiving the UA response. The data link layer entity shall notify layer 3 by means of the appropriate confirm primitive.

5.5.5.2 Different transmitted and received commands

If the transmitted and received unnumbered commands (SABME or DISC) are different, the data link layer entities shall issue a DM response at the earliest possible opportunity. Upon receipt of a DM response with the F bit set to 1, the data link layer shall enter the *TEI-assigned* state and notify layer 3 by means of the appropriate primitive. The entity receiving the DISC command will issue a

DL-RELEASE indication primitive, while the other entity will issue a DL-RELEASE confirm primitive.

5.5.6 Unsolicited DM response and SABME or DISC command

When a DM response with the F bit set to 0 is received by a data link layer entity, a collision between a transmitted SABME or DISC command and the unsolicited DM response may have occurred. This is typically caused by a user equipment applying a protocol procedure according to X.25 LAPB [7] to ask for a mode-setting command.

In order to avoid misinterpretation of the DM response received, a data link layer entity shall always send its SABME or DISC command with the P bit set to 1.

A DM response with the F bit set to 0 colliding with an SABME or DISC command shall be ignored.

5.6 Procedures for information transfer in multiple frame operation

The procedures which apply to the transmission of I frames are defined below.

NOTE – The term "transmission of an I frame" refers to the delivery of an I frame by the data link layer to the physical layer.

5.6.1 Transmitting I frames

Information received by the data link layer entity from layer 3 by means of a DL-DATA request primitive shall be transmitted in an I frame. The control field parameters N(S) and N(R) shall be assigned the values of V(S) and V(R), respectively. V(S) shall be incremented by 1 at the end of the transmission of the I frame.

If timer T200 is not running at the time of transmission of an I frame, it shall be started. If time T200 expires, the procedures defined in 5.6.7 shall be followed.

If V(S) is equal to V(A) plus k (where k is the maximum number of outstanding I frames – see 5.9.5), the data link layer entity shall not transmit any new I frames, but may retransmit an I frame as a result of the error recovery procedures as described in 5.6.4 and 5.6.7.

When the network side or user side is in the own receiver busy condition, it may still transmit I frames, provided that a peer receiver busy condition does not exist.

NOTE – Any DL-DATA request primitives received whilst in the timer recovery condition shall be queued.

5.6.2 Receiving I frames

Independent of a timer recovery condition, when a data link layer entity is not in an own receiver busy condition and receives a valid I frame whose N(S) is equal to the current V(R), the data link layer entity shall:

- pass the information field of this frame to layer 3 using the DL-DATA indication primitive;
- increment by 1 its V(R) and act as indicated below.

5.6.2.1 P bit set to 1

If the P bit of the received I frame was set to 1, the data link layer entity shall respond to its peer in one of the following ways:

- if the data link layer entity receiving the I frame is still not in an own receiver busy condition, it shall send an RR response with the F bit set to 1;
- if the data link layer entity receiving the I frame enters the own receiver busy condition upon receipt of the I frame, it shall send an RNR response with the F bit set to 1.

5.6.2.2 P bit set to 0

If the P bit of the received I frame was set to 0 and:

- a) if the data link layer entity is still not in an own receiver busy condition:
 - if no I frame is available for transmission or if an I frame is available for transmission but a peer receiver busy condition exists, the data link layer entity shall transmit an RR response with the F bit set to 0; or
 - if an I frame is available for transmission and no peer receiver busy condition exists, the data link layer entity shall transmit the I frame with the value of N(R) set to the current value of V(R) as defined in 5.6.1; or
- b) if, on receipt of this I frame, the data link layer entity is now in an own receiver busy condition, it shall transmit an RNR response with the F bit set to 0.

When the data link layer entity is in an own receiver busy condition, it shall process any received I frame according to 5.6.6.

5.6.3 Sending and receiving acknowledgements

5.6.3.1 Sending acknowledgements

Whenever a data link layer entity transmits an I frame or a supervisory frame, N(R) shall be set equal to V(R).

5.6.3.2 Receiving acknowledgements

On receipt of a valid I frame or supervisory frame (RR, RNR, or REJ), even in the own receiver busy, or timer recovery conditions, the data link layer entity shall treat the N(R) contained in this frame as an acknowledgement for all the I frames it has transmitted with an N(S) up to and including the received N(R) – 1. V(A) shall be set to N(R). When not in the timer recovery condition the data link layer entity shall reset the timer T200 on receipt of a valid I frame or supervisory frame with the N(R) higher than V(A) (actually acknowledging some I frames), or an REJ frame with an N(R) equal to V(A).

NOTE 1 – If a supervisory frame or an I frame with the P bit set to 1 has been transmitted and not acknowledged by a supervisory frame response with the F bit set to 1, timer T200 shall not be reset.

NOTE 2 – Upon receipt of a valid I frame, timer T200 shall not be reset if the data link layer entity is in the peer receiver busy condition.

If timer T200 has been reset by the receipt of an I, RR, or RNR frame, and if there are outstanding I frames still unacknowledged, the data link layer entity shall restart timer T200. If timer T200 then expires, the data link layer entity shall follow the recovery procedure as defined in 5.6.7 with respect to the unacknowledged I frames.

If timer T200 has been reset by the receipt of an REJ frame, the data link layer entity shall follow the retransmission procedures in 5.6.4.

5.6.4 Receiving REJ frames

On receipt of a valid REJ frame, the data link layer entity shall act as follows:

- a) if it is not in the timer recovery condition:
 - clear an existing peer receiver busy condition;
 - set its V(S) and its V(A) to the value of the N(R) contained in the REJ frame control field;
 - stop timer T200;

- start timer T203, if implemented;
 - if it was an REJ command frame with the P bit set to 1, transmit an appropriate supervisory response frame (see Note 2 in 5.6.5) with the F bit set to 1.
 - transmit the corresponding I frame as soon as possible, as defined in 5.6.1, taking into account the items 1) to 3) below and the paragraph following items 1) to 3) and
 - notify a protocol violation to the connection management entity by means of the MDL-ERROR indication primitive, if it was an REJ response frame with the F bit set to 1.
- b) if it is in the timer recovery condition and it was an REJ response frame with the F bit set to 1:
- clear an existing peer receiver busy condition;
 - set its V(S) and its V(A) to the value N(R) contained in the REJ frame control field;
 - stop timer T200;
 - start timer T203, if implemented;
 - enter the multiple-frame-established state; and
 - transmit the corresponding I frame as soon as possible, as defined in 5.6.1, taking into account the items 1) to 3) below and the paragraph following items 1) to 3).
- c) if it is in the timer recovery condition and it was an REJ frame other than an REJ response frame with the F bit set to 1:
- clear an existing peer receiver busy condition;
 - set its V(A) to the value of the N(R) contained in the REJ frame control field; and
 - if it was an REJ command frame with the P bit set to 1, transmit an appropriate supervisory response frame with the F bit set to 1 (see Note 2 in 5.6.5).

Transmission of I frames shall take account of the following:

- 1) if the data link layer entity is transmitting a supervisory frame when it receives the REJ frame, it shall complete that transmission before commencing transmission of the requested I frame;
- 2) if the data link layer entity is transmitting an SABME command, a DISC command, a UA response or a DM response when it receives the REJ frame, it shall ignore the request for retransmission; and
- 3) if the data link layer entity is not transmitting a frame when the REJ is received, it shall immediately commence transmission of the requested I frame.

All outstanding unacknowledged I frames, commencing with the I frame identified in the received REJ frame, shall be transmitted. Other I frames not yet transmitted may be transmitted following the retransmitted I frames.

5.6.5 Receiving RNR frames

After receiving a valid RNR command or response, if the data link layer entity is not engaged in a mode-setting operation, it shall set a peer receiver busy condition and then:

- if it was an RNR command with the P bit set to 1, it shall respond with an RR response with the F bit set to 1 if the data link layer entity is not in an own receiver busy condition, and shall respond with an RNR response with the F bit set to 1 if the data link layer entity is in an own receiver busy condition; and
- if it was an RNR response with the F bit set to 1, an existing timer recovery condition shall be cleared and the N(R) contained in this RNR response shall be used to update V(S).

The data link layer entity shall take note of the peer receiver busy condition and not transmit any I frames to the peer which has indicated the busy condition.

NOTE 1 – The N(R) received in any RR or RNR command frame (irrespective of the setting of the P bit) will not be used to update the V(S).

The data link layer entity shall then:

- treat the N(R) contained in the received RNR frame as an acknowledgement for all the I frames that have been (re)transmitted with an N(S) up to and including N(R) – 1, and set its V(A) to the value of the N(R) contained in the RNR frame; and
- restart timer T200 unless a supervisory response frame with the F bit set to 1 is still expected.

If timer T200 expires, the data link layer entity shall:

- i) if it is not yet in a timer recovery condition, enter the timer recovery condition and reset the retransmission count variable; or
- ii) if it is already in a timer recovery condition, continue as indicated below.

The data link layer entity shall then:

- a) if the value of the retransmission count variable is less than N200:
 - transmit an appropriate supervisory command (see Note 2) with a P bit set to 1;
 - restart timer T200;
 - add one to its retransmission count variable; and
- b) if the value of the retransmission count variable is equal to N200, initiate a re-establishment procedure as defined in 5.7, and indicate this by means of the MDL-ERROR indication primitive to the connection management entity.

The data link layer entity receiving the supervisory frame with the P bit set to 1 shall respond, at the earliest opportunity, with an appropriate supervisory response frame (see Note 2) with the F bit set to 1, to indicate whether or not its own receiver busy condition still exists.

Upon receipt of the supervisory response with the F bit set to 1, the data link layer entity shall reset timer T200, and:

- if the response is an RR or REJ response, the peer receiver busy condition is cleared and the data link layer entity may transmit new I frames or retransmit I frames as defined in 5.6.1 or 5.6.4, respectively; or
- if the response is an RNR response, the data link layer entity receiving the response shall proceed according to this 5.6.5, first paragraph.

If a supervisory command (RR, RNR, or REJ) with the P bit set to 0 or 1, or a supervisory response frame (RR, RNR, or REJ) with the F bit set to 0 is received during the enquiry process, the data link layer entity shall:

- if the supervisory frame is an RR or REJ command frame or an RR or REJ response frame with the F bit set to 0, clear the peer receiver busy condition and if the supervisory frame received was a command with the P bit set to 1, transmit the appropriate supervisory response frame (see Note 2) with the F bit set to 1. However, the transmission or retransmission of I frames shall not be undertaken until the appropriate supervisory response frame with the F bit set to 1 is received or until expiry of timer T200; or
- if the supervisory frame is an RNR command frame or an RNR response frame with the F bit set to 0, retain the peer receiver busy condition and if the supervisory frame received was

an RNR command with the P bit set to 1, transmit the appropriate supervisory response frame (see Note 2) with the F bit set to 1.

Upon receipt of an SABME command, the data link layer entity shall clear the peer receiver busy condition.

NOTE 2 – The appropriate supervisory frame for the circumstances indicated is defined below:

- a) If the data link layer entity is not in an own receiver busy condition and is in a Reject exception condition [that is, an N(S) sequence error has been received, and an REJ frame has been transmitted, but the requested I frame has not been received], the appropriate supervisory frame is the RR frame.
- b) If the data link layer entity is not in an own receiver busy condition but is in an N(S) sequence error exception condition [that is, an N(S) sequence error has been received but an REJ frame has not been transmitted], the appropriate supervisory frame is the REJ frame.
- c) If the data link layer entity is in its own receiver busy condition, the appropriate supervisory frame is the RNR frame.
- d) Otherwise, the appropriate supervisory frame is the RR frame.

5.6.6 Data link layer own receiver busy condition

When the data link layer entity enters an own receiver busy condition, it shall transmit an RNR frame at the earliest opportunity.

The RNR frame may be either:

- an RNR response with the F bit set to 0; or
- if this condition is entered on receiving a command frame with the P bit set to 1, an RNR response with the F bit set to 1; or
- if this condition is entered on expiry of timer T200, an RNR command with the P bit set to 1.

All received I frames with the P bit set to 0 shall be discarded, after updating V(A).

All received supervisory frames with the P/F bit set to 0 shall be processed, including updating V(A).

All received I frames with the P bit set to 1 shall be discarded, after updating V(A). However, an RNR response frame with the F bit set to 1 shall be transmitted.

All received supervisory frames with the P bit set to 1 shall be processed including updating V(A). An RNR response with the F bit set to 1 shall be transmitted.

To indicate to the peer data link layer entity the clearance of the own receiver busy condition, the data link layer entity shall transmit an RR frame or, if a previously detected N(S) sequence error has not yet been reported, an REJ frame with the N(R) set to the current value of V(R).

The transmission of an SABME command or a UA response (in reply to an SABME command) also indicates to the peer data link layer entity the clearance of the own receiver busy condition.

5.6.7 Waiting acknowledgement

The data link layer entity shall maintain an internal retransmission count variable.

If timer T200 expires, the data link layer entity shall:

- if it is not yet in the timer recovery condition, enter the timer recovery condition and reset the retransmission count variable; or
- if it is already in the timer recovery condition, continue as indicated below.

The data link layer entity shall then:

- a) if the value of the retransmission count variable is less than N200:
 - add one to its retransmission count variable; and
 - restart timer T200; and either
 - transmit an appropriate supervisory command (see Note 2 in 5.6.5) with the P bit set to 1; or
 - retransmit the last transmitted I frame $[V(S) - 1]$ with the P bit set to 1; or
- b) if the value of the retransmission count variable is equal to N200, initiate a re-establishment procedure as defined in 5.7 and indicate this by means of the MDL-ERROR indication primitive to the connection management entity.

The following paragraph applies only for a data link layer which is in the timer recovery condition since the case of receiving acknowledgement in the multiple frame established state is described in 5.6.3.2.

The timer recovery condition is cleared only if the data link layer entity receives a valid supervisory frame response with the F bit set to 1. If the N(R) of this received supervisory frame is within the range from its current V(A) to its current V(S) inclusive, it shall set its V(S) to the value of the received N(R). Timer T200 shall be reset if the received supervisory frame response is an RR or REJ response with the F bit set to 1. The data link layer entity shall resume then with I frame transmission or retransmission, as appropriate. Timer T200 shall be reset and restarted if the received supervisory response is an RNR response with the F bit set to 1, to proceed with the enquiry process according to 5.6.5.

5.7 Re-establishment of multiple frame operation

5.7.1 Criteria for re-establishment

The criteria for re-establishing the multiple frame mode of operation are defined in this subclause by the following conditions:

- the receipt, while in the multiple-frame mode of operation, of an SABME;
- the receipt of a DL-ESTABLISH request primitive from layer 3 (see 5.5.1.1);
- the occurrence of N200 retransmission failures while in the timer recovery condition (see 5.6.7);
- the occurrence of a frame rejection condition as identified in 5.8.5;
- the receipt, while in the multiple-frame mode of operation, of an FRMR response frame (see 5.8.6);
- the receipt, while in the multiple-frame mode of operation, of an unsolicited DM response with the F bit set to 0 (see 5.8.7);
- the receipt, while in the timer-recovery condition, of a DM response with the F bit set to 1.

5.7.2 Procedures

In all re-establishment situations, the data link layer entity shall follow the procedures defined in 5.5.1. All locally-generated conditions for re-establishment will cause the transmission of the SABME.

In the case of data link layer and peer initiated re-establishment, the data link layer entity shall also:

- issue an MDL-ERROR indication primitive to the connection management entity; and

- if $V(S) > V(A)$ prior to re-establishment, issue a DL-ESTABLISH indication primitive to layer 3, and discard all I queues.

In case of layer 3 initiated re-establishment or if a DL-ESTABLISH request primitive occurs pending re-establishment, the DL-ESTABLISH confirm primitive shall be used.

5.8 Exception condition reporting and recovery

Exception conditions may occur as the result of physical layer errors or data link layer procedural errors.

The error recovery procedures which are available to effect recovery following the detection of an exception condition at the data link layer are defined in this subclause.

The actions to be taken by the connection management entity on receipt of an MDL-ERROR indication primitive are defined in Appendix II.

5.8.1 N(S) sequence error

An N(S) sequence error exception condition occurs in the receiver when a valid I frame is received which contains an N(S) value which is not equal to the V(R) at the receiver. The information field of all I frames whose N(S) does not equal V(R) shall be discarded.

The receiver shall not acknowledge [nor increment its V(R)] the I frame causing the sequence error, nor any I frames which may follow, until an I frame with the correct N(S) is received.

A data link layer entity which receives one or more I frames having sequence errors but otherwise error-free, or subsequent supervisory frames (RR, RNR, and REJ), shall use the control field information contained in the N(R) field and the P or F bit to perform data link control functions; for example, to receive acknowledgement of previously transmitted I frames and to cause the data link layer entity to respond if the P bit is set to 1. Therefore, the retransmitted I frame may contain an N(R) field value and P bit that are updated from, and therefore different from, the ones contained in the originally transmitted I frame.

The REJ frame is used by a receiving data link layer entity to initiate an exception condition recovery (retransmission) following the detection of an N(S) sequence error.

Only one REJ exception condition for a given direction of information transfer shall be established at a time.

A data link layer entity receiving an REJ command or response shall initiate sequential transmission (retransmission) of I frames starting with the I frame indicated by the N(R) contained in the REJ frame.

An REJ exception condition is cleared when the requested I frame is received or when an SABME or DISC command is received.

An optional procedure for the retransmission of an REJ response frame is described in Appendix I.

5.8.2 N(R) sequence error

An N(R) sequence error exception condition occurs in the transmitter when a valid supervisory frame or I frame is received which contains an invalid N(R) value.

A valid N(R) is one that is in the range $V(A) \leq N(R) \leq V(S)$.

The information field contained in an I frame which is correct in sequence and format may be delivered to layer 3 by means of the DL-DATA indication primitive.

The data link layer entity shall inform the connection management entity of this exception condition by means of the MDL-ERROR indication primitive, and initiate re-establishment according to 5.7.2.

5.8.3 Timer recovery condition

If a data link layer entity, due to a transmission error, does not receive a single I frame or the last I frame(s) in a sequence of I frames, it will not detect an out-of-sequence exception condition and therefore will not transmit an REJ frame.

The data link layer entity which transmitted the unacknowledged I frame(s) shall, on the expiry of timer T200, take appropriate recovery action as defined in 5.6.7 to determine at which I frame retransmission must begin.

5.8.4 Invalid frame condition

Any frame received which is invalid (as defined in 2.9) shall be discarded, and no action shall be taken as a result of that frame.

5.8.5 Frame rejection condition

A frame rejection condition results from one of the following conditions:

- a) the receipt of an undefined frame (see 3.6.1, third paragraph);
- b) the receipt of a supervisory or unnumbered frame with incorrect length;
- c) the receipt of an invalid N(R); or
- d) the receipt of a frame with an information field which exceeds the maximum established length.

Upon occurrence of a frame rejection condition whilst in the multiple frame operation, the data link layer entity shall:

- issue an MDL-ERROR indication primitive; and
- initiate re-establishment (see 5.7.2).

Upon occurrence of a frame rejection condition during establishment or release from multiple frame operation, or whilst a data link is not established, the data link layer entity shall:

- issue an MDL-ERROR indication primitive; and
- discard the frame.

NOTE – For satisfactory operation it is essential that a receiver is able to discriminate between invalid frames, as defined in 2.9, and frames with an I-field which exceeds the maximum established length [see d) of 3.6.11]. An unbounded frame may be assumed, and thus discarded, if two times the longest permissible frame plus two octets are received without a flag detection.

5.8.6 Receipt of an FRMR response frame

Upon receipt of an FRMR response frame in the multiple-frame mode of operation, the data link layer entity shall:

- issue an MDL-ERROR indication primitive; and
- initiate re-establishment (see 5.7.2).

5.8.7 Unsolicited response frames

The action to be taken on the receipt of an unsolicited response frame is defined in Table 9.

The data link layer entity shall assume possible duplicate-TEI assignment on the receipt of an unsolicited UA response and shall inform layer management.

5.8.8 Duplicate assignment of a TEI value

A data link layer entity shall assume duplicate assignment of a TEI value and initiate recovery as specified below by:

- a) the receipt of a UA response frame whilst in the *multiple-frame-established* state;
- b) the receipt of a UA response frame whilst in the *timer recovery* state;
- c) the receipt of a UA response frame whilst in the *TEI-assigned* state.

Table 9/Q.921 – Actions taken on receipt of unsolicited response frames

Unsolicited response frame	TEI assigned	Awaiting Establishment	Awaiting Release	Multiple frame modes of operation	
				Established mode	Time recovery condition
UA response F = 1	MDL-ERROR indication	Solicited	Solicited	MDL-ERROR indication	MDL-ERROR indication
UA response F = 0	MDL-ERROR indication	MDL-ERROR indication	MDL-ERROR indication	MDL-ERROR indication	MDL-ERROR indication
DM response F = 1	Ignore	Solicited	Solicited	MDL-ERROR indication	Re-establish MDL-ERROR indication
DM response F = 0	Establish	Ignore	Ignore	Re-establish MDL-ERROR indication	Re-establish MDL-ERROR indication
Supervisory response F = 1	Ignore	Ignore	Ignore	MDL-ERROR indication	Solicited
Supervisory response F = 0	Ignore	Ignore	Ignore	Solicited	Solicited

A data link layer entity, after assuming duplicate assignment of a TEI value shall inform the connection management entity by means of the MDL-ERROR indication primitive.

5.9 List of system parameters

The system parameters listed below are associated with each individual SAP.

A method of assigning these parameters is defined in 5.4.

The term default implies that the value defined should be used in the absence of any assignment or negotiation of alternative values.

5.9.1 Timer T200

The default value for timer T200 at the end of which transmission of a frame may be initiated according to the procedures described in 5.6 shall be one second.

NOTE 1 – The proper operation of the procedure requires that timer T200 be greater than the maximum time between transmission of command frames and the reception of their corresponding response or acknowledgement frames.

NOTE 2 – When an implementation includes multiple terminals on the user side together with a satellite connection in the transmission path, a value of T200 greater than 1 second may be necessary. A value of 2.5 seconds is suggested.

NOTE 3 – In certain digital sections (e.g. involving satellites), the default value of timer T200 may be too small to ensure proper operation. To accommodate such configurations, it is recommended that user and network equipment allow selection of alternate values of timer T200 or implement the automatic negotiation of data link parameters procedures of Appendix IV.

5.9.2 Maximum number of retransmissions (N200)

The maximum number of retransmissions of a frame (N200) is a system parameter. The default value of N200 shall be 3.

5.9.3 Maximum number of octets in an information field (N201)

The maximum number of octets in an information field (N201) is a system parameter. (See also 2.5.)

- For an SAP supporting signalling, the default value shall be 260 octets.
- For SAPs supporting packet information, the default value shall be 260 octets.

5.9.4 Maximum number of transmissions of the TEI Identity request message (N202)

The maximum number of transmissions of a TEI Identity request message (when the user requests a TEI) is a system parameter. The default value of N202 shall be 3.

5.9.5 Maximum number of outstanding I frames (k)

The maximum number (*k*) of sequentially numbered I frames that may be outstanding (that is, unacknowledged) at any given time is a system parameter which shall not exceed 127, for extended (modulo 128) operation.

- For an SAP supporting basic access (16 kbit/s) signalling, the default value shall be 1.
- For an SAP supporting primary rate (64 kbit/s) signalling, the default value shall be 7.
- For an SAP supporting basic access (16 kbit/s) packet information, the default value shall be 3.
- For an SAP supporting primary rate (64 kbit/s) packet information, the default value shall be 7.

NOTE – In certain digital sections (e.g. involving satellites), for the 64 kbit/s D-channel, the value of *k* may not be large enough to assure efficient operation. To accommodate such configurations, it is recommended that user and network equipment allow selection of alternative values of *k*, or implement the data link layer parameter negotiation procedures of Appendix IV.

5.9.6 Timer T201

The minimum time between retransmission of the TEI Identity check messages (T201) is a system parameter which shall be set to T200 seconds.

5.9.7 Timer T202

The minimum time between the transmission of TEI Identity request messages is a system parameter (T202) which shall be set to 2 seconds.

5.9.8 Timer T203

Timer T203 represents the maximum time allowed without frames being exchanged. The default value of timer T203 shall be 10 seconds.

Table 10 provides an overview of these system parameters by depicting which procedures, link types and user or network side data link layer entities use them and by indicating the recommended default or fixed values, respectively.

5.10 Data link layer monitor function

5.10.1 General

The procedural elements defined in clause 5 allow for the supervision of the data link layer resource. This subclause describes procedures which may be used to provide this supervision function. The use of this function is optional.

5.10.2 Data link layer supervision in the multiple-frame-established state

The procedures specified herein propose a solution which is already identified in the HDLC classes of procedures. The connection verification is a service provided by data link layer to layer 3. This implies that layer 3 is informed in case of a failure only. Furthermore, the procedure may be incorporated in the "normal" exchange of information and may become more efficient than a procedure based on the involvement of layer 3.

Table 10/Q.921 – System parameters

	k	T200	T201	T202	T203	N200	N201	N202
Point-to-point data link procedure on a D-channel at 16 kbit/s	1	1 s	Not applicable	Not applicable	10 s	3	260	Not applicable
	3	1 s	Not applicable	Not applicable	10 s	3	260	Not applicable
Point-to-point data link procedure on a D-channel at 64 kbit/s	7	1 s	Not applicable	Not applicable	10 s	3	260	Not applicable
	7	1 s	Not applicable	Not applicable	10 s	3	260	Not applicable
TEI assignment procedure (SAPI = 63)	Not applicable	Not applicable	Not applicable	2 s	Not applicable	Not applicable	Not applicable	3
	Not applicable	Not applicable	1 s	Not applicable	Not applicable	Not applicable	Not applicable	Not applicable

The procedure is based on supervisory command frames (RR command, RNR command) and timer T203, and operates in the multiple-frame-established state as follows.

If there are no frames being exchanged on the data link connection (neither new nor outstanding I frames, nor supervisory frames with a P bit set to 1), there is no means to detect a faulty data link connection condition, or a user equipment having been unplugged. Timer T203 represents the maximum time allowed without frames being exchanged.

If timer T203 expires, a supervisory command with a P bit set to 1 is transmitted. Such a procedure is protected against transmission errors by making use of the normal timer T200 procedure including retransmission count and N200 attempts.

5.10.3 Connection verification procedures

5.10.3.1 Start timer T203

The timer T203 is started:

- when the *multiple-frame-established* state is entered; and
- in the *multiple-frame-established* state whenever timer T200 is stopped. (See Note in 5.10.3.2.)

Upon receiving an I or supervisory frame, timer T203 will be restarted if timer T200 is not to be started.

5.10.3.2 Stop timer T203

The timer T203 is stopped:

- when, in the *multiple-frame-established* state, the timer T200 is started (see Note); and
- upon leaving the *multiple-frame-established* state.

NOTE – These two conditions mean that timer T203 is only started whenever timer T200 is stopped and not restarted.

5.10.3.3 Expiry of timer T203

If timer T203 expires, the data link layer entity will act as follows (it should be noted that timer T200 is neither running nor expired):

- a) set the retransmission count variable to 0;
- b) enter *timer recovery* state;
- c) transmit a supervisory command with the P bit set to 1 as follows:
 - if there is not a receiver busy condition (own receiver not busy), transmit an RR command; or
 - if there is a receiver busy condition (own receiver busy), transmit an RNR command; and
- d) start timer T200; and
- e) send MDL-ERROR indication primitive to connection management after N200 retransmissions.

ANNEX A

Provision of point-to-point signalling connections

In certain applications it may be advantageous to have a single point-to-point signalling connection at layer 3. This implementation option uses the value 0 for the TEI. Use of the value 0 in such applications does not preclude using that value in other applications or networks.

In certain networks, for both user and network sides, an arrangement is permitted for point-to-point configurations using for signalling a single point-to-point data link connection only. In such a configuration, the following requirements are defined:

- a) support of signalling:
 - the TEI value 0 shall be used in combination with SAPI 0; and
 - two peer-to-peer layer 3 signalling entities shall communicate over a single point-to-point data link connection within the SAP identified by the SAPI value 0, making use of the acknowledged information transfer service provided by layer 2;
- b) peer-to-peer layer 2 management procedures shall not be used, independent of the number of SAPs in use, if:
 - one, and not more than one, TEI is used; and
 - the same TEI is used for all SAPs;otherwise, peer-to-peer layer 2 management procedures according to 5.3 shall be used.

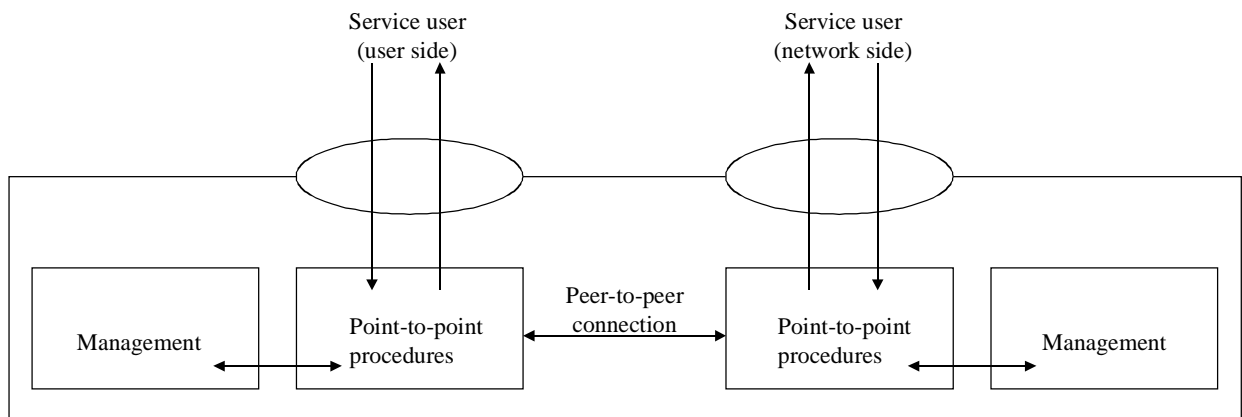
ANNEX B

SDL for point-to-point procedures

B.1 General

The purpose of this Annex is to provide one example of an SDL representation of the point-to-point procedures of the data link layer, to assist in the understanding of this Recommendation. This representation does not describe all of the possible actions of the data link layer entity, as a non-partitioned representation was selected in order to minimize its complexity. The SDL representation does not therefore constrain implementations from exploiting the full scope of the procedures as presented within the text of this Recommendation. The text description of the procedures is definitive.

The representation is a peer-to-peer model of the point-to-point procedures of the data link layer and is applicable to the data link layer entities at both the user and network sides for all ranges of TEI values. See Figure B.1.



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Figure B.1/Q.921 – Peer-to-peer of the point-to-point procedures

B.2 An overview of the states of the point-to-point data link layer entity

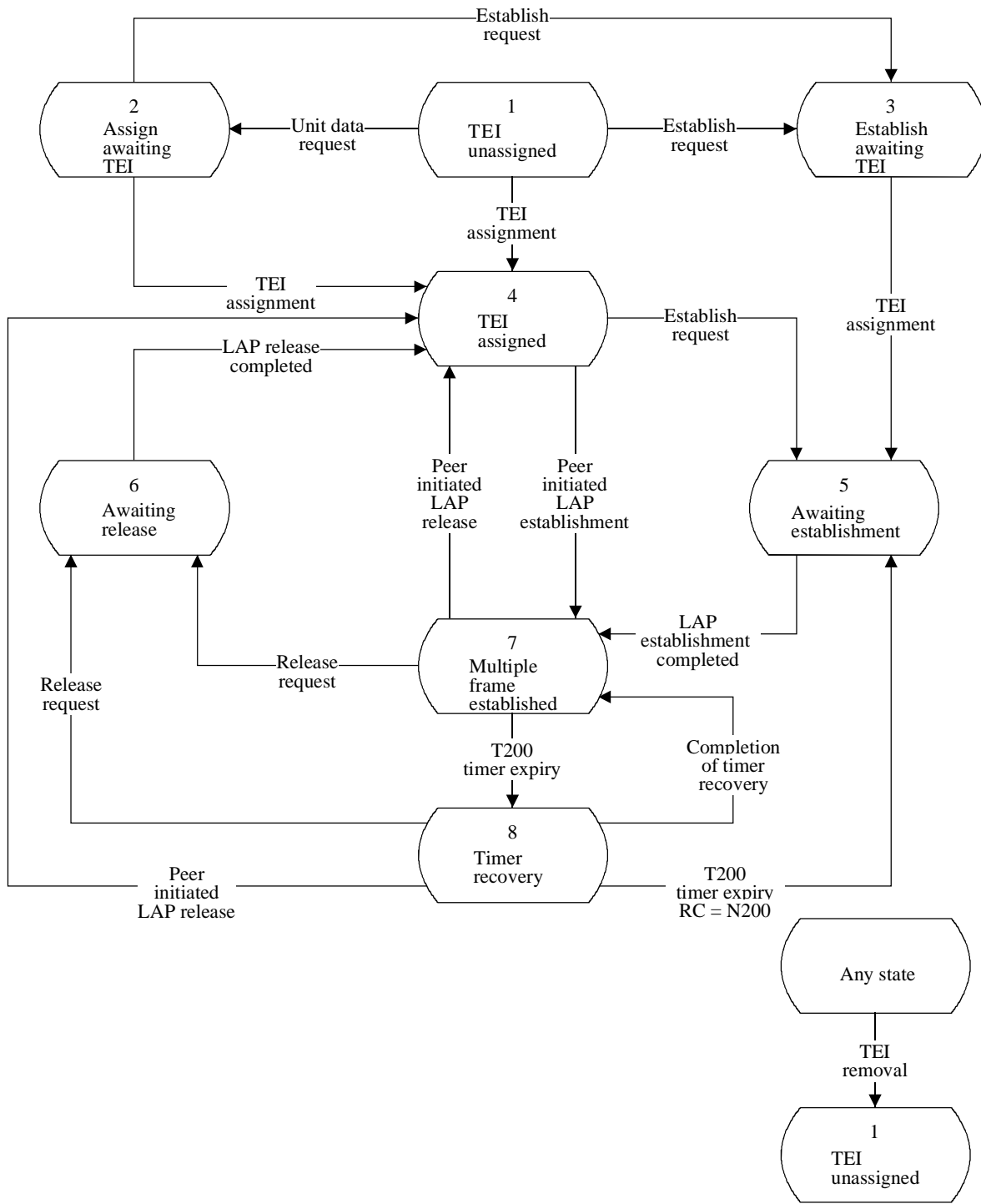
The SDL representation of the point-to-point procedures are based on an expansion of the three basic states identified in 3.4.2/Q.920 [1] to the following 8 states:

- State 1: *TEI unassigned.*
- State 2: *Assign awaiting TEI.*
- State 3: *Establish awaiting TEI.*
- State 4: *TEI assigned.*
- State 5: *Awaiting establishment.*
- State 6: *Awaiting release.*
- State 7: *Multiple frame established.*
- State 8: *Timer recovery.*

An overview of the inter-relationship of these states is provided in Figure B.2. This overview is incomplete, and serves only as an introduction to the SDL representation. All data link layer entities are conceptually initiated in the *TEI unassigned* state (state 1), and will interact with the layer management in order to request a TEI value. TEI assignment initiated by a Unit data request will cause the data link layer entity to move to the *TEI assigned* state (state 4) via the *assign awaiting TEI* state (state 2). Initiation by an Establishment request will cause a transition to the *awaiting establishment* state (state 5) via the *establish awaiting TEI* state (state 3). Direct TEI assignment will cause an immediate transition to the *TEI assigned* state (state 4). In states 4-8, Unit data requests can be directly serviced by the data link layer entity. The receipt of an Establish request in the *TEI assigned* state (state 4) will cause the initiation of the establishment procedures and the transition to the *awaiting establishment* state (state 5). Completion of the LAP establishment procedures takes the data link layer entity into the *multiple frame established* state (state 7). Peer initiated establishment causes a direct transition from the *TEI assigned* state (state 4) to the *multiple frame established* state (state 7). In the *multiple frame established* state (state 7), Acknowledged data transfer requests can be serviced directly subject to the restrictions of the procedures. Expiry of timer T200, which is used in both the flow control and data transfer aspects of the data link layer entity's procedures, initiates the transition to the *timer recovery* state (state 8). Completion of the timer recovery procedures will return the data link layer entity to the *multiple frame established* state (state 7). In states 7 and 8 of the SDL representation, the following conditions which are identified within this Recommendation are observed:

- a) peer receiver busy;
- b) reject exception;
- c) own receiver busy.

In addition, other conditions are used in order to avoid identification of additional states. The complete combination of both of these categories of conditions with the 8 states of the SDL representation is the basis for the state transition table description of the data link layer entity. A peer initiated LAP release will take the data link layer entity directly into the *TEI assigned* state (state 4), whilst a Release request will be via the *awaiting release* state (state 6). TEI removal will cause a transition to the *TEI unassigned* state (state 1).

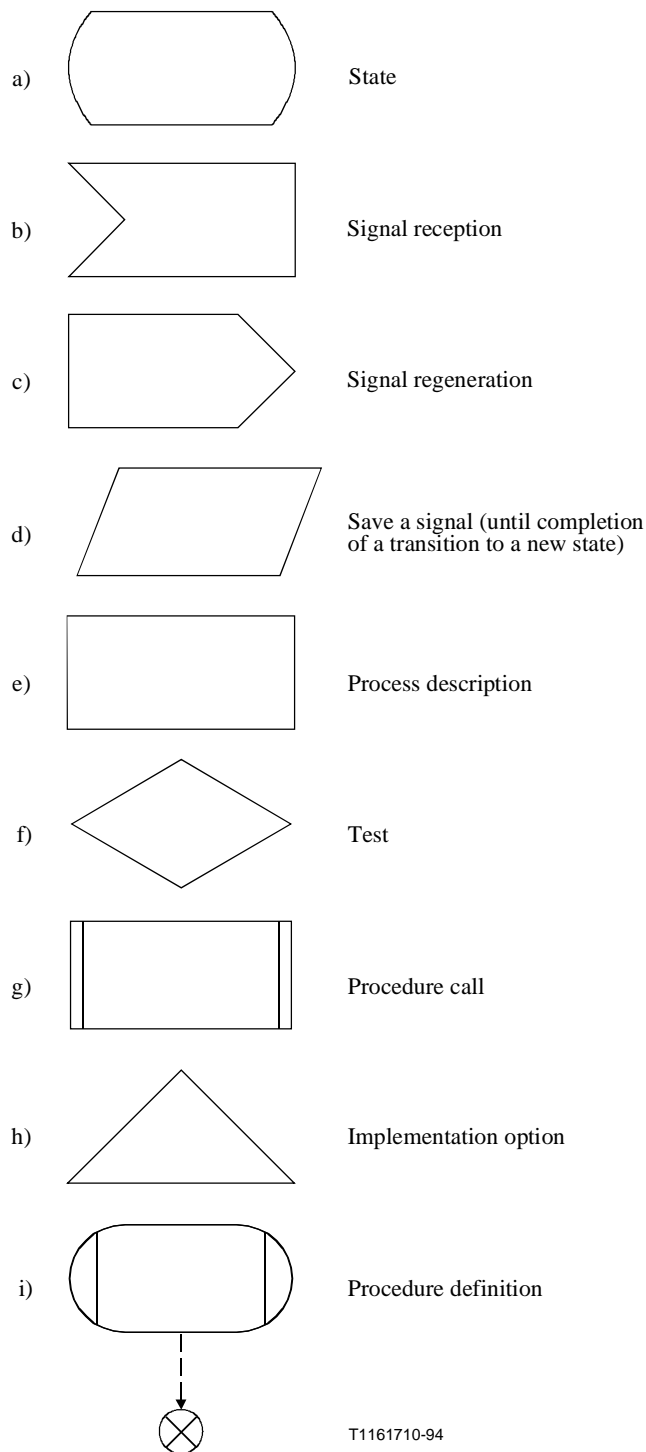


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Figure B.2/Q.921 – An overview of the states of the point-to-point procedures

B.3 Cover notes

The following symbols and abbreviations are used within this description. A full description of the symbols and their meaning and application can be found in the Series-Z Recommendations (Fascicles X.1 to X.5, *Blue Book*).



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- j) *** To mark an event or signal required as a result of the representation approach adopted, which is local to the data link layer entity
- k) RC Retransmission Counter
- l) (A-O) The codes used in the MDL-ERROR indication signals are defined in Table II.1. When multiple codes are shown, only one applies.

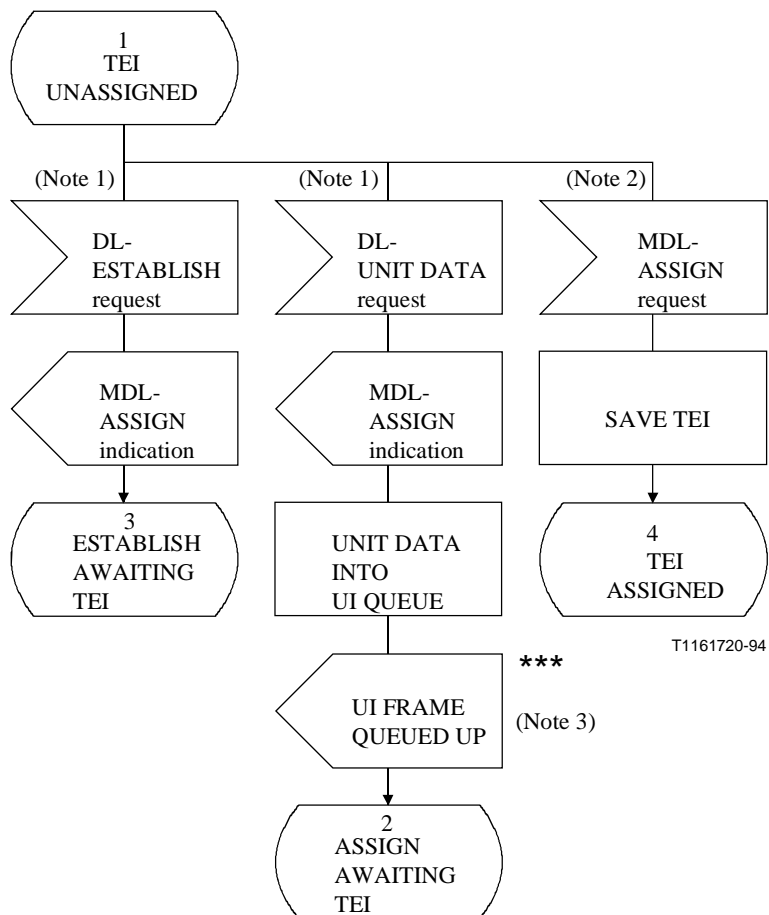
B.4 The use of queues

To enable a satisfactory representation of the data link layer entity, conceptual queues for the UI frame and I frame transmission have been explicitly brought out. These conceptual queues are finite but unbounded and should in no way restrict the implementation of the point-to-point procedures.

Two additional signals have been provided in order to cause the servicing of these queues to be initiated – UI frame queued up and I frame queued up.

B.5 SDL representation

See Figures B.3 to B.9.

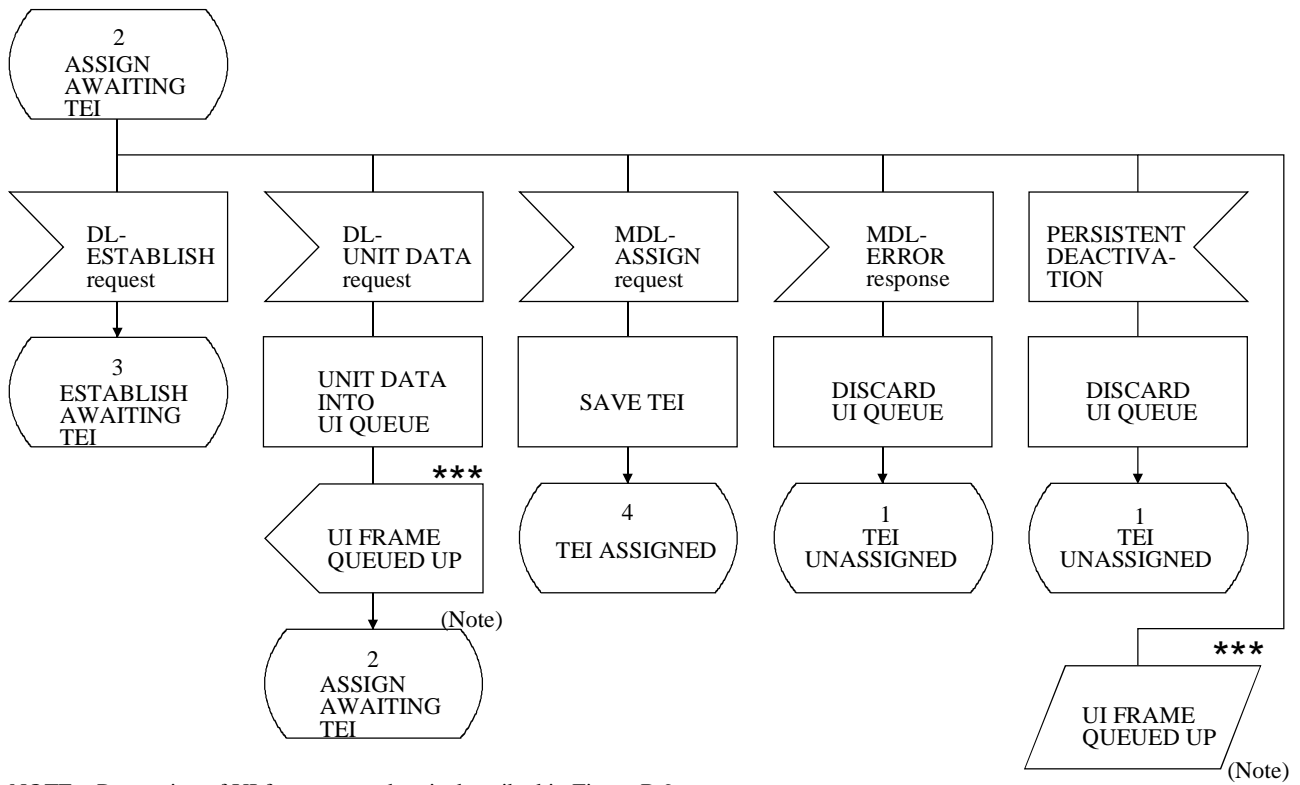


NOTE 1 – The use of these events on the network side is for further study.

NOTE 2 – This function may be implemented over a geographically distributed architecture. This primitive may occur on initialization for fixed TEIs at the network side, or as appropriate in order to correctly process a frame carrying a fixed TEI.

NOTE 3 – Processing of UI frame queued up is described in Figure B.9.

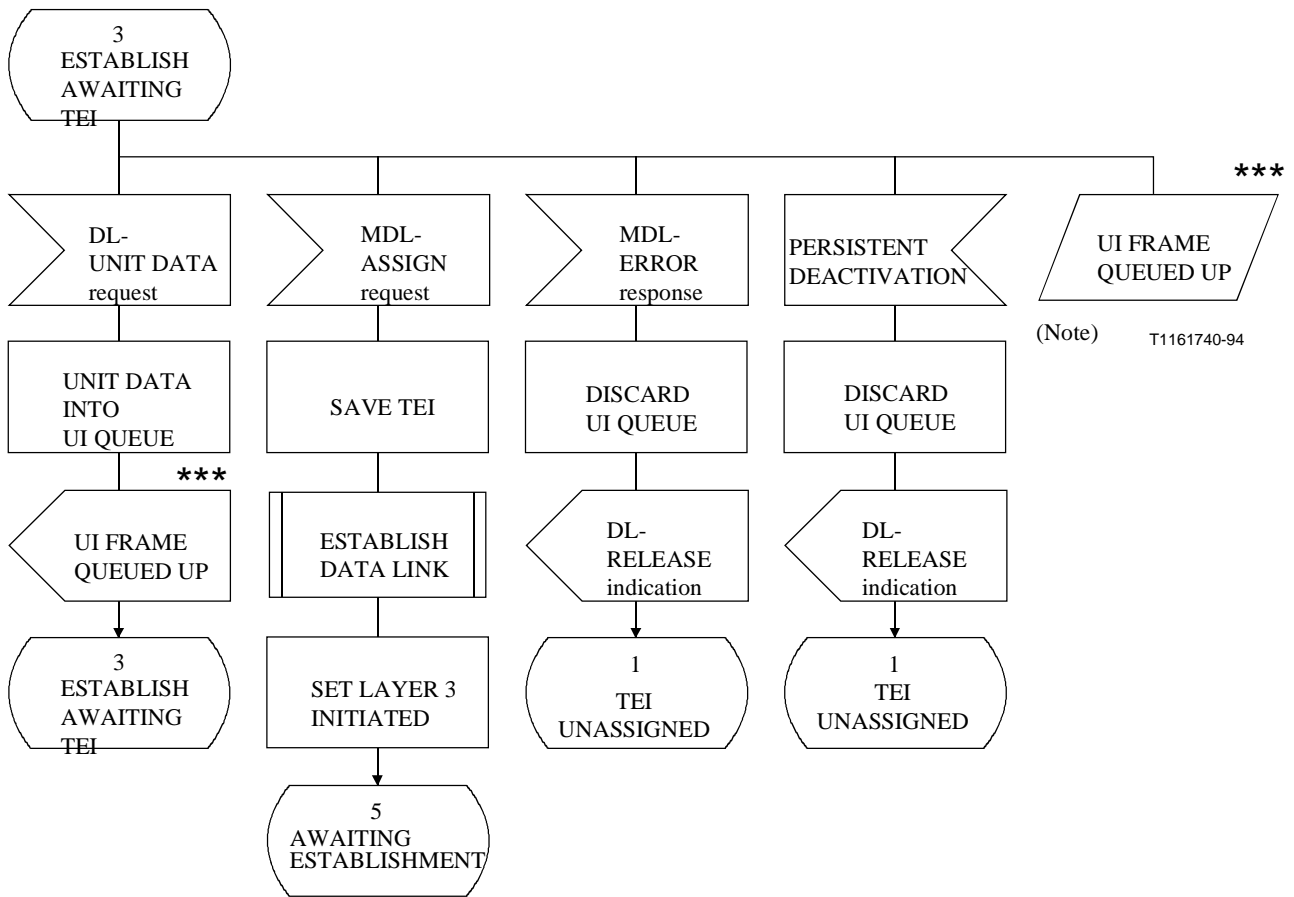
Figure B.3/Q.921 (sheet 1 of 3)



NOTE – Processing of UI frame queued up is described in Figure B.9.

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Figure B.3/Q.921 (sheet 2 of 3)



NOTE – Le traitement des trames UI mises en attente fait l'objet de la Figure B.9.

Figure B.3/Q.921 (sheet 3 of 3)

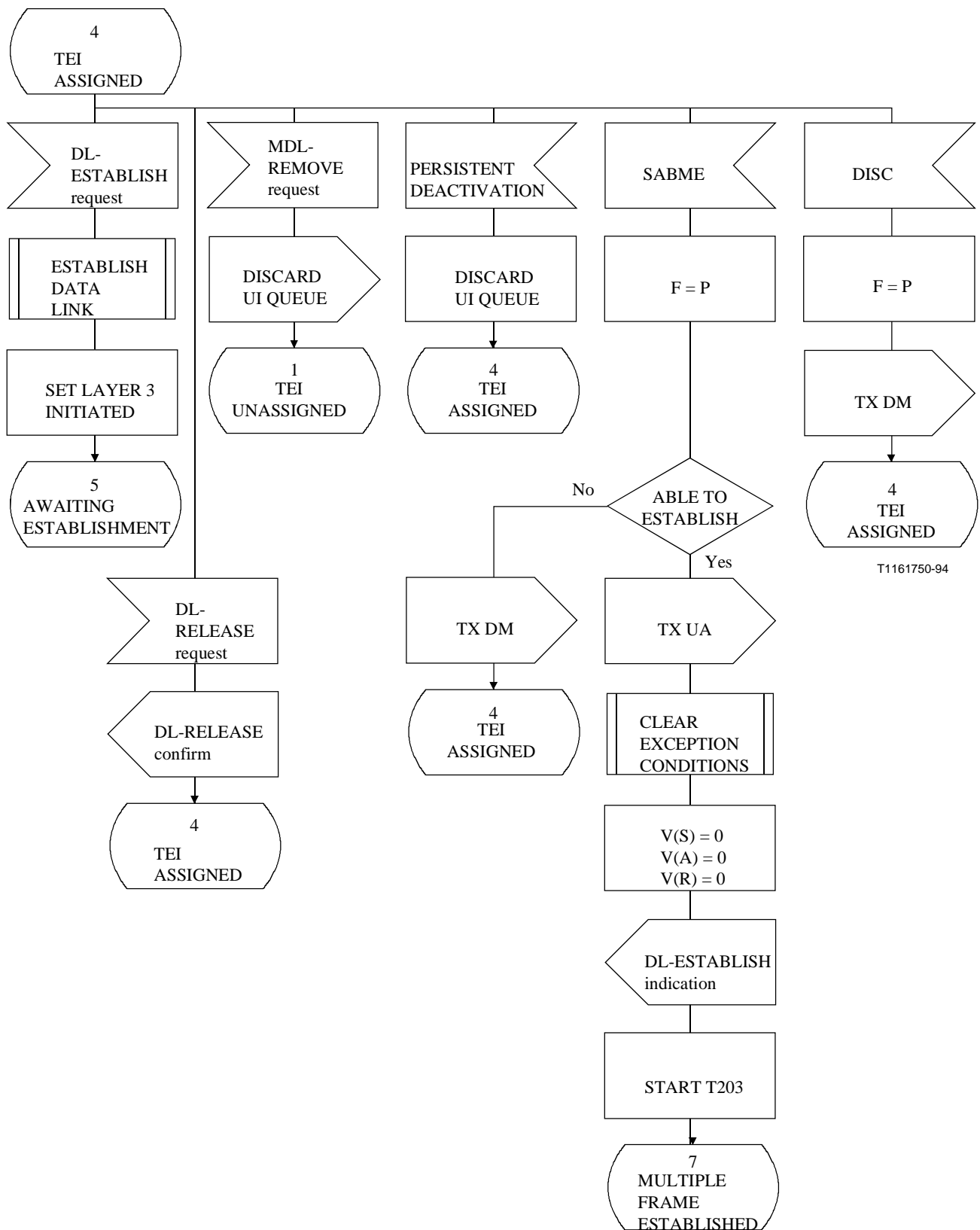
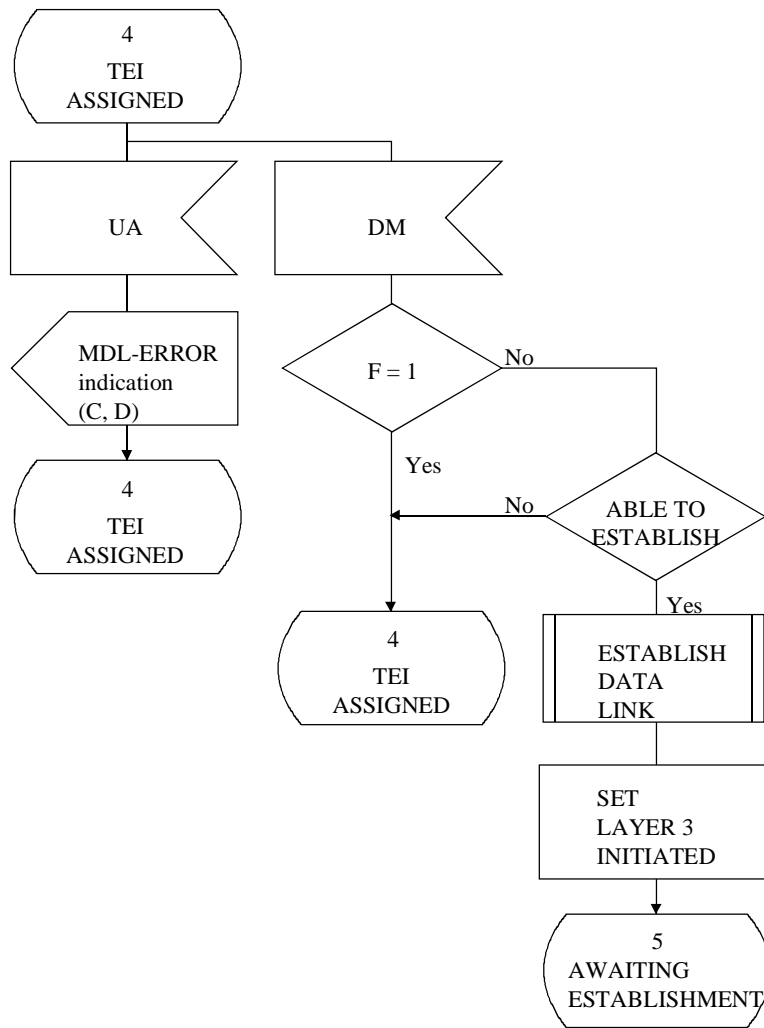
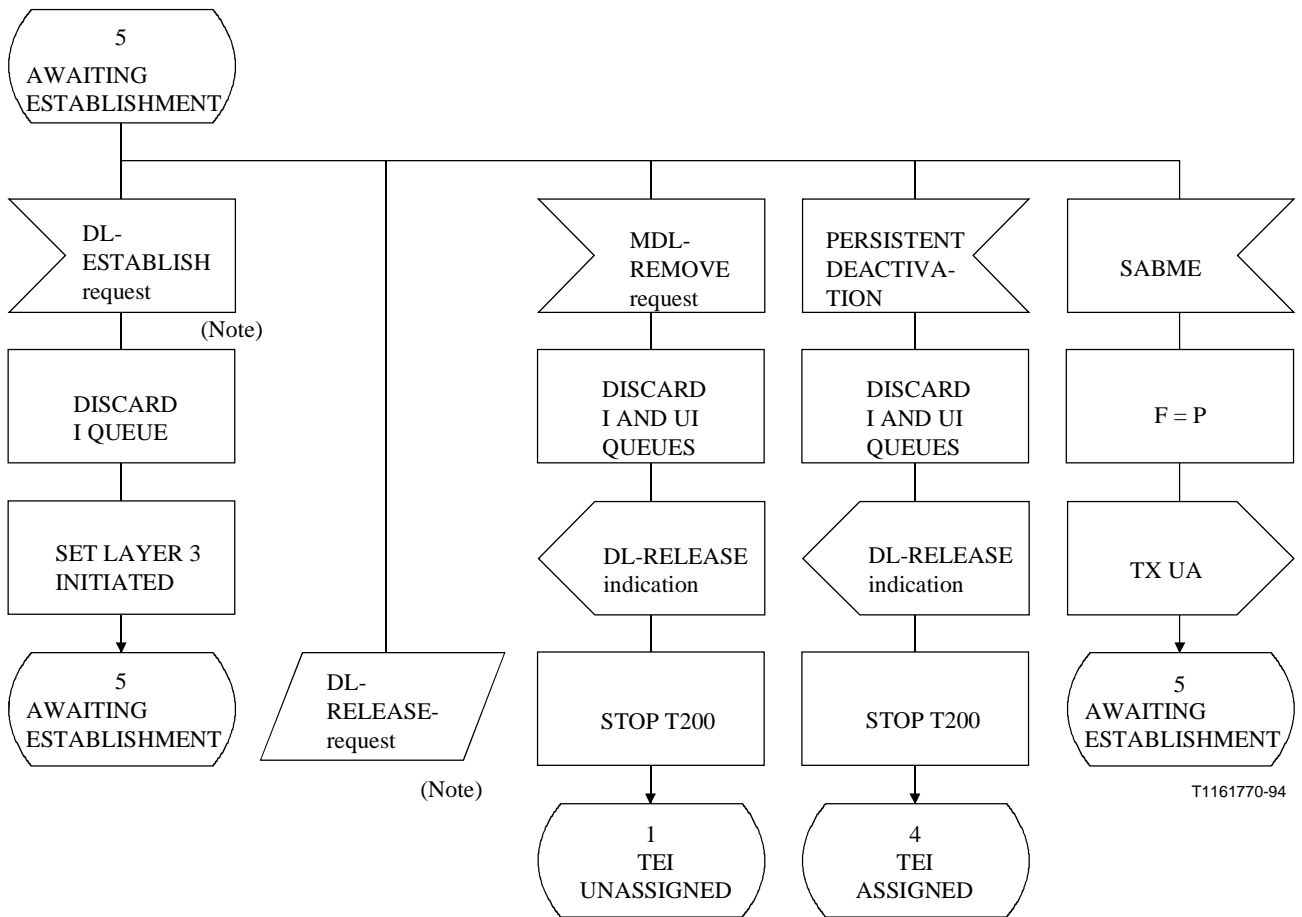


Figure B.4/Q.921 (sheet 1 of 2)



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Figure B.4/Q.921 (sheet 2 of 2)



T1161770-94

NOTE – Only possible in cases of Layer 2 initiated re-establishment.

Figure B.5/Q.921 (sheet 1 of 3)

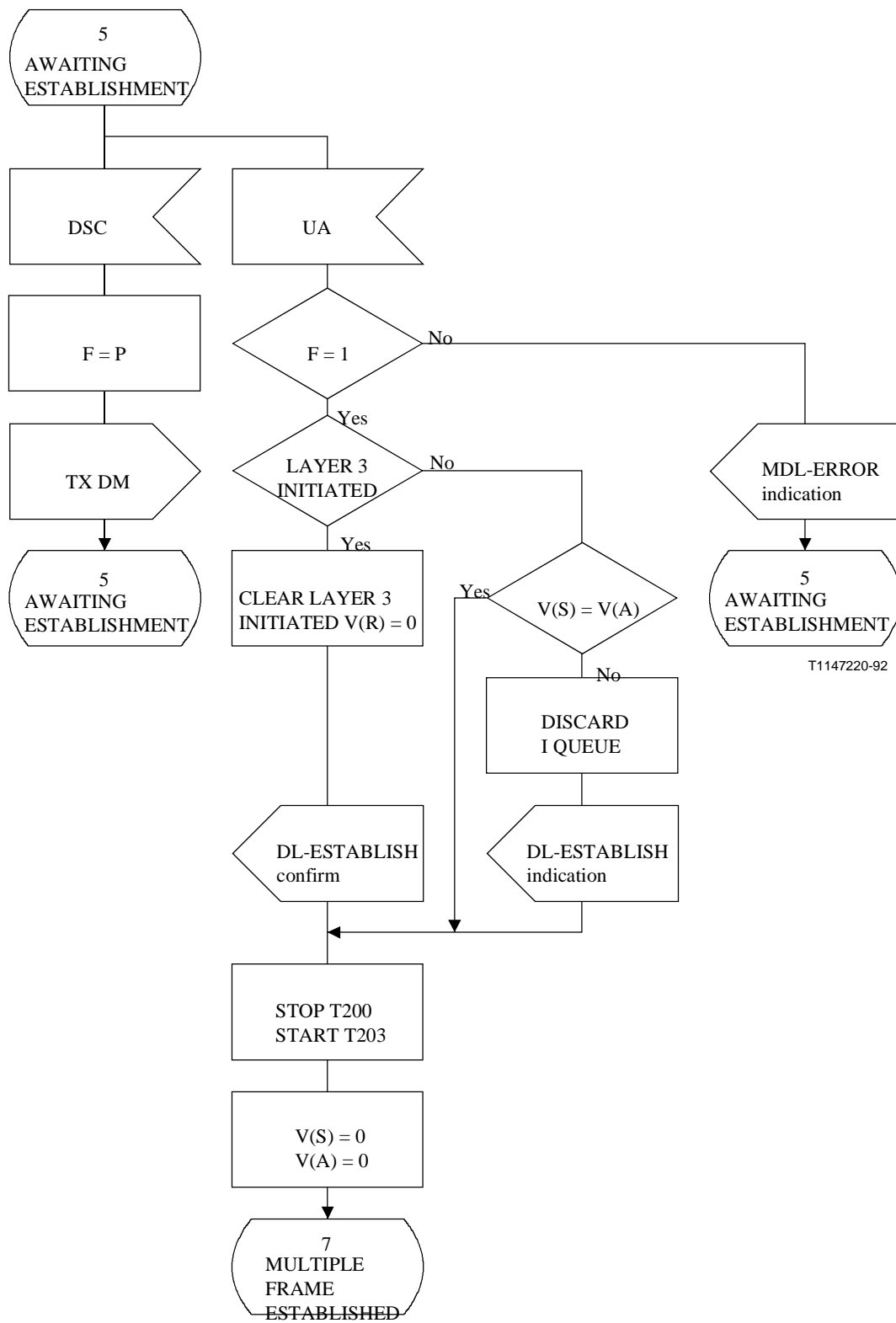
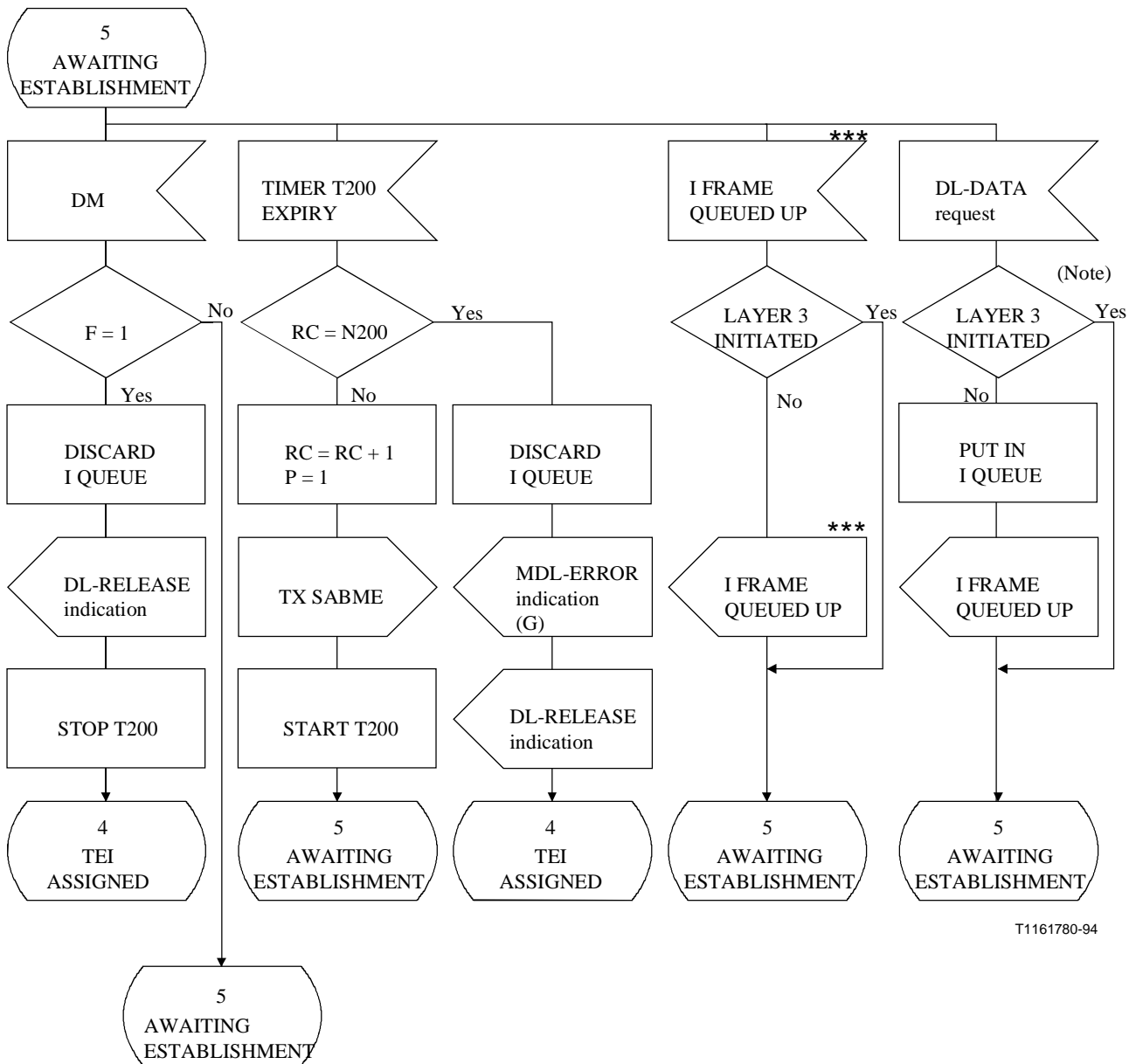


Figure B.5/Q.921 (sheet 2 of 3)



T1161780-94

NOTE – Only possible in cases of Layer 2 initiated re-establishment.

Figure B.5/Q.921 (sheet 3 of 3)

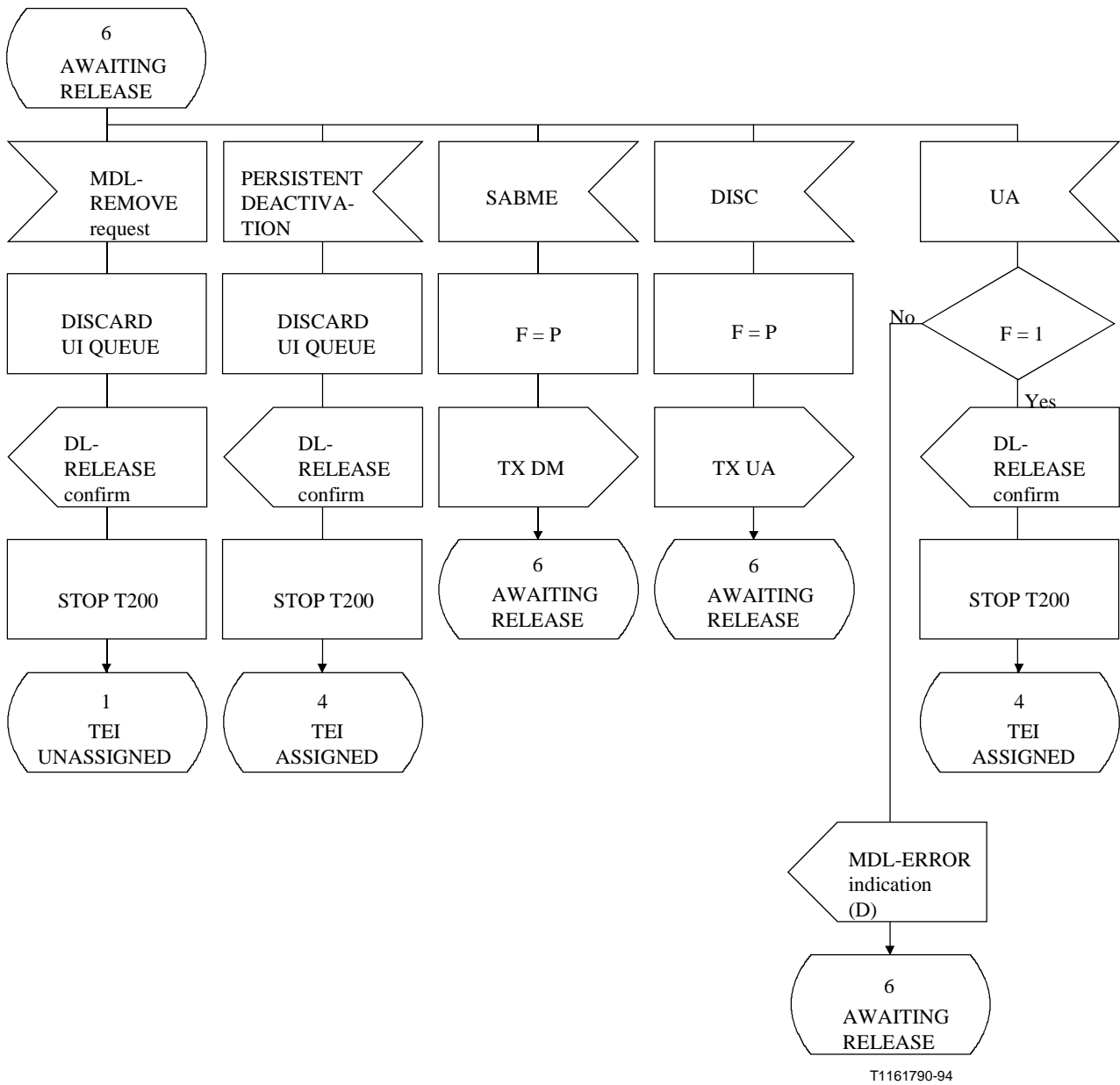
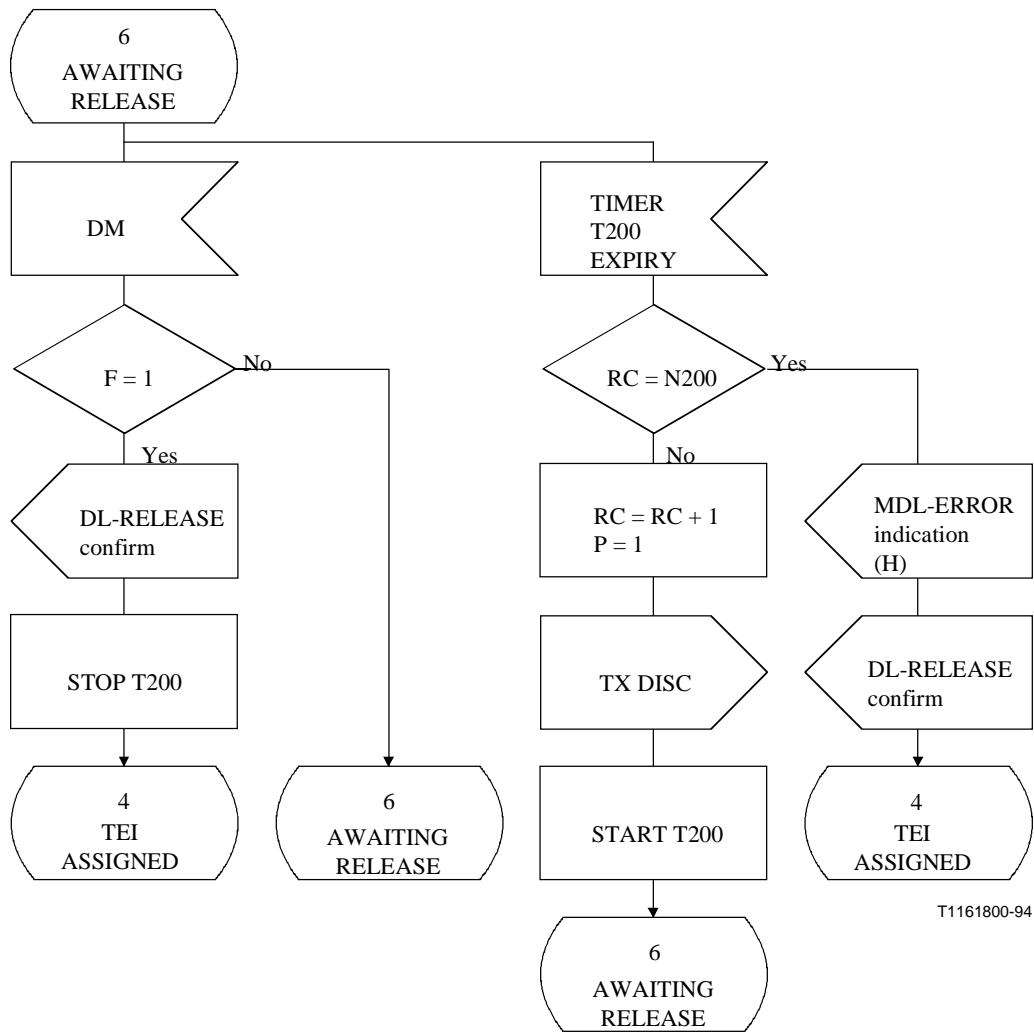
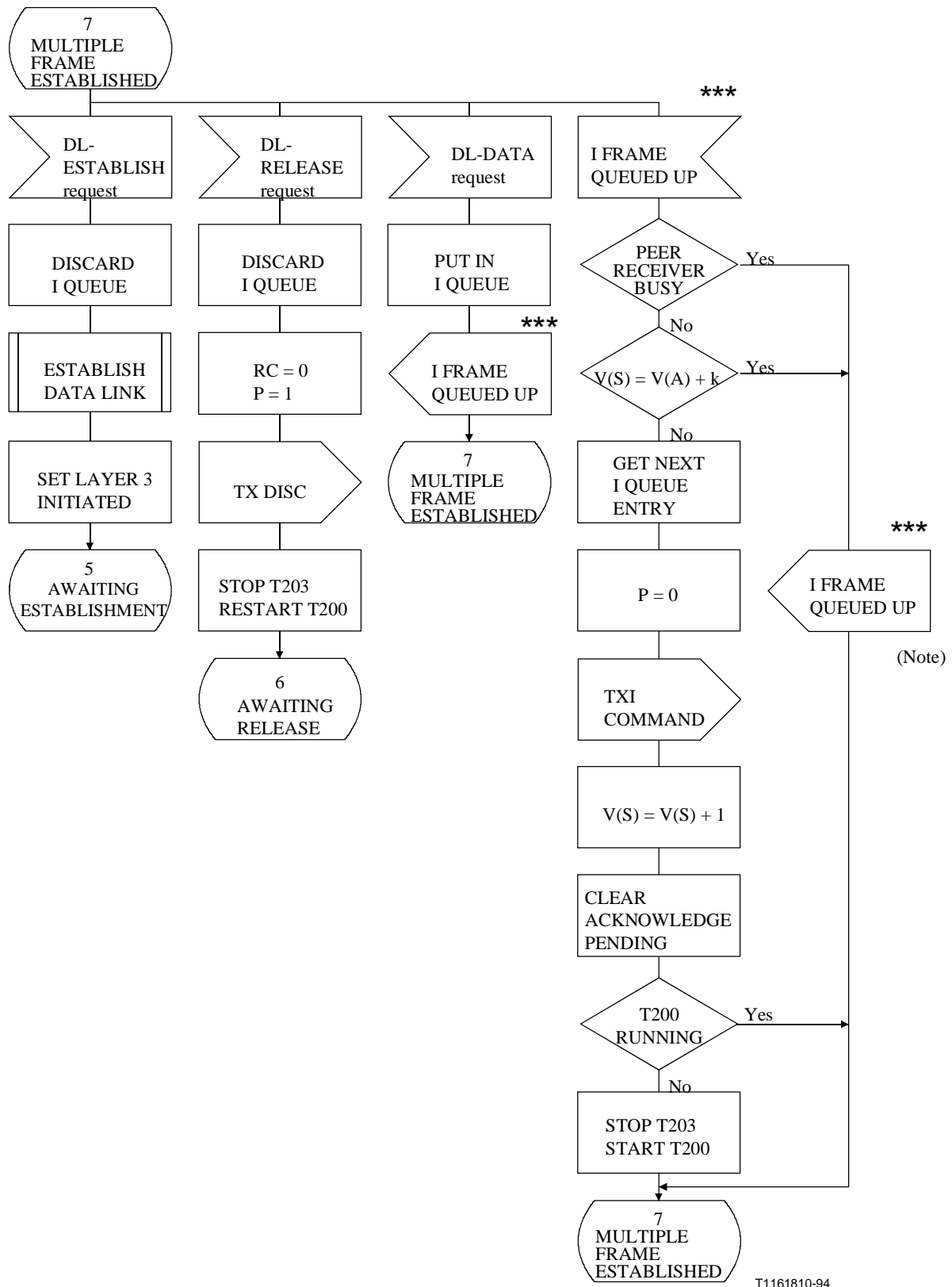


Figure B.6/Q.921 (sheet 1 of 2)



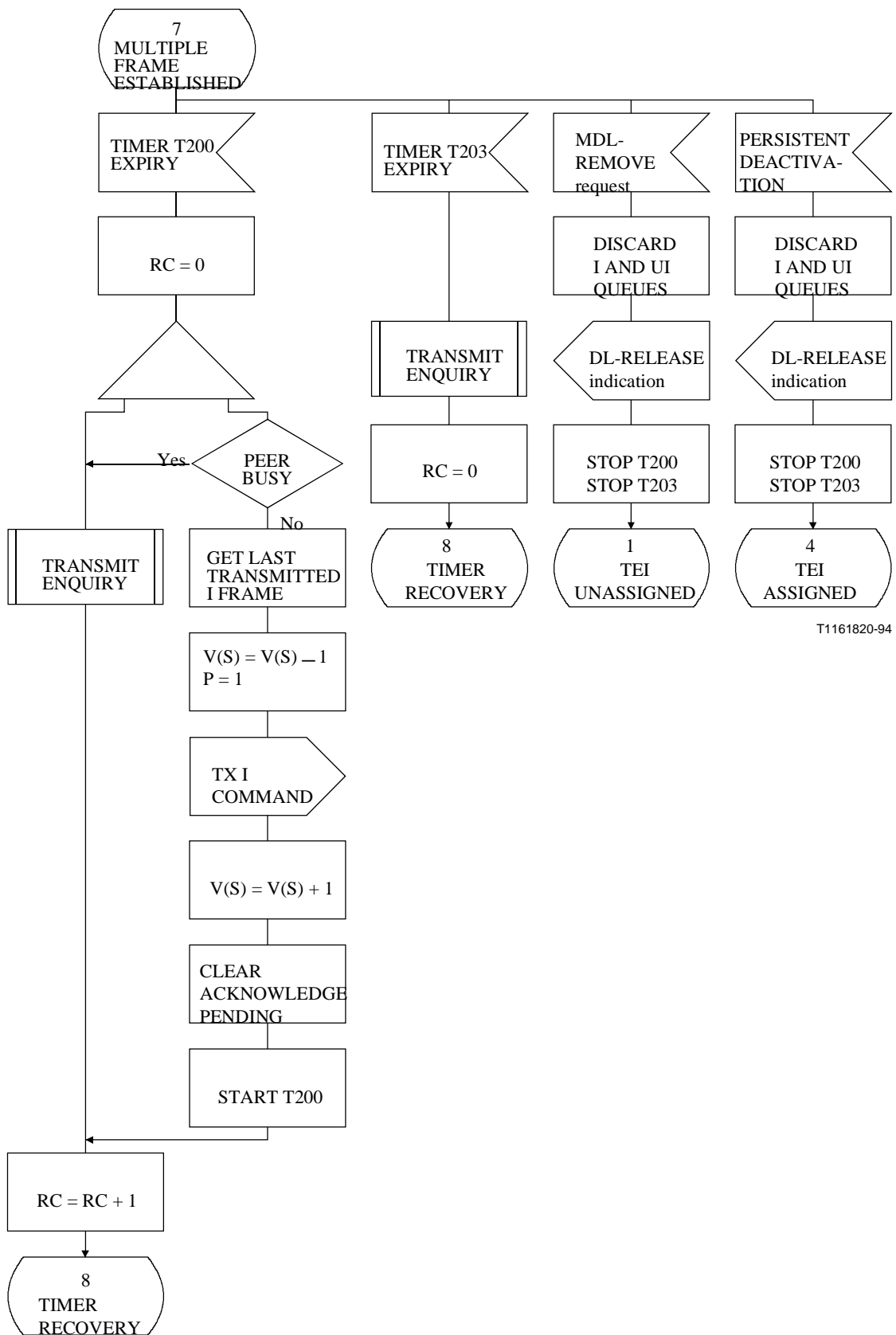
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Figure B.6/Q.921 (sheet 2 of 2)



NOTE – The regeneration of this signal does not affect the sequence integrity of the I queue.

Figure B.7/Q.921 (sheet 1 of 10)



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Figure B.7/Q.921 (sheet 2 of 10)

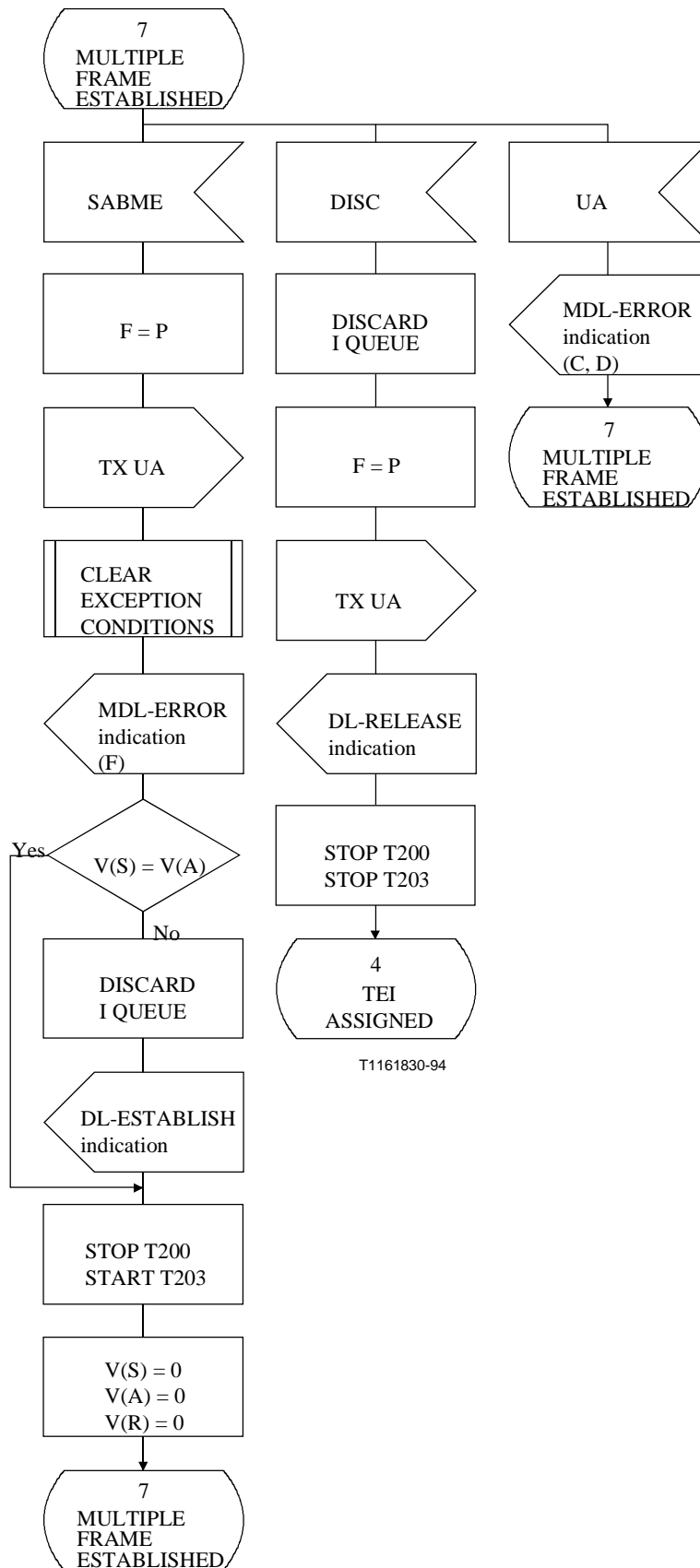
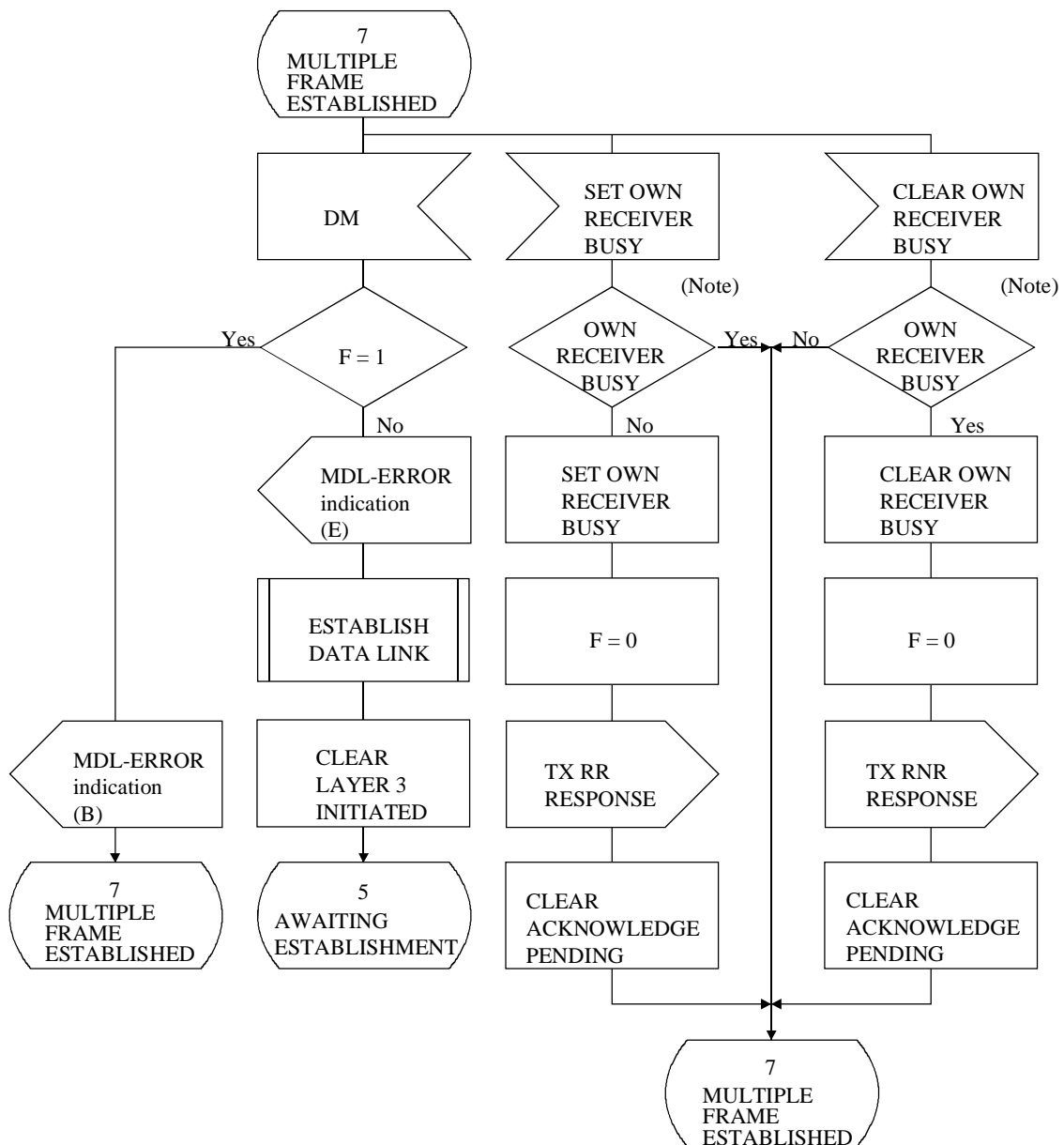


Figure B.7/Q.921 (sheet 3 of 10)



T1161840-94

NOTE – These signals are generated outside of this SDL representation, and may be generated by the connection management entity.

Figure B.7/Q.921 (sheet 4 of 10)

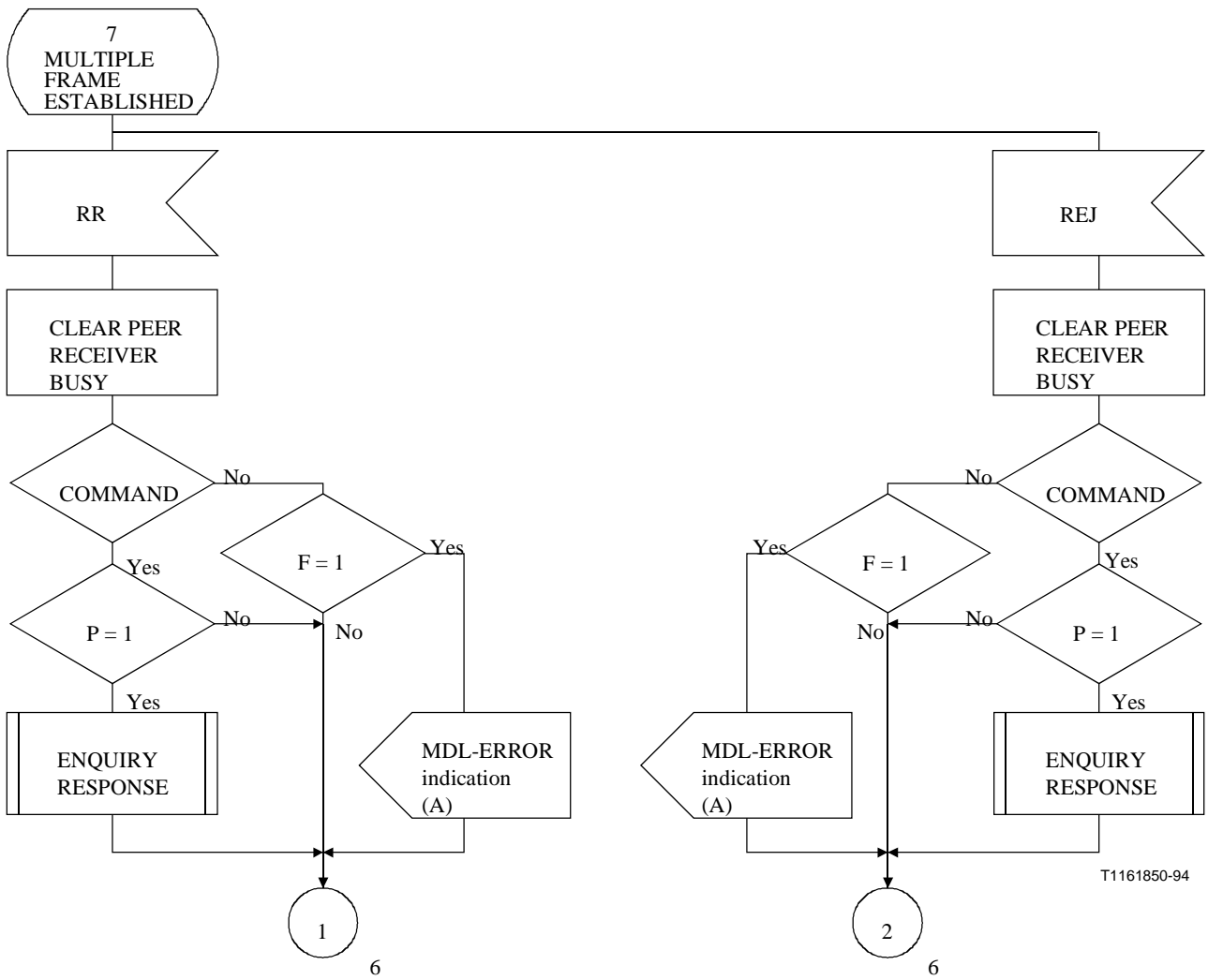


Figure B.7/Q.921 (sheet 5 of 10)

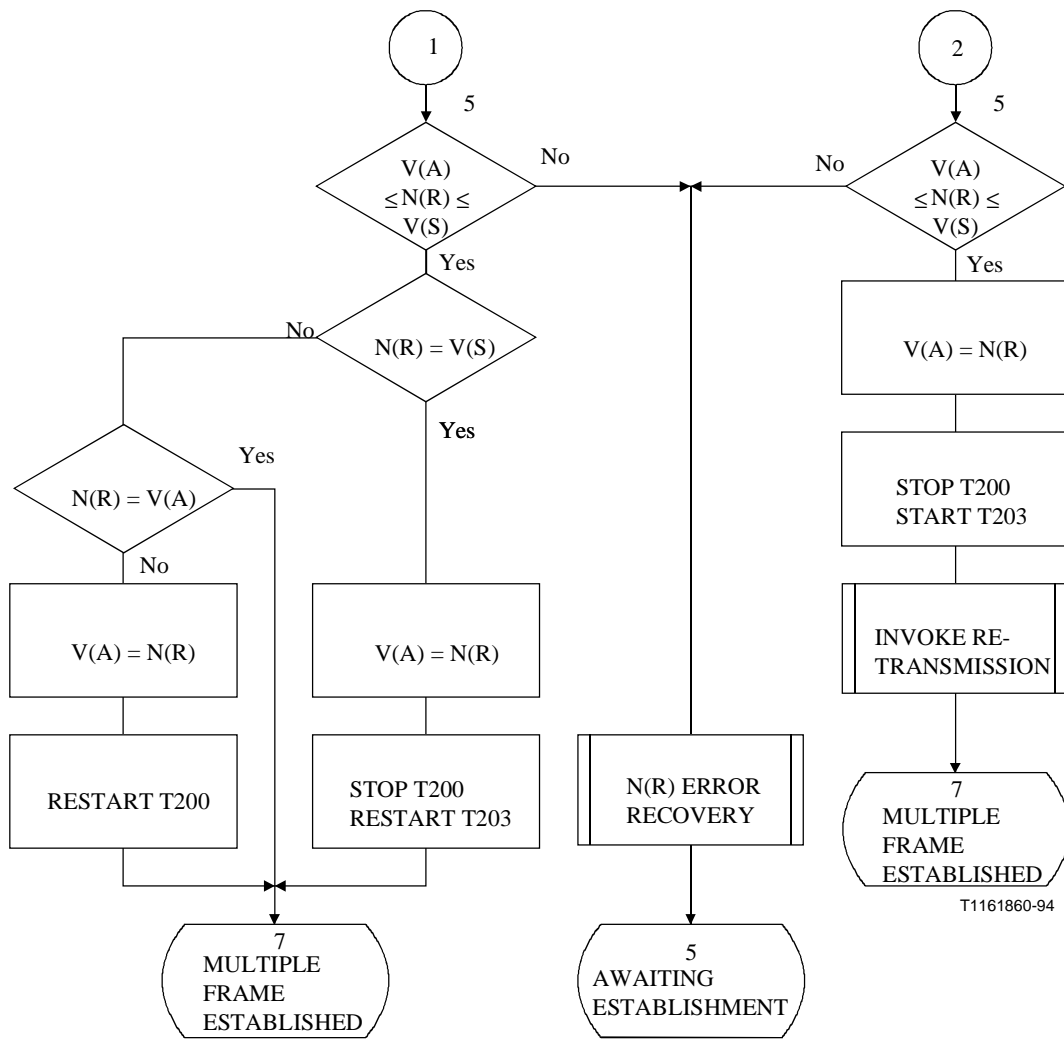
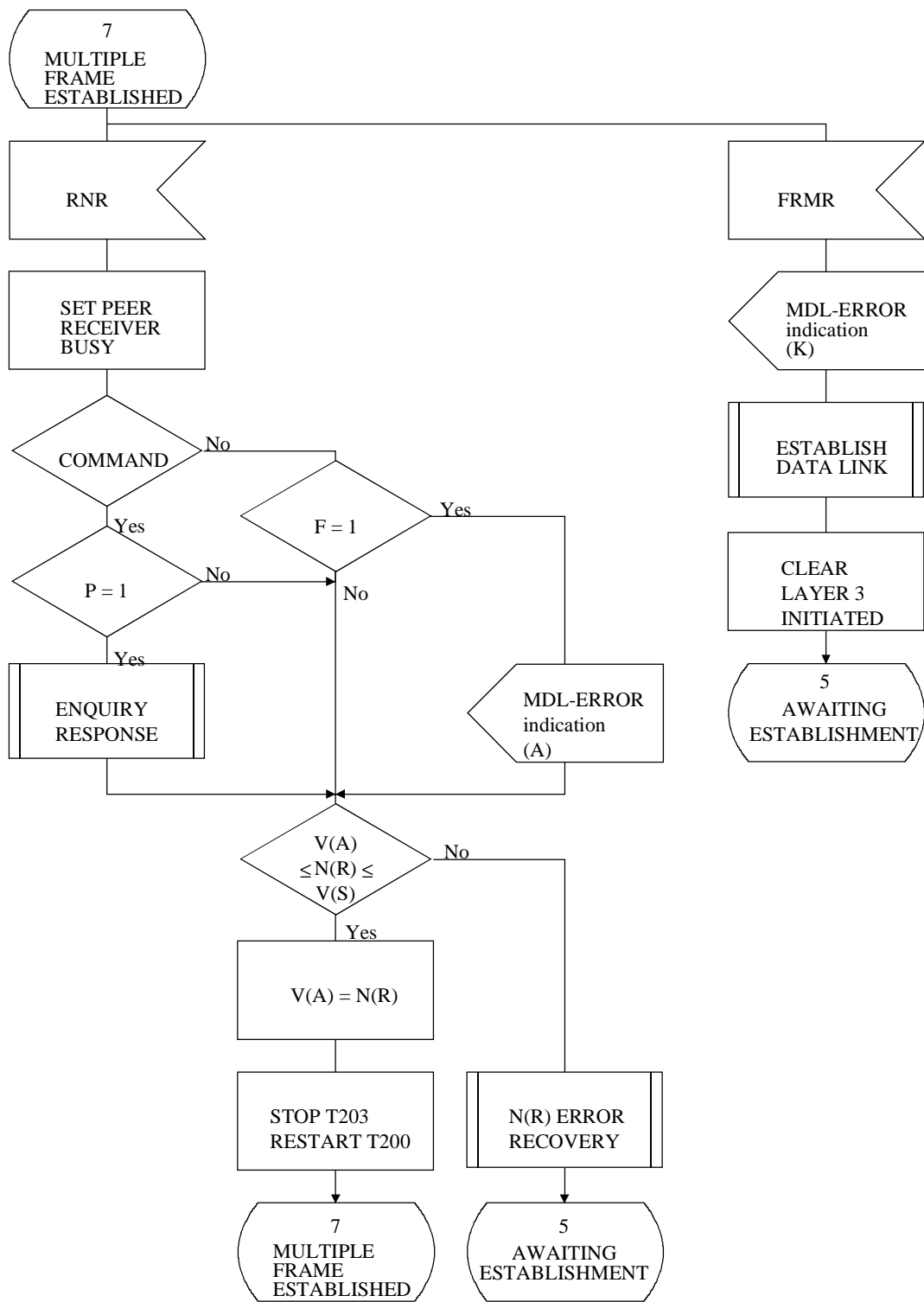
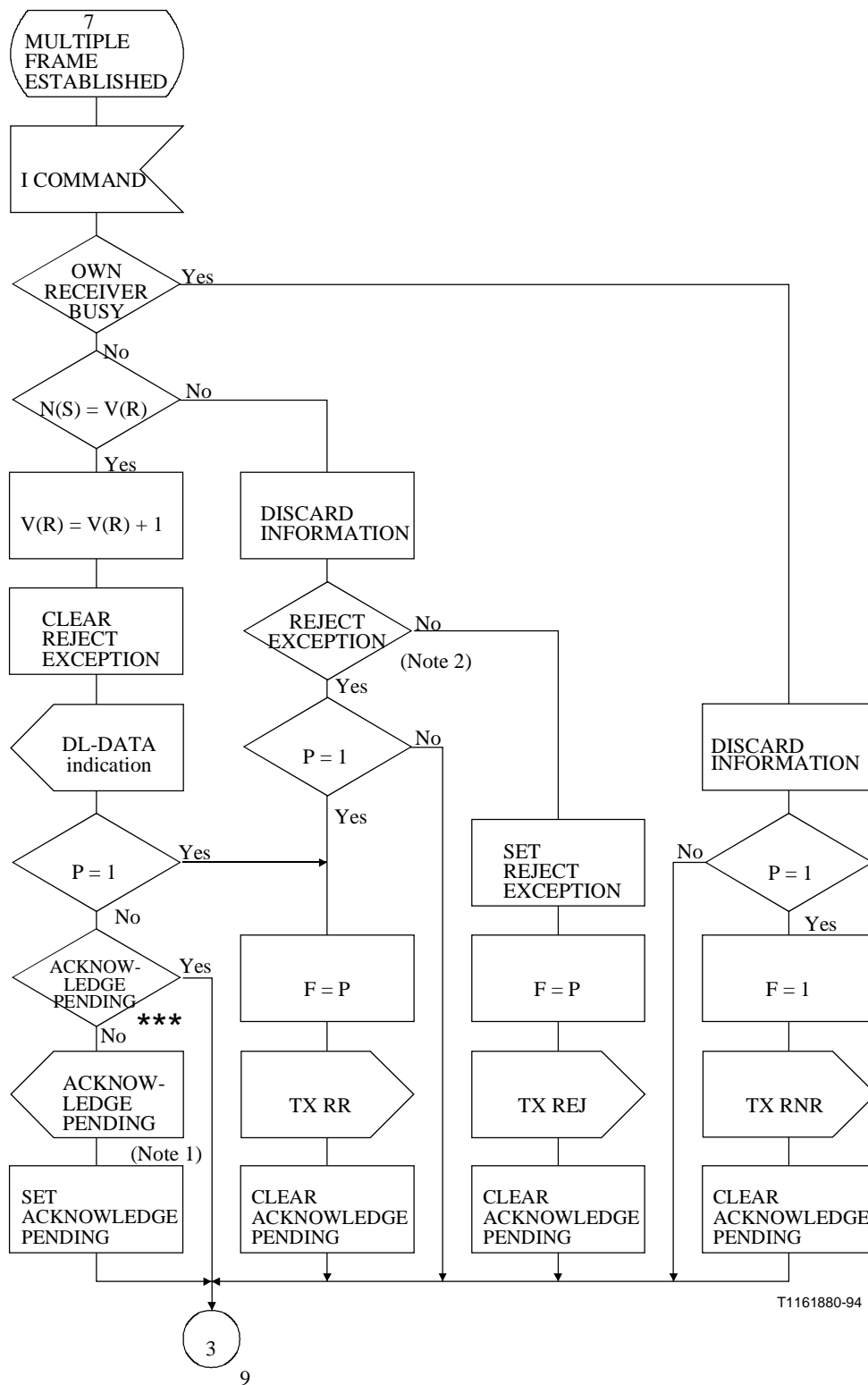


Figure B.7/Q.921 (sheet 6 of 10)



T1161870-94

Figure B.7/Q.921 (sheet 7 of 10)



T1161880-94

NOTE 1 – Processing of acknowledge pending is described on sheet 10 of this Figure.

NOTE 2 – This SDL representation does not include the optional procedure in Appendix I.

Figure B.7/Q.921 (sheet 8 of 10)

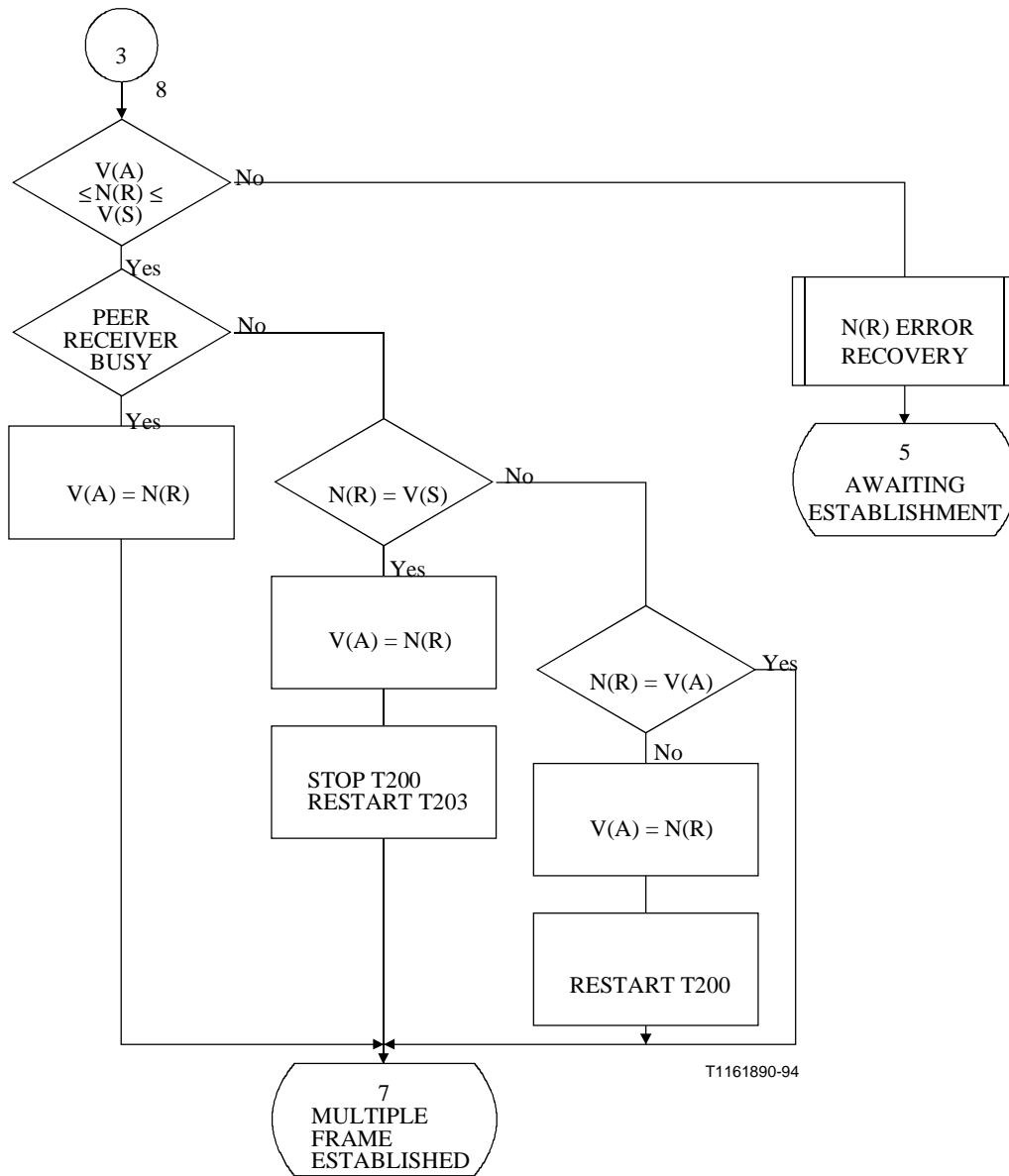
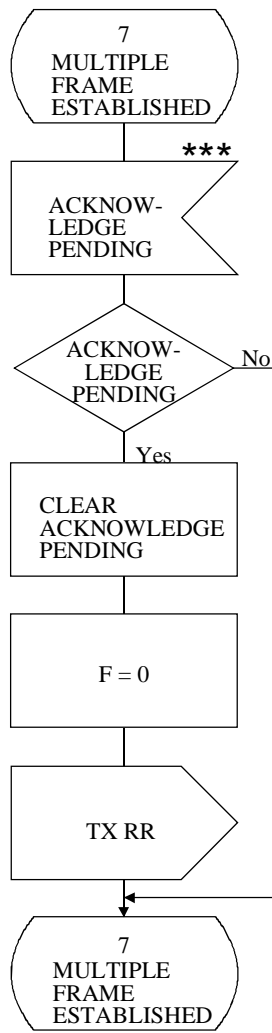
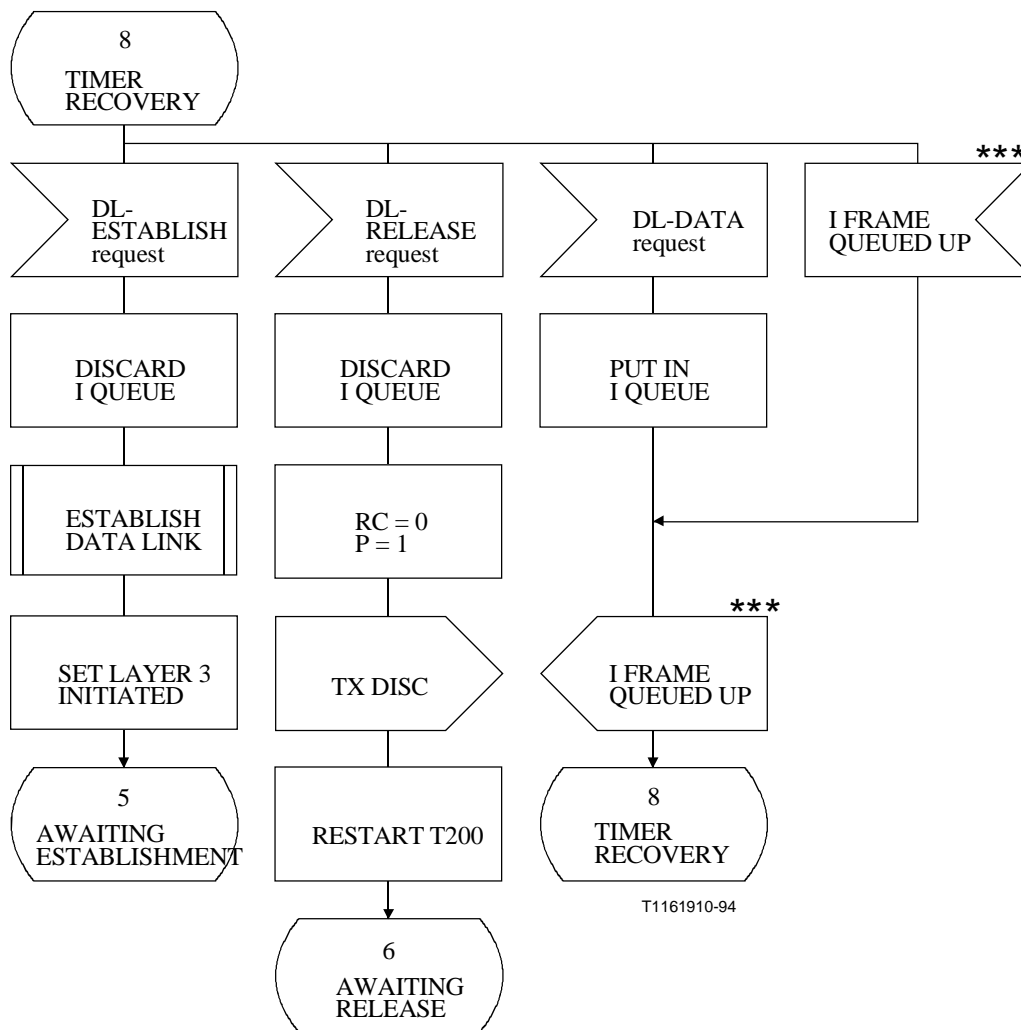


Figure B.7/Q.921 (sheet 9 of 10)



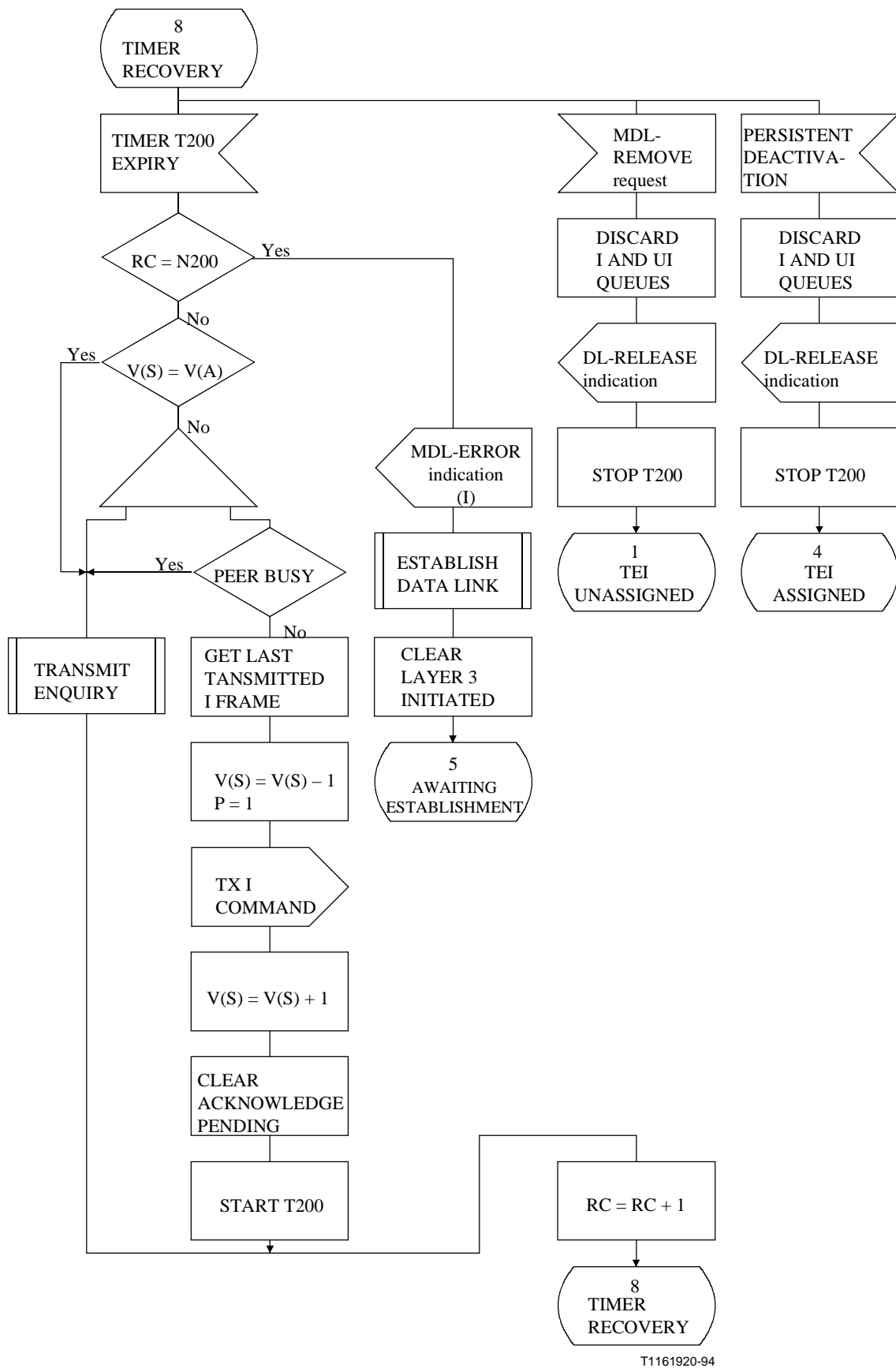
T1161900-94

Figure B.7/Q.921 (sheet 10 of 10)



T1161910-94

Figure B.8/Q.921 (sheet 1 of 9)



T1161920-94

Figure B.8/Q.921 (sheet 2 of 9)

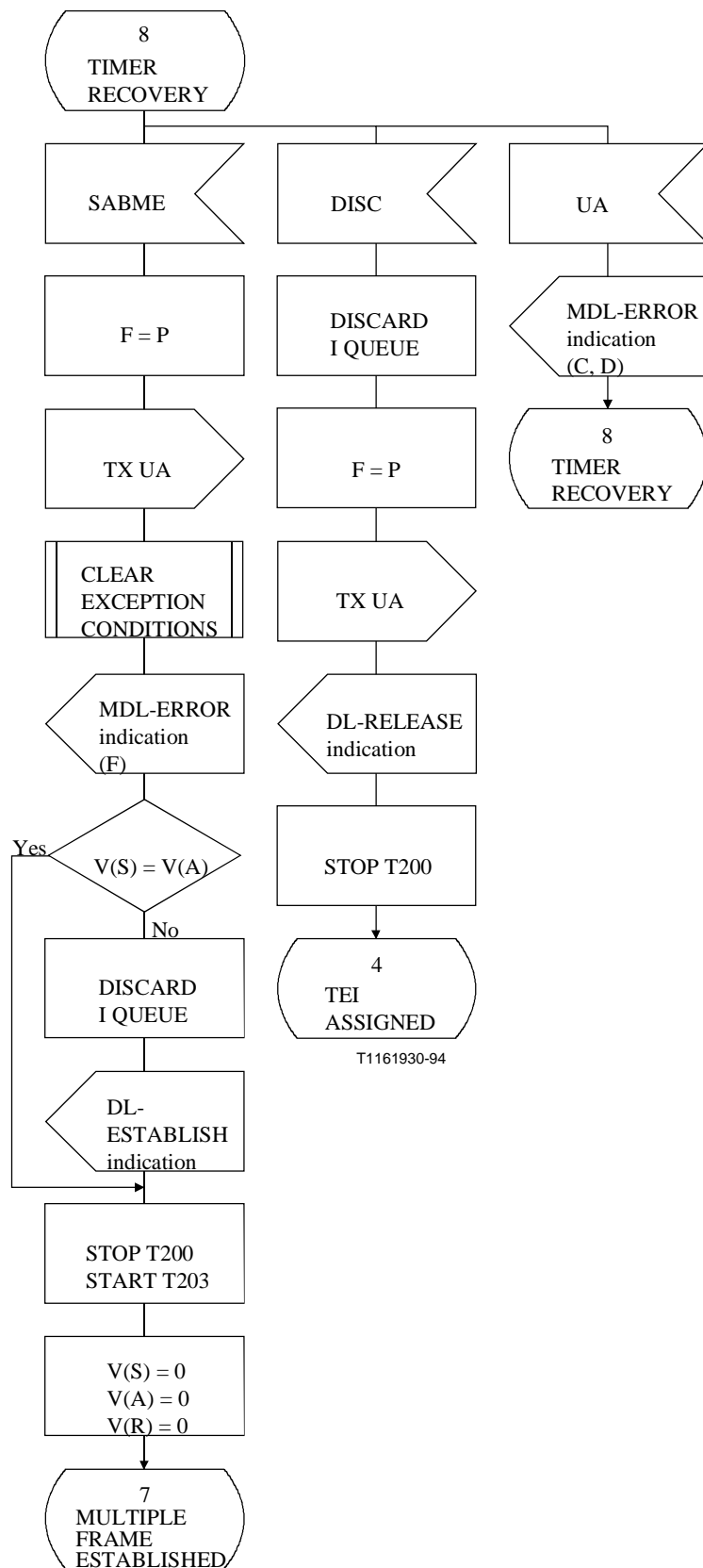
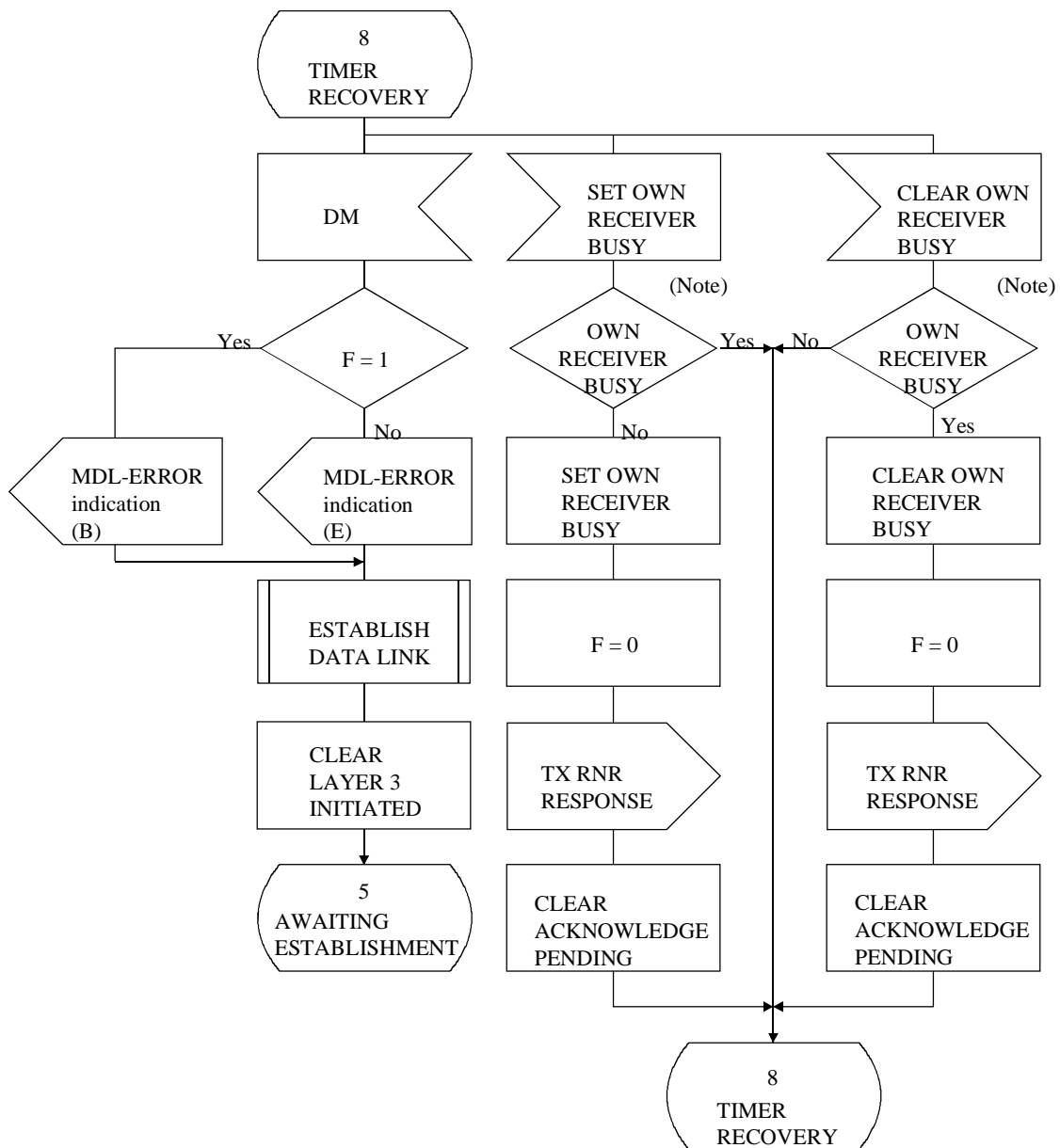


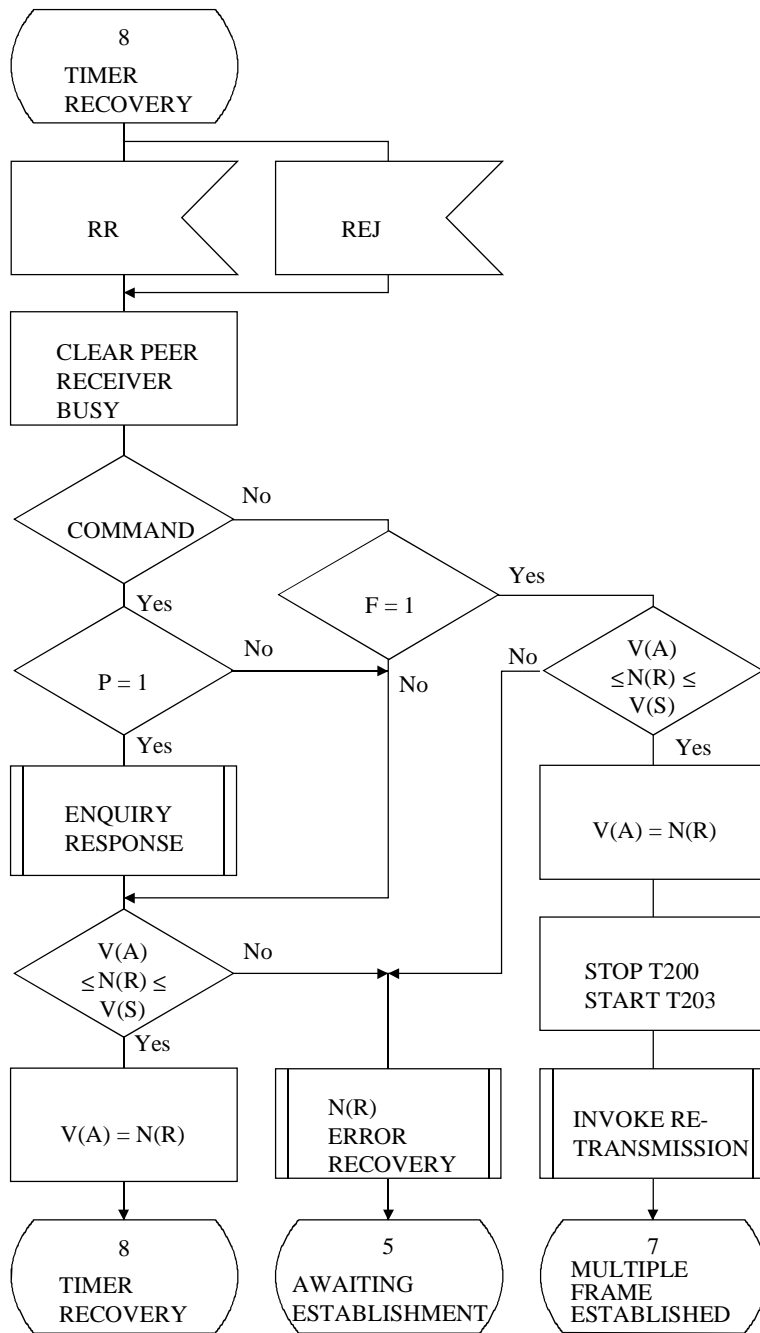
Figure B.8/Q.921 (sheet 3 of 9)



T1161940-94

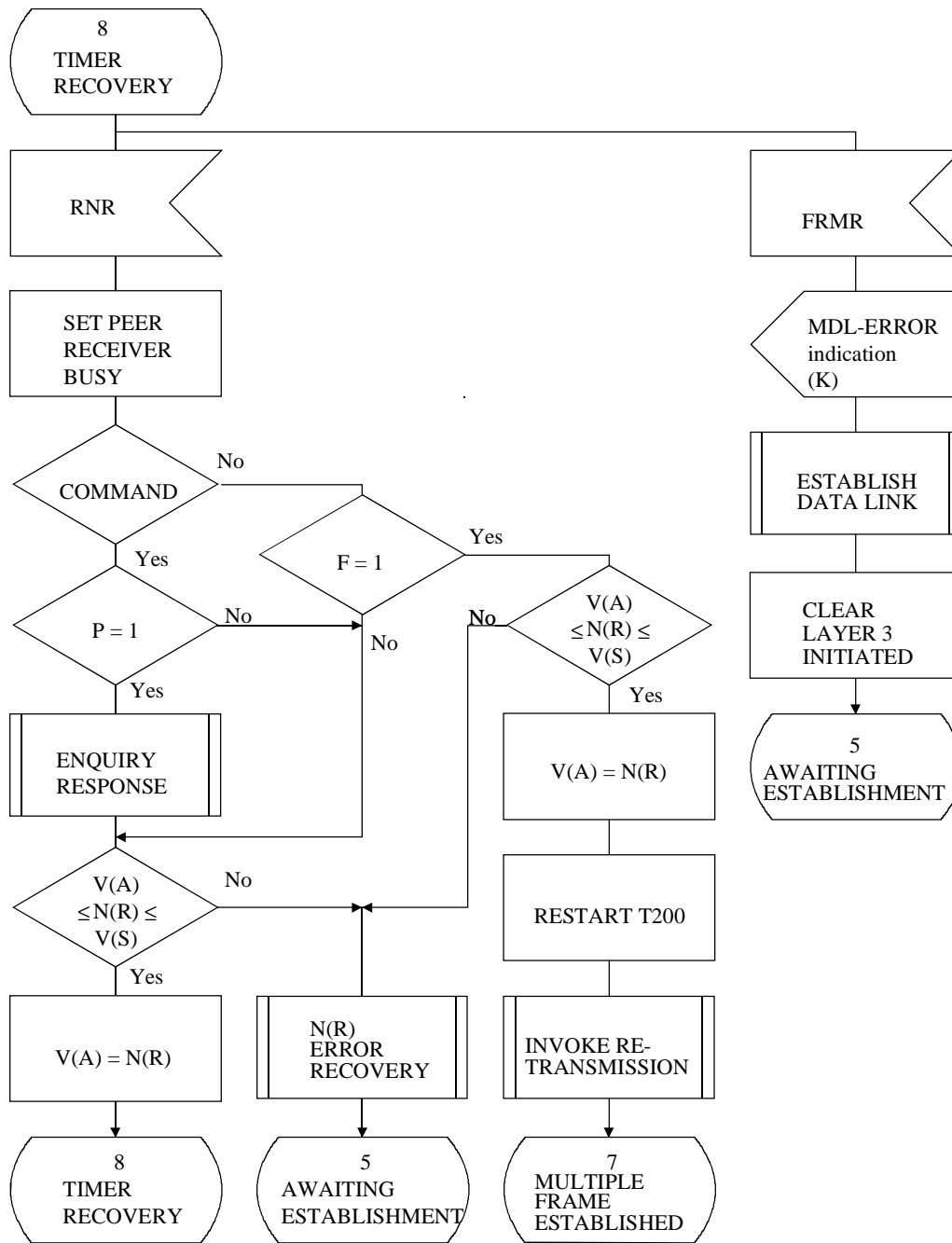
NOTE – These signals are generated outside of this SDL representation, and may be generated by the connection management entity.

Figure B.8/Q.921 (sheet 4 of 9)



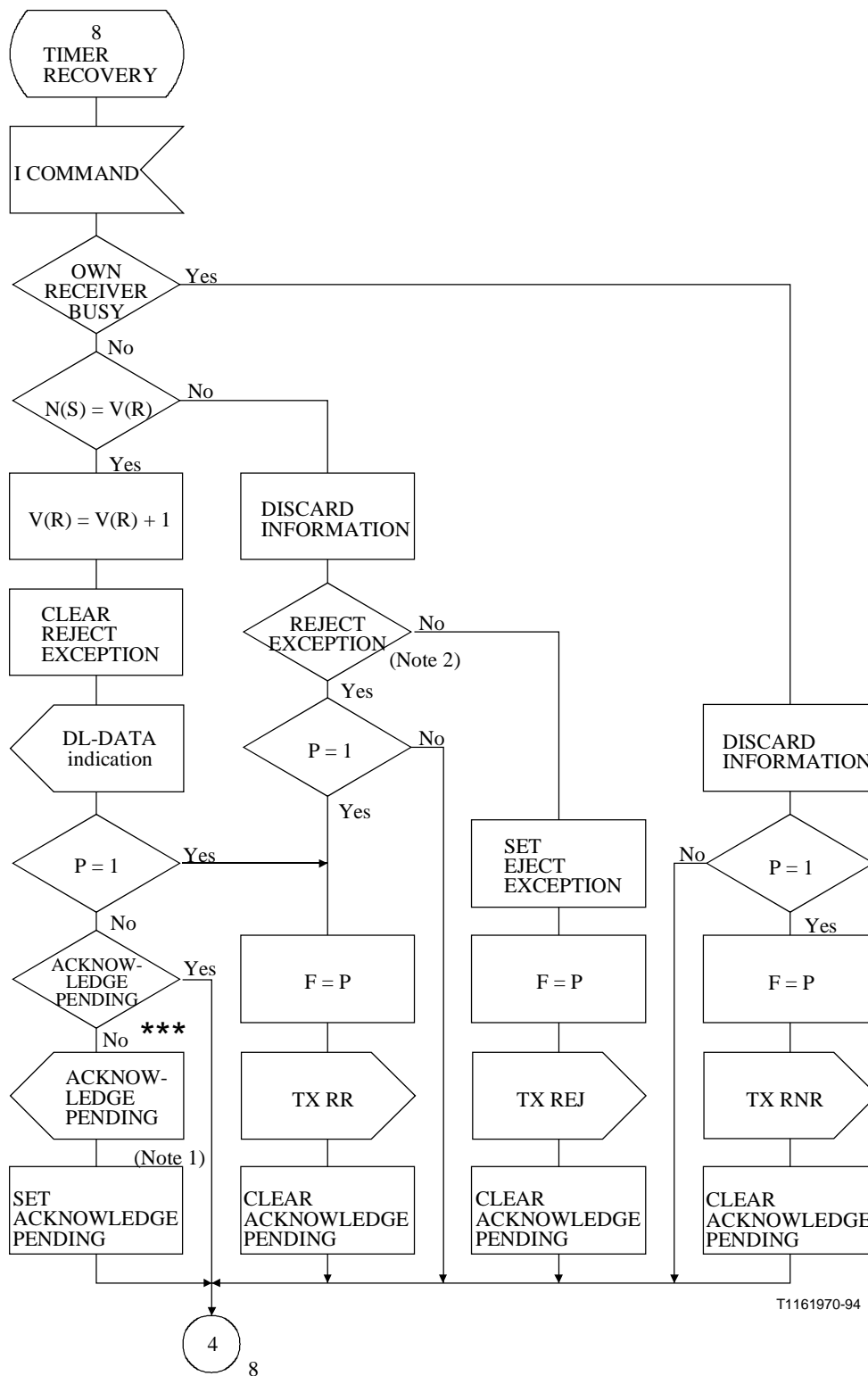
T1161950-94

Figure B.8/Q.921 (sheet 5 of 9)



T1161960-94

Figure B.8/Q.921 (sheet 6 of 9)



NOTE 1 – Processing of Acknowledge Pending is described on sheet 9 of this Figure.

NOTE 2 – This SDL representation does not include the optional procedure in Appendix I.

Figure B.8/Q.921 (sheet 7 of 9)

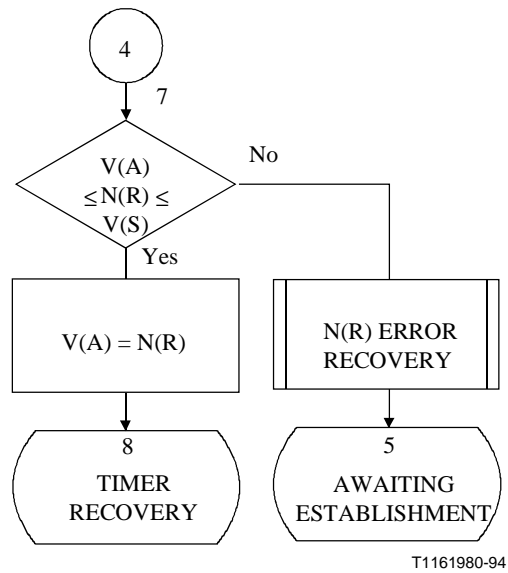


Figure B.8/Q.921 (sheet 8 of 9)

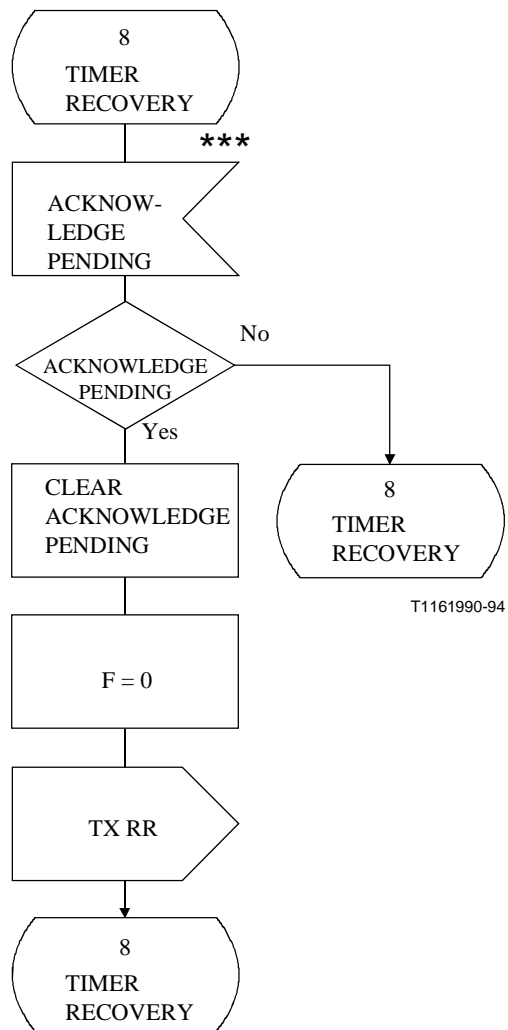
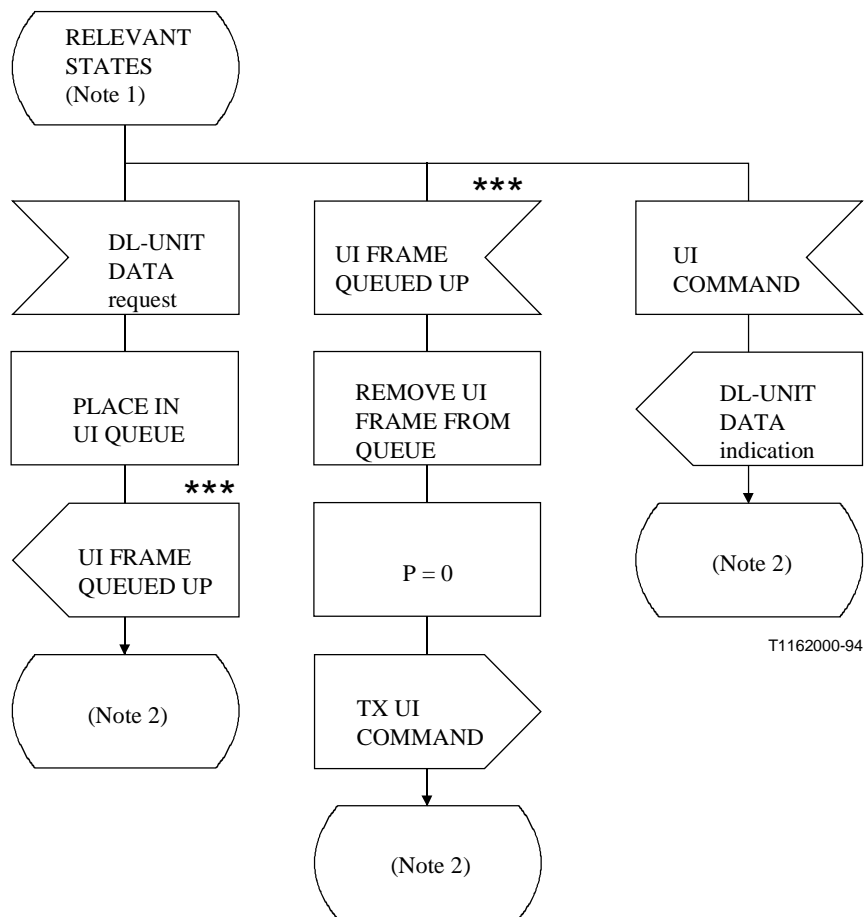


Figure B.8/Q.921 (sheet 9 of 9)

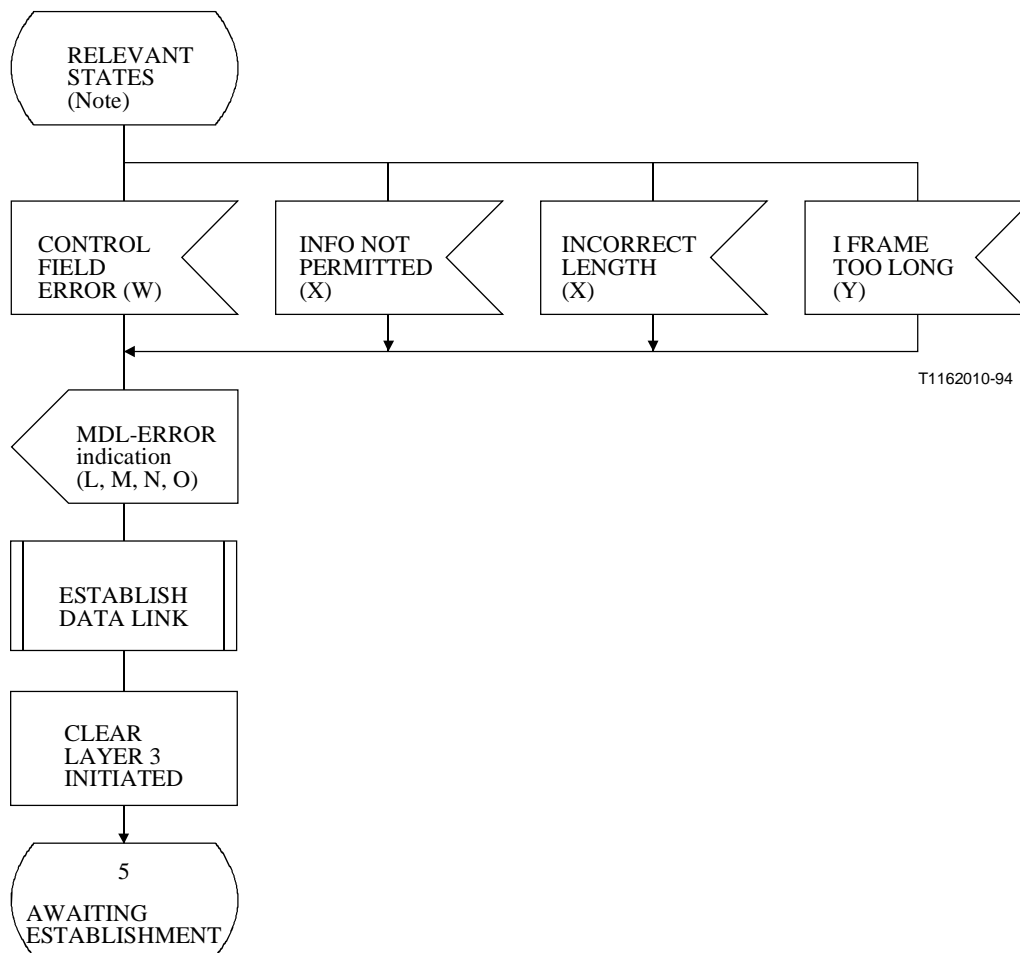


NOTE 1 – The relevant states are as follows:

- 4 *TEI-assigned*
- 5 *Awaiting-establishment*
- 6 *Awaiting-release*
- 7 *Multiple-frame-established*
- 8 *Timer-recovery.*

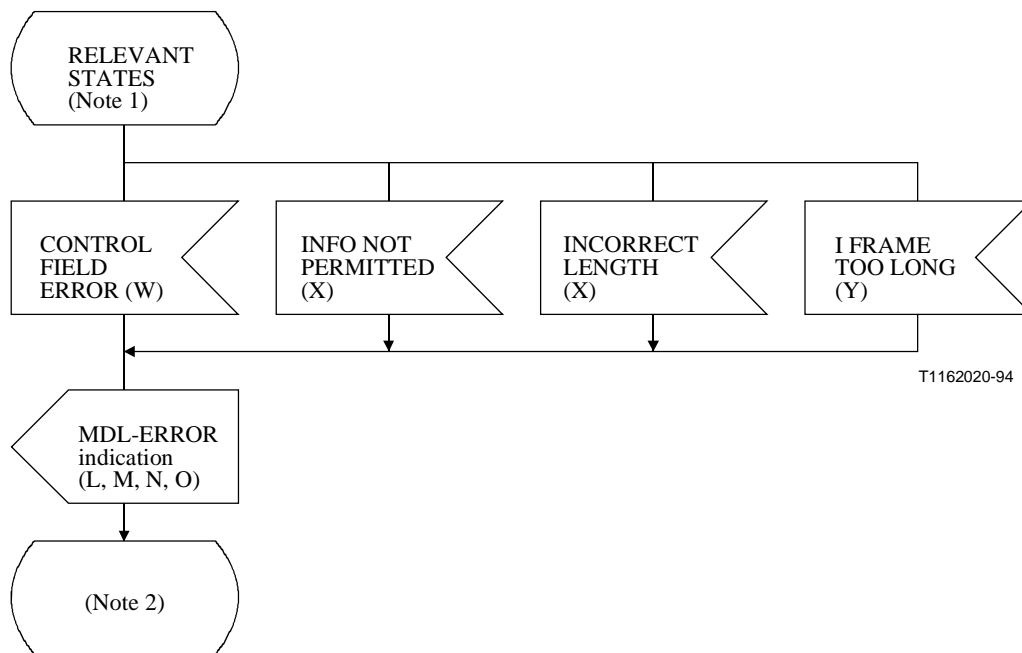
NOTE 2 – The data link layer returns to the state it was in prior to the events shown.

Figure B.9/Q.921 (sheet 1 of 5)



NOTE – The relevant states are as follows:
 1 *Multiple-frame-established*
 8 *Timer-recovery.*

Figure B.9/Q.921 (sheet 2 of 5)



NOTE 1 – The relevant states are as follows:

- 4 *TEI-assigned*
- 5 *Awaiting-establishment*
- 6 *Awaiting-release.*

NOTE 2 – The data link layer returns to the state it was in prior to the events shown.

Figure B.9/Q.921 (sheet 3 of 5)

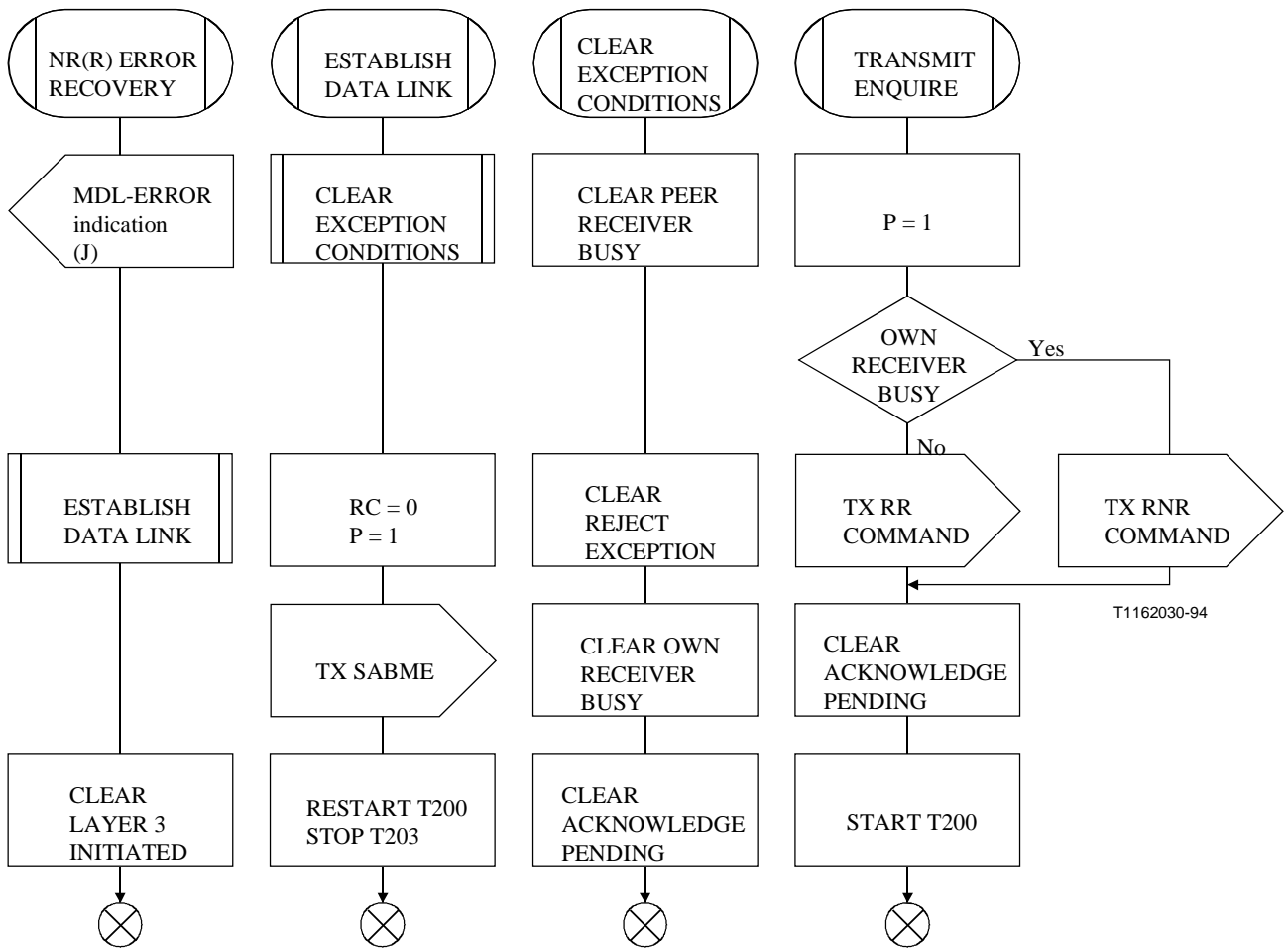
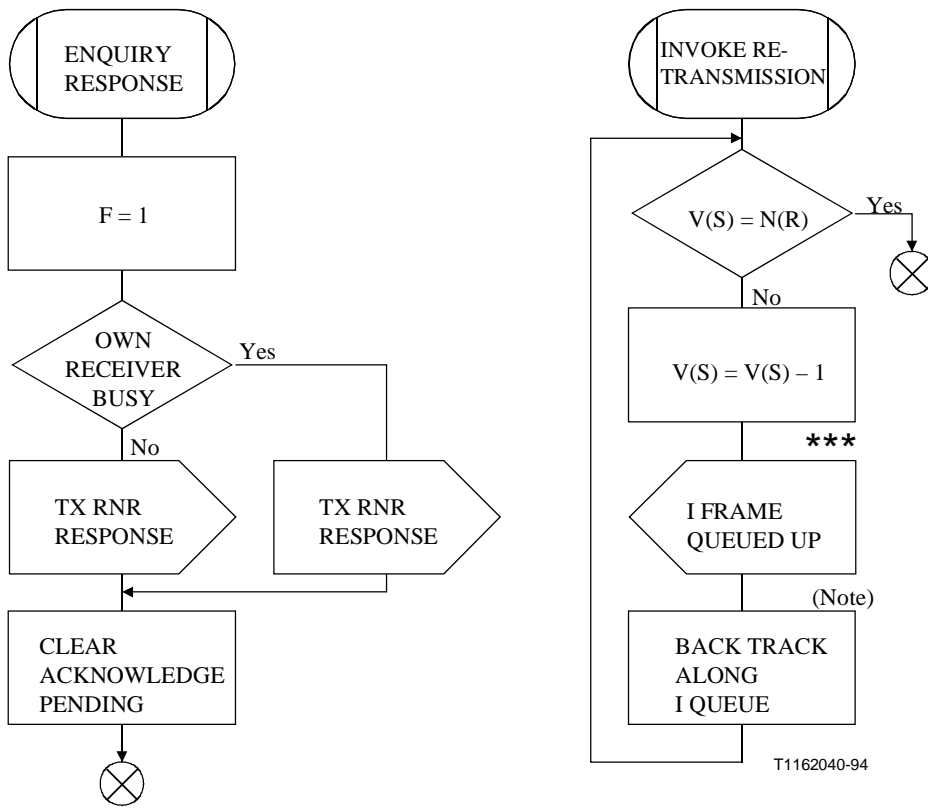


Figure B.9/Q.921 (sheet 4 of 5)

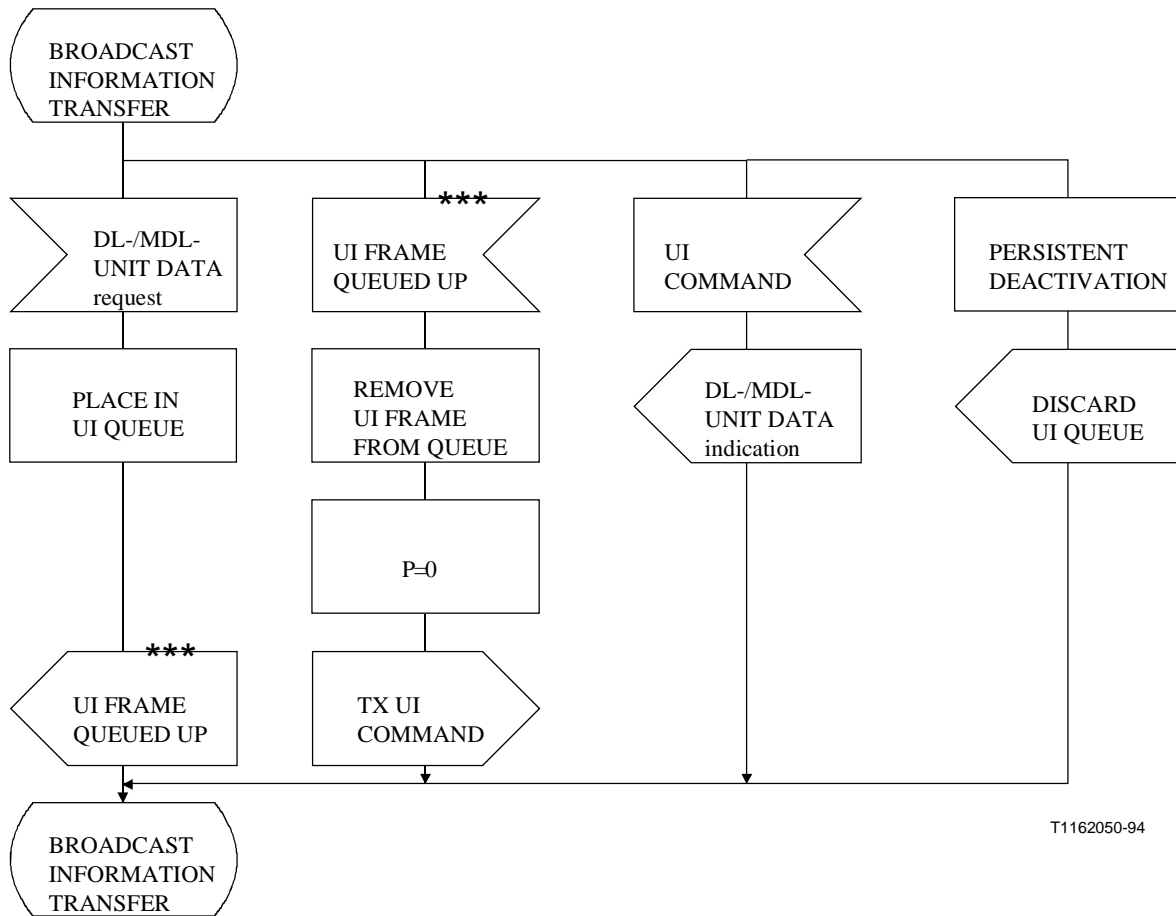


NOTE – The generation of the correct number of signals in order to cause the required retransmission of I frames does not alter their sequence integrity.

Figure B.9/Q.921 (sheet 5 of 5)

ANNEX C

SDL representation of the broadcast procedures



T1162050-94

Figure C.1/Q.921

ANNEX D

State transition table of the point-to-point procedures of the data link layer

D.1 The state transition table presented in Tables D.1 to D.3 is based on the eight basic states (see B.2) recognized in the SDL representation and the related transmitter and receiver conditions.

The state transition table relinquishes to any partitioning of the procedures. It is conceptual and does not prevent a designer from partitioning in his implementation. Moreover, all the processes related to primitive procedures, the management of queues and the exchange of information between adjacent layers are conceptual, not visible from outside of the system and would not impose any constraints on the implementation.

The eight basic states apply to both the transmitter and the receiver within one data link layer entity. However, some of the conditions are confined to the transmitter (e.g. "peer receiver busy"), whilst some are confined to the receiver (e.g. "REJ recovery"). This implies, if the concept of non-

partitioning is adopted, that each transmitter condition has to be combined with each receiver condition resulting in composite states. This state transition table comprises 24 composite states representing the 8 basic states and the related combinations of transmitter and receiver conditions.

Events are defined as follows:

- a) primitives;
- b) repertoire of frames to be received:
 - unnumbered frames (SABME, DISC, UA, DM, UI, FRMR),
 - supervisory frames (RR, REJ, RNR),
 - information frame (I);
- c) internal events (servicing of queues, expiry of timers, receiver busy condition).

The actions to be taken when an event occurs whilst in a specific state comprise:

- i) transitioning to another state;
- ii) transmitting peer-to-peer frames;
- iii) issuing primitives;
- iv) setting timers;
- v) setting counters;
- vi) updating variables;
- vii) setting P/F bit;
- viii) discarding contents of queues.

D.2 Key to the state transition table

D.2.1 Definition of a cell of the state transition table

	STATE	X defines the transition to the next state X empty indicates "remain in the current state"
EVENT	ACTIONS X	

D.2.2 Key to the contents of a cell

- | Impossible by the definition of the data link layer service.
 - / Impossible by the definition of the peer-to-peer data link procedures.
 - No action, no state change.
- V(S) = V(A) = N(R) Collective term for the two actions V(S) = N(R) and V(A) = N(R).
- Timer T200 Start timer T200 if not already running.
- TX ACK The acknowledgement of the received I frame may be conveyed by an I frame associated with the information flow in the opposite direction or a supervisory response frame, as appropriate.
- "DISCARD" Indicates the discarding of the information contained in the information field of the I frame.

(A-O)

The codes used in MDL-ERROR indication signals are defined in Table II.1. When multiple codes are shown, only one applies.

The action

A	

indicates

A	A
A	A

NOTE – In general, this state transition table does not prevent an implementation from using N(R) to acknowledge more than one I frame.

See Tables D.1 to D.3

Table D.1/Q.921 (sheet 1 of 10) – State transition table: Receiving primitive

BASIC STATE	TEI UNASSIGNED	ASSIGNED AWAITING TEI	ESTABLISH AWAITING TEI	TEI ASSIGNED	AWAITING ESTABLISHMENT			AWAITING RELEASE
					Establish	Re-establish	Pending release (note)	
TRANSMITTER CONDITION								
RECEIVER CONDITION								
STATE NUMBER	1	2	3	4	5.0	5.1	5.2	6
DL-ESTABLISH request	MDL-ASS ind 3	3		RC = 0 TX SABME P = 1 START T200 5.0		DISC, I QUEUE 5.0		
DL-RELEASE request				DL-REL conf		5.2		
DL-DATA request						DATA INTO I QUEUE		
I FRAME IN QUEUE $V(S) < V(A) + k$						LEAVE I FRAME IN QUEUE		
I FRAME IN QUEUE $V(S) = V(A) + k$								
DL-UNIT DATA request	MDL-ASS ind UNIT DATA INTO UI QUEUE 2	UNIT DATA INTO UI QUEUE						
UI FRAME IN QUEUE		LEAVE UI FRAME IN QUEUE		TX UI P = 0				

Table D.1/Q.921 (sheet 1 of 10) – State transition table: Receiving primitive (concluded)

BASIC STATE	TEI UNASSIGNED	ASSIGNED AWAITING TEI	ESTABLISH AWAITING TEI	TEI ASSIGNED	AWAITING ESTABLISHMENT			AWAITING RELEASE
					Establish	Re-establish	Pending release (note)	
TRANSMITTER CONDITION								
RECEIVER CONDITION								
STATE NUMBER	1	2	3	4	5.0	5.1	5.2	6
MDL-ASSIGN request	STORE TEI VALUE 4		STORE TEI VALUE RC = 0 TX SABME P = 1 START T200 5.0					
MDL-REMOVE request				DISC. UI QUEUE 1	DL-REL ind DISC. UI QUEUE STOP T200 1	DL-REL ind DISC. I and UI QUEUES STOP T200 1	DL-REL conf DISC. I and UI QUEUES STOP T200 1	DL-REL conf DISC. UI QUEUE STOP T200 1
MDL-ERROR response		DISC. UI QUEUE 1	DL-REL ind DISC. UI QUEUE 1					
PERSISTENT DEACTIVATION	–	DISC. UI QUEUE 1	DL-REL ind DISC. UI QUEUE 1	DISC. UI QUEUE	DL-REL ind DISC. UI QUEUE STOP T200 4	DL-REL ind DISC. I and UI QUEUES STOP T200 4	DL-REL conf DISC. I and UI QUEUES STOP T200 4	DL-REL conf DISC. UI QUEUE STOP T200 4

NOTE – The transmitter condition "pending release" may occur only in cases of layer 2 initiated re-establishment.

Table D.1/Q.921 (sheet 2 of 10) – State transition table: Receiving unnumbered frame with correct format

BASIC STATE	TEI UNASSIGNED	ASSIGNED AWAITING TEI	ESTABLISH AWAITING TEI	TEI ASSIGNED	AWAITING ESTABLISHMENT			AWAITING RELEASE
					Establish	Re-establish	Pending release	
TRANSMITTER CONDITION								
RECEIVER CONDITION								
STATE NUMBER	1	2	3	4	5.0	5.1	5.2	6
SABME P = 1 ABLE TO ENTER STATE 7.0	/	/	/	DL-EST ind V(S,R,A) = 0 TX UA F = 1 START T203 7.0	TX UA F = 1			TX DM F = 1
SABME P = 1 UNABLE TO ENTER STATE 7.0	/	/	/	TX DM F = 1	/	/	/	/
SABME P = 0 ABLE TO ENTER STATE 7.0	/	/	/	DL-EST ind V(S,R,A) = 0 TX UA F = 0 START T203 7.0	TX UA F = 0			TX DM F = 0
SABME P = 0 UNABLE TO ENTER STATE 7.0	/	/	/	TX DM F = 0	/	/	/	/
DISC P = 1	/	/	/	TX DM F = 1				TX UA F = 1
DISC P = 0	/	/	/	TX DM F = 0				TX UA F = 0
UA F = 1 V(S) = V(A)	/	/	/	MDL-ERR ind(C)	V(S,R,A) = 0 DL-EST conf STOP T200 START T203 7.0	V(S,R,A) = 0 STOP T200 START T203 7.0	DISC I QUEUE RC = 0 TX DISC P = 1 RESTART T200 6	DL-REL conf STOP T200 4

Table D.1/Q.921 (sheet 2 of 10) – State transition table: Receiving unnumbered frame with correct format (concluded)

BASIC STATE	TEI UNASSIGNED	ASSIGNED AWAITING TEI	ESTABLISH AWAITING TEI	TEI ASSIGNED	AWAITING ESTABLISHMENT			AWAITING RELEASE
					Establish	Re-establish	Pending release	
TRANSMITTER CONDITION								
RECEIVER CONDITION								
STATE NUMBER	1	2	3	4	5.0	5.1	5.2	6
UA F = 1 V(S) ≠ V(A)	/	/	/			DISC I QUEUE V(S,R,A) = 0 DL-EST ind STOP T200 START T203 7.0		
UA F = 0	/	/	/	MDL-ERR ind(D)				
DM F = 1	/	/	/	–	DL-REL ind STOP T200 4	DL-REL ind DISC I QUEUE STOP T200 4	DL-REL conf DISC I QUEUE STOP T200 4	DL-REL conf STOP T200 4
DM F = 0 ABLE TO ENTER STATE 7.0	/	/	/	RC = 0 TX SABME P = 1 START T200 5.1	–	–	–	–
DM F = 0 UNABLE TO ENTER STATE 7.0	/	/	/	–	/	/	/	/
UI command	/	/	/	DL-UNIT DATA ind				

Table D.1/Q.921 (sheet 3 of 10) – State transition table: Receiving FRMR unnumbered frame with correct format

BASIC STATE	TEI UNASSIGNED	ASSIGNED AWAITING TEI	ESTABLISH AWAITING TEI	TEI ASSIGNED	AWAITING ESTABLISHMENT			AWAITING RELEASE
					Establish	Re-establish	Pending release	
TRANSMITTER CONDITION								
RECEIVER CONDITION								
STATE NUMBER	1	2	3	4	5.0	5.1	5.2	6
FRMR response rejecting SABME	/	/	/	/	–	–	–	/
FRMR response rejecting DISC	/	/	/	/	/	/	/	–
FRMR response rejecting UA	/	/	/	–	–	–	–	–
FRMR response rejecting DM	/	/	/	–	–	–	–	–
FRMR response rejecting I command	/	/	/	/	–	–	–	–
FRMR response rejecting S frame	/	/	/	/	–	–	–	–
FRMR response rejecting FRMR	/	/	/	/	/	/	/	/

Table D.1/Q.921 (sheet 4 of 10) – State transition table: Receiving RR unnumbered frame with correct format

BASIC STATE	TEI UNASSIGNED	ASSIGNED AWAITING TEI	ESTABLISH AWAITING TEI	TEI ASSIGNED	AWAITING ESTABLISHMENT			AWAITING RELEASE
					Establish	Re-establish	Pending release	
TRANSMITTER CONDITION								
RECEIVER CONDITION								
STATE NUMBER	1	2	3	4	5.0	5.1	5.2	6
RR command P = 1	/	/	/	–	–	–	–	–
RR command P = 0	/	/	/	–	–	–	–	–
RR response F = 0	/	/	/	–	–	–	–	–
RR response F = 1	/	/	/	–	–	–	–	–

Table D.1/Q.921 (sheet 5 of 10) – State transition table: Receiving REJ supervisory frame with correct format

BASIC STATE	TEI UNASSIGNED	ASSIGNED AWAITING TEI	ESTABLISH AWAITING TEI	TEI ASSIGNED	AWAITING ESTABLISHMENT			AWAITING RELEASE
					Establish	Re-establish	Pending release	
TRANSMITTER CONDITION								
RECEIVER CONDITION								
STATE NUMBER	1	2	3	4	5.0	5.1	5.2	6
REJ command P = 1	/	/	/	–	–	–	–	–
REJ command P = 0	/	/	/	–	–	–	–	–
REJ response F = 0	/	/	/	–	–	–	–	–
REJ response F = 1	/	/	/	–	–	–	–	–

Table D.1/Q.921 (sheet 6 of 10) – State transition table: Receiving RNR supervisory frame with correct format

BASIC STATE	TEI UNASSIGNED	ASSIGNED AWAITING TEI	ESTABLISH AWAITING TEI	TEI ASSIGNED	AWAITING ESTABLISHMENT			AWAITING RELEASE
					Establish	Re-establish	Pending release	
TRANSMITTER CONDITION								
RECEIVER CONDITION								
STATE NUMBER	1	2	3	4	5.0	5.1	5.2	6
RNR command P = 1	/	/	/	–	–	–	–	–
RNR command P = 0	/	/	/	–	–	–	–	–
RNR response F = 0	/	/	/	–	–	–	–	–
RNR response F = 1	/	/	/	–	–	–	–	–

Table D.1/Q.921 (sheet 7 of 10) – State transition table: Receiving I command frame with correct format acknowledging all outstanding I frames or containing an N(R) which satisfies $V(A) < N(R) < V(S)$

BASIC STATE	TEI UNASSIGNED	ASSIGNED AWAITING TEI	ESTABLISH AWAITING TEI	TEI ASSIGNED	AWAITING ESTABLISHMENT			AWAITING RELEASE
					Establish	Re-establish	Pending release	
TRANSMITTER CONDITION								
RECEIVER CONDITION								
STATE NUMBER	1	2	3	4	5.0	5.1	5.2	6
I command P = 1 N(S) = V(R) N(R) = V(S)	/	/	/	-	-	-	-	-
I command P = 0 N(S) = V(R) N(R) = V(S)	/	/	/	-	-	-	-	-
I command P = 1 N(S) ≠ V(R) N(R) = V(S)	/	/	/	-	-	-	-	-
I command P = 0 N(S) ≠ V(R) N(R) = V(S)	/	/	/	-	-	-	-	-
I command P = 1 N(S) = V(R) $V(A) < N(R) < V(S)$	/	/	/	-	-	-	-	-
I command P = 0 N(S) = V(R) $V(A) < N(R) < V(S)$	/	/	/	-	-	-	-	-
I command P = 1 N(S) ≠ V(R) $V(A) < N(R) < V(S)$	/	/	/	-	-	-	-	-
I command P = 0 N(S) ≠ V(R) $V(A) < N(R) < V(S)$	/	/	/	-	-	-	-	-

Table D.1/Q.921 (sheet 8 of 10) – State transition table: Receiving I command frame with correct format containing an N(R) which satisfies $V(A) = N(R) < V(S)$, or an N(R) error

BASIC STATE	TEI UNASSIGNED	ASSIGNED AWAITING TEI	ESTABLISH AWAITING TEI	TEI ASSIGNED	AWAITING ESTABLISHMENT			AWAITING RELEASE
					Establish	Re-establish	Pending release	
TRANSMITTER CONDITION								
RECEIVER CONDITION								
STATE NUMBER	1	2	3	4	5.0	5.1	5.2	6
I command P = 1 N(S) = V(R) V(A) = N(R) < V(S)	/	/	/	-	-	-	-	-
I command P = 0 N(S) = V(R) V(A) = N(R) < V(S)	/	/	/	-	-	-	-	-
I command P = 1 N(S) ≠ V(R) V(A) = N(R) < V(S)	/	/	/	-	-	-	-	-
I command P = 0 N(S) ≠ V(R) V(A) = N(R) < V(S)	/	/	/	-	-	-	-	-
I command P = 1 N(S) = V(R) N(R) error	/	/	/	-	-	-	-	-
I command P = 0 N(S) = V(R) N(R) error	/	/	/	-	-	-	-	-
I command P = 1 N(S) ≠ V(R) N(R) error	/	/	/	-	-	-	-	-
I command P = 0 N(S) ≠ V(R) N(R) error	/	/	/	-	-	-	-	-

Table D.1/Q.921 (sheet 9 of 10) – State transition table: Internal events (expiry of timers, receiver busy condition)

BASIC STATE	TEI UNASSIGNED	ASSIGNED AWAITING TEI	ESTABLISH AWAITING TEI	TEI ASSIGNED	AWAITING ESTABLISHMENT			AWAITING RELEASE
					Establish	Re-establish	Pending release	
TRANSMITTER CONDITION								
RECEIVER CONDITION								
STATE NUMBER	1	2	3	4	5.0	5.1	5.2	6
T200 TIME-OUT RC < N200	/	/	/	/	RC = RC + 1 TX SABME P = 1 START T200			RC = RC + 1 TX DISC P = 1 START T200
T200 TIME-OUT RC = N200	/	/	/	/	DL-REL ind MDL-ERR ind(G) 4	DISC. I QUEUE DL-REL ind MDL-ERR ind(G) 4	DISC. I QUEUE DL-REL conf MDL-ERR ind(G) 4	DL-REL conf MDL-ERR ind(H) 4
T203 TIME-OUT	/	/	/	/	/	/	/	/
SET OWN RECEIVER BUSY (Note)	/	/	/	/	/	/	/	/
CLEAR OWN RECEIVER BUSY (Note)	/	/	/	/	/	/	/	/
NOTE – These signals are generated outside the procedures specified in this state transition table, and may be generated by the connection management entity.								

Table D.1/Q.921 (sheet 10 of 10) – State transition table: Receiving frame with incorrect format or frame not implemented

BASIC STATE	TEI UNASSIGNED	ASSIGNED AWAITING TEI	ESTABLISH AWAITING TEI	TEI ASSIGNED	AWAITING ESTABLISHMENT			AWAITING RELEASE
					Establish	Re-establish	Pending release	
TRANSMITTER CONDITION								
RECEIVER CONDITION								
STATE NUMBER	1	2	3	4	5.0	5.1	5.2	6
SABME incorrect length	/	/	/	MDL-ERR ind(N)				
DISC incorrect length	/	/	/					
UA incorrect length	/	/	/					
DM incorrect length	/	/	/					
FRMR incorrect length	/	/	/					
Supervisory frame RR, REJ, RNR incorrect length	/	/	/					
N201 error	/	/	/	MDL-ERR ind(O)				
Undefined command and response frames	/	/	/	MDL-ERR ind(L)				
I field not permitted	/	/	/	MDL-ERR ind(M)				

Table D.2/Q.921 (sheet 1 of 10) – State transition table: Receiving primitive

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
DL-ESTABLISH request	DISC I QUEUE RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.0							
DL-RELEASE request	DISC I QUEUE RC = 0 TX DISC P = 1 STOP T203 RESTART T200 6							
DL-DATA request	DATA INTO I QUEUE							
I FRAME IN QUEUE $V(S) < V(A) + k$	TX I P = 0 $V(S) = V(S) + 1$ STOP T203 TIMER T200					LEAVE I FRAME IN QUEUE		
I FRAME IN QUEUE $V(S) = V(A) + k$	LEAVE I FRAME IN QUEUE							
DL-UNIT DATA request	UNIT DATA INTO UI QUEUE							

Table D.2/Q.921 (sheet 1 of 10) – State transition table: Receiving primitive (concluded)

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
UI FRAME IN QUEUE	TX UI P = 0							
MDL-ASSIGN request								
MDL-REMOVE request	DL-REL ind DISC I and UI QUEUES STOP T200 STOP T203 1							
MDL-ERROR response								
PERSISTENT DEACTIVATION	DL-REL ind DISC I and UI QUEUES STOP T200 STOP T203 4							

Table D.2/Q.921 (sheet 2 of 10) – State transition table: Receiving unnumbered frame with correct format

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
SABME P = 1 V(S) = V(A)	MDL-ERR ind(F) V(S,R,A) = 0 TX UA F = 1 STOP T200 START T203	MDL-ERR ind(F) V(S,R,A) = 0 TX UA F = 1 STOP T200 START T203	7.0					
SABME P = 1 V(S) ≠ V(A)	DL-EST ind MDL-ERR ind(F) DISC I QUEUE V(S,R,A) = 0 TX UA F = 1 STOP T200 START T203	DL-EST ind MDL-ERR ind(F) DISC I QUEUE V(S,R,A) = 0 TX UA F = 1 STOP T200 START T203	7.0					
SABME P = 0 V(S) = V(A)	MDL-ERR ind(F) V(S,R,A) = 0 TX UA F = 0 STOP T200 START T203	MDL-ERR ind(F) V(S,R,A) = 0 TX UA F = 0 STOP T200 START T203	7.0					

Table D.2/Q.921 (sheet 2 of 10) – State transition table: Receiving unnumbered frame with correct format (concluded)

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
SABME P = 0 V(S) ≠ V(A)	DL-EST ind MDL-ERR ind(F) DISC I QUEUE V(S,R,A) = 0 TX UA F = 0 STOP T200 START T203	DL-EST ind MDL-ERR ind(F) DISC I QUEUE V(S,R,A) = 0 TX UA F = 0 STOP T200 START T203 7.0						
DISC P = 1	DL-REL ind DISC I QUEUE TX UA F = 1 STOP T200, T203 4							
DISC P = 0	DL-REL ind DISC I QUEUE TX UA F = 0 STOP T200, T203 4							
UA F = 1	MDL-ERR ind(C)							
UA F = 0	MDL-ERR ind(D)							
DM F = 1	MDL-ERR ind(B)							
DM F = 0	MDL-ERR ind(E) RC = 0 TX SABME P = 1 STOP T200 RESTART T203 5.1				MDL-ERR ind(E) RC = 0 TX SABME P = 1 RESTART T200 5.1			
UI command	DL-UNIT DATA ind							

Table D.2/Q.921 (sheet 3 of 10) – State transition table: Receiving FRMR unnumbered frame with correct format

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
FRMR response rejecting SABME	/	/	/	/	/	/	/	/
FRMR response rejecting DISC	/	/	/	/	/	/	/	/
FRMR response rejecting UA	MDL-ERR ind(K) RC = 0 TX SABME P = 1 STOP T200 RESTART T203 5.1				MDL-ERR ind(K) RC = 0 TX SABME P = 1 RESTART T200 5.1			
FRMR response rejecting DM	/	/	/	/	/	/	/	/
FRMR response rejecting I command	MDL-ERR ind(K) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1				MDL-ERR ind(K) RC = 0 TX SABME P = 1 RESTART T200 5.1			
FRMR response rejecting S frame								
FRMR response rejecting FRMR	/	/	/	/	/	/	/	/

Table D.2/Q.921 (sheet 4 of 10) – State transition table: Receiving RR supervisory frame with correct format

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
RR command P = 1 N(R) = V(S)	TX RR F = 1 STOP T200 RESTART T203 V(A) = N(R)		TX RNR F = 1 STOP T200 RESTART T203 V(A) = N(R)		TX RR F = 1 STOP T200 START T203 V(A) = N(R) 7.0	TX RR F = 1 STOP T200 START T203 V(A) = N(R) 7.1	TX RNR F = 1 STOP T200 START T203 V(A) = N(R) 7.2	TX RNR F = 1 STOP T200 START T203 V(A) = N(R) 7.3
RR command P = 0 N(R) = V(S)	STOP T200 RESTART T203 V(A) = N(R)				STOP T200 START T203 V(A) = N(R) 7.0	STOP T200 START T203 V(A) = N(R) 7.1	STOP T200 START T203 V(A) = N(R) 7.2	STOP T200 START T203 V(A) = N(R) 7.3
RR response F = 0 N(R) = V(S)								
RR response F = 1 N(R) = V(S)	MDL-ERR ind(A) STOP T200 RESTART T203 V(A) = V(R)				MDL-ERR ind(A) STOP T200 START T203 V(A) = V(R) 7.0	MDL-ERR ind(A) STOP T200 START T203 V(A) = V(R) 7.1	MDL-ERR ind(A) STOP T200 START T203 V(A) = V(R) 7.2	MDL-ERR ind(A) STOP T200 START T203 V(A) = V(R) 7.3
RR command P = 1 V(A) < N(R) < V(S)	TX RR F = 1 RESTART T200 V(A) = N(R)		TX RNR F = 1 RESTART T200 V(A) = N(R)		TX RR F = 1 RESTART T200 V(A) = N(R) 7.0	TX RR F = 1 RESTART T200 V(A) = N(R) 7.1	TX RNR F = 1 RESTART T200 V(A) = N(R) 7.2	TX RR F = 1 RESTART T200 V(A) = N(R) 7.3
RR command P = 0 V(A) < N(R) < V(S)	RESTART T200 V(A) = N(R)				RESTART T200 V(A) = N(R) 7.0	RESTART T200 V(A) = N(R) 7.1	RESTART T200 V(A) = N(R) 7.2	RESTART T200 V(A) = N(R) 7.3
RR response F = 0 V(A) < N(R) < V(S)								
RR response F = 1 V(A) < N(R) < V(S)	MDL-ERR ind(A) RESTART T200 V(A) = N(R)				MDL-ERR ind(A) RESTART T200 V(A) = N(R) 7.0	MDL-ERR ind(A) RESTART T200 V(A) = N(R) 7.1	MDL-ERR ind(A) RESTART T200 V(A) = N(R) 7.2	MDL-ERR ind(A) RESTART T200 V(A) = N(R) 7.3
RR command P = 1 V(A) = N(R) < V(S)	TX RR F = 1		TX RNR F = 1		TX RR F = 1 7.0	TX RR F = 1 7.1	TX RNR F = 1 7.2	TX RNR F = 1 7.3
RR command P = 0 V(A) = N(R) < V(S)	–	–	–	–	7.0	7.1	7.2	7.3

Table D.2/Q.921 (sheet 4 of 10) – State transition table: Receiving RR supervisory frame with correct format (concluded)

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
RR response F = 0 V(A) = N(R) < V(S)	–	–	–	–				
RR response F = 1 V(A) = N(R) < V(S)	MDL-ERR ind(A)				MDL-ERR ind(A) 7.0	MDL-ERR ind(A) 7.1	MDL-ERR ind(A) 7.2	MDL-ERR ind(A) 7.3
RR command P = 1 N(R) error	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1		TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1		TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	
RR command P = 0 N(R) error	MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1				MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1			
RR response F = 0 N(R) error								
RR response F = 1 N(R) error	MDL-ERR ind(A) MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1				MDL-ERR ind(A) MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1			

Table D.2/Q.921 (sheet 5 of 10) – State transition table: Receiving REJ supervisory frame with correct format

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
REJ command P = 1 N(R) = V(S) (Note)	TX RR F = 1 V(A) = N(R) STOP T200 RESTART T203		TX RNR F = 1 V(A) = N(R) STOP T200 RESTART T203		TX RR F = 1 V(A) = N(R) STOP T200 START T203 7.0	TX RR F = 1 V(A) = N(R) STOP T200 START T203 7.1	TX RNR F = 1 V(A) = N(R) STOP T200 START T203 7.2	TX RNR F = 1 V(A) = N(R) STOP T200 START T203 7.3
REJ command P = 0 N(R) = V(S) (Note)	V(A) = N(R) STOP T200 RESTART T203				V(A) = N(R) STOP T200 START T203 7.0	V(A) = N(R) STOP T200 START T203 7.1	V(A) = N(R) STOP T200 START T203 7.2	V(A) = N(R) STOP T200 START T203 7.3
REJ response F = 0 N(R) = V(S) (Note)								
REJ response F = 1 N(R) = V(S) (Note)	MDL-ERR ind(A) V(A) = N(R) STOP T200 RESTART T203				MDL-ERR ind(A) V(A) = N(R) STOP T200 START T203 7.0	MDL-ERR ind(A) V(A) = N(R) STOP T200 START T203 7.1	MDL-ERR ind(A) V(A) = N(R) STOP T200 START T203 7.2	MDL-ERR ind(A) V(A) = N(R) STOP T200 START T203 7.3
REJ command P = 1 V(A) ≤ N(R) < V(S)	TX RR F = 1 V(S) = V(A) = N(R) STOP T200 START T203		TX RNR F = 1 V(S) = V(A) = N(R) STOP T200 START T203		TX RR F = 1 V(S) = V(A) = N(R) STOP T200 START T203 7.0	TX RR F = 1 V(S) = V(A) = N(R) STOP T200 START T203 7.1	TX RNR F = 1 V(S) = V(A) = N(R) STOP T200 START T203 7.2	TX RNR F = 1 V(S) = V(A) = N(R) STOP T200 START T203 7.3
REJ command P = 0 V(A) ≤ N(R) < V(S)	V(S) = V(A) = N(R) STOP T200 START T203				V(S) = V(A) = N(R) STOP T200 START T203 7.0	V(S) = V(A) = N(R) STOP T200 START T203 7.1	V(S) = V(A) = N(R) STOP T200 START T203 7.2	V(S) = V(A) = N(R) STOP T200 START T203 7.3
REJ response F = 0 V(A) ≤ N(R) < V(S)								
REJ response F = 1 V(A) ≤ N(R) < V(S)	MDL-ERR ind(A) V(S) = V(A) = N(R) STOP T200 START T203				MDL-ERR ind(A) V(S) = V(A) = N(R) STOP T200 START T203 7.0	MDL-ERR ind(A) V(S) = V(A) = N(R) STOP T200 START T203 7.1	MDL-ERR ind(A) V(S) = V(A) = N(R) STOP T200 START T203 7.2	MDL-ERR ind(A) V(S) = V(A) = N(R) STOP T200 START T203 7.3

Table D.2/Q.921 (sheet 5 of 10) – State transition table: Receiving REJ supervisory frame with correct format (concluded)

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
REJ command P = 1 N(R) error	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1		TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1		TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	
REJ command P = 0 N(R) error	MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1				MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1			
REJ response F = 0 N(R) error								
REJ response F = 1 N(R) error	MDL-ERR ind(A) MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1				MDL-ERR ind(A) MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1			

NOTE – This event is impossible by the definition of the peer-to-peer data link procedures. However, it would not harm the information transfer, if actions according to this table are taken.

Table D.2/Q.921 (sheet 6 of 10) – State transition table: Receiving RNR supervisory frame with correct format

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
RNR command P = 1 N(R) = V(S)	TX RR F = 1 STOP T203 RESTART T200 V(A) = N(R) 7.4	TX RR F = 1 STOP T203 RESTART T200 V(A) = N(R) 7.5	TX RNR F = 1 STOP T203 RESTART T200 V(A) = N(R) 7.6	TX RNR F = 1 STOP T203 RESTART T200 V(A) = N(R) 7.7	TX RR F = 1 RESTART T200 V(A) = N(R)		TX RNR F = 1 RESTART T200 V(A) = N(R)	
RNR command P = 0 N(R) = V(S)	STOP T203 RESTART T200 V(A) = N(R) 7.4	STOP T203 RESTART T200 V(A) = N(R) 7.5	STOP T203 RESTART T200 V(A) = N(R) 7.6	STOP T203 RESTART T200 V(A) = N(R) 7.7	RESTART T200 V(A) = N(R)			
RNR response F = 0 N(R) = V(S)								
RNR response F = 1 N(R) = V(S)	MDL-ERR ind(A) STOP T203 RESTART T200 V(A) = N(R) 7.4	MDL-ERR ind(A) STOP T203 RESTART T200 V(A) = N(R) 7.5	MDL-ERR ind(A) STOP T203 RESTART T200 V(A) = N(R) 7.6	MDL-ERR ind(A) STOP T203 RESTART T200 V(A) = N(R) 7.7	MDL-ERR ind(A) RESTART T200 V(A) = N(R)			
RNR command P = 1 V(A) ≤ N(R) < V(S)	TX RR F = 1 RESTART T200 V(A) = N(R) 7.4	TX RR F = 1 RESTART T200 V(A) = N(R) 7.5	TX RNR F = 1 RESTART T200 V(A) = N(R) 7.6	TX RNR F = 1 RESTART T200 V(A) = N(R) 7.7	TX RR F = 1 RESTART T200 V(A) = N(R)		TX RNR F = 1 RESTART T200 V(A) = N(R)	
RNR command P = 0 V(A) ≤ N(R) < V(S)	RESTART T200 V(A) = N(R) 7.4	RESTART T200 V(A) = N(R) 7.5	RESTART T200 V(A) = N(R) 7.6	RESTART T200 V(A) = N(R) 7.7	RESTART T200 V(A) = N(R)			
RNR response F = 0 V(A) ≤ N(R) < V(S)								
RNR response F = 1 V(A) ≤ N(R) < V(S)	MDL-ERR ind(A) RESTART T200 V(A) = N(R) 7.4	MDL-ERR ind(A) RESTART T200 V(A) = N(R) 7.5	MDL-ERR ind(A) RESTART T200 V(A) = N(R) 7.6	MDL-ERR ind(A) RESTART T200 V(A) = N(R) 7.7	MDL-ERR ind(A) RESTART T200 V(A) = N(R)			

Table D.2/Q.921 (sheet 6 of 10) – State transition table: Receiving RNR supervisory frame with correct format (concluded)

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
RNR command N(R) error P = 1	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1		TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1		TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	
RNR command N(R) error P = 0	MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1				MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1			
RNR response N(R) error F = 0								
RNR response N(R) error F = 1	MDL-ERR ind(A) MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1				MDL-ERR ind(A) MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1			

Table D.2/Q.921 (sheet 7 of 10) – State transition table: Receiving I command frame with correct format acknowledging all outstanding I frames or containing an N(R) which satisfies $V(A) < N(R) < V(S)$

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
I command N(S) = V(R) N(R) = V(S) P = 1	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 STOP T200 RESTART T203 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 STOP T200 RESTART T203 V(A) = N(R) 7.0	"DISCARD" TX RNR F = 1 STOP T200 RESTART T203 V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R) 7.4	"DISCARD" TX RNR F = 1 V(A) = N(R)	
I command N(S) = V(R) N(R) = V(S) P = 0	V(R) = V(R) + 1 DL-DATA ind TX ACK STOP T200 RESTART T203 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX ACK STOP T200 RESTART T203 V(A) = N(R) 7.0	"DISCARD" STOP T200 RESTART T203 V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R) 7.4	"DISCARD" V(A) = N(R)	
I command N(S) ≠ V(R) N(R) = V(S) P = 1	"DISCARD" TX REJ F = 1 STOP T200 RESTART T203 V(A) = N(R) 7.1	"DISCARD" TX RR F = 1 STOP T200 RESTART T203 V(A) = N(R)	"DISCARD" TX RNR F = 1 STOP T200 RESTART T203 V(A) = N(R)		"DISCARD" TX REJ F = 1 V(A) = N(R) 7.5	"DISCARD" TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	
I command N(S) ≠ V(R) N(R) = V(S) P = 0	"DISCARD" TX REJ F = 0 STOP T200 RESTART T203 V(A) = N(R) 7.1	"DISCARD" STOP T200 RESTART T203 V(A) = N(R)			"DISCARD" TX REJ F = 0 V(A) = N(R) 7.5	"DISCARD" V(A) = N(R)		
I command N(S) = V(R) V(A) < N(R) < V(S) P = 1	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 RESTART T200 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 RESTART T200 V(A) = N(R) 7.0	"DISCARD" TX RNR F = 1 RESTART T200 V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R) 7.4	"DISCARD" TX RNR F = 1 V(A) = N(R)	

Table D.2/Q.921 (sheet 7 of 10) – State transition table: Receiving I command frame with correct format acknowledging all outstanding I frames or containing an N(R) which satisfies $V(A) < N(R) < V(S)$ (concluded)

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
I command P = 0 N(S) = V(R) V(A) < N(R) < V(S)	V(R) = V(R) + 1 DL-DATA ind TX ACK RESTART T200 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX ACK RESTART T200 V(A) = N(R)	"DISCARD" RESTART T200 V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)	
I command P = 1 N(S) ≠ V(R) V(A) < N(R) < V(S)	"DISCARD" TX REJ F = 1 RESTART T200 V(A) = N(R)	"DISCARD" TX RR F = 1 RESTART T200 V(A) = N(R)	"DISCARD" TX RNR F = 1 RESTART T200 V(A) = N(R)		"DISCARD" TX REJ F = 1 V(A) = N(R)	"DISCARD" TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	
I command P = 0 N(S) ≠ V(R) V(A) < N(R) < V(S)	"DISCARD" TX REJ F = 0 RESTART T200 V(A) = N(R)	"DISCARD" RESTART T200 V(A) = N(R)			"DISCARD" TX REJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)		

Table D.2/Q.921 (sheet 8 of 10) – State transition table: Receiving I command frame with correct format containing an N(R) which satisfies $V(A) = N(R) < V(S)$, or an N(R) error

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
I command P = 1 N(S) = V(R) V(A) = N(R) < V(S)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 7.0	"DISCARD" TX RNR F = 1		V(R) = V(R) + 1 DL-DATA ind TX RR F = 1	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 7.4	"DISCARD" TX RNR F = 1	
I command P = 0 N(S) = V(R) V(A) = N(R) < V(S)	V(R) = V(R) + 1 DL-DATA ind TX ACK	V(R) = V(R) + 1 DL-DATA ind TX ACK 7.0	"DISCARD"		V(R) = V(R) + 1 DL-DATA ind TX RR F = 0	V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 7.4	"DISCARD"	
I command P = 1 N(S) ≠ V(R) V(A) = N(R) < V(S)	"DISCARD" TX REJ F = 1 7.1	"DISCARD" TX RR F = 1	"DISCARD" TX RNR F = 1		"DISCARD" TX REJ F = 1 7.5	"DISCARD" TX RR F = 1	"DISCARD" TX RNR F = 1	
I command P = 0 N(S) ≠ V(R) V(A) = N(R) < V(S)	"DISCARD" TX REJ F = 0 7.1	"DISCARD"			"DISCARD" TX REJ F = 0 7.5	"DISCARD"		
I command P = 1 N(S) = V(R) N(R) error	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1		"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1		V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	
I command P = 0 N(S) = V(R) N(R) error	V(R) = V(R) + 1 DL-DATA ind MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1		"DISCARD" MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1		V(R) = V(R) + 1 DL-DATA ind MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		"DISCARD" MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	

Table D.2/Q.921 (sheet 8 of 10) – State transition table: Receiving I command frame with correct format containing an N(R) which satisfies $V(A) = N(R) < V(S)$, or an N(R) error (*concluded*)

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
I command N(S) ≠ V(R) N(R) error P = 1	"DISCARD" TX REJ F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1	"DISCARD" TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1		"DISCARD" TX REJ F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	"DISCARD" TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	
I command N(S) ≠ V(R) N(R) error P = 0	"DISCARD" TX REJ F = 0 MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1			"DISCARD" TX REJ F = 0 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 SABME P = 1 RESTART T200 5.1		

Table D.2/Q.921 (sheet 9 of 10) – State transition table: Internal events (expiry of timers, receiver busy condition)

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
T200 TIME-OUT RC < N200	RC = 0 either V(S) = V(S) - 1 TX I P = 1 V(S) = V(S) + 1 or TX RR P = 1 then RC = RC + 1 START T200 8.0	RC = 0 either V(S) = V(S) - 1 TX I P = 1 V(S) = V(S) + 1 or TX RR P = 1 then RC = RC + 1 START T200 8.1	RC = 0 either V(S) = V(S) - 1 TX I P = 1 V(S) = V(S) + 1 or TX RNR P = 1 then RC = RC + 1 START T200 8.2	RC = 0 either V(S) = V(S) - 1 TX I P = 1 V(S) = V(S) + 1 or TX RNR P = 1 then RC = RC + 1 START T200 8.3	RC = 0 TX RR P = 1 RC = RC + 1 START T200 8.4	RC = 0 TX RR P = 1 RC = RC + 1 START T200 8.5	RC = 0 TX RNR P = 1 RC = RC + 1 START T200 8.6	RC = 0 TX RNR P = 1 RC = RC + 1 START T200 8.7
T200 TIME-OUT RC = N200	/	/	/	/	/	/	/	/
T203 TIME-OUT	RC = 0 TX RR P = 1 START T200 8.0	RC = 0 TX RR P = 1 START T200 8.1	RC = 0 TX RNR P = 1 START T200 8.2	RC = 0 TX RNR P = 1 START T200 8.3	/	/	/	/
SET OWN RECEIVER BUSY (Note)	TX RNR F = 0 7.2	TX RNR F = 0 7.3	-	-	TX RNR F = 0 7.6	TX RNR F = 0 7.7	-	-
CLEAR OWN RECEIVER BUSY (Note)	-	-	TX RR F = 0 7.0	TX RR F = 0 7.1	-	-	TX RR F = 0 7.4	TX RR F = 0 7.5
NOTE – These signals are generated outside the procedures specified in this state transition table, and may be generated by the connection management entity.								

Table D.2/Q.921 (sheet 10 of 10) – State transition table: Receiving frame with incorrect format or frame not implemented

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
SABME incorrect length	MDL-ERR ind(N) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1							
DISC incorrect length								
UA incorrect length								
DM incorrect length								
FRMR incorrect length								
Supervisory frame RR, REJ, RNR incorrect length								
N201 error	MDL-ERR ind(O) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1							
Undefined command and response frames	MDL-ERR ind(L) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1							
I field not permitted	MDL-ERR ind(M) RC = 0 TX SABME P = 1 STOP T203 RESTART T200 5.1							

Table D.3/Q.921 (sheet 1 of 10) – State transition table: Receiving primitive

BASIC STATE	TIMER RECOVERY							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
DL-ESTABLISH request	DISC, I QUEUE RC = 0 TX SABME P = 1 RESTART T200 5.0							
DL-RELEASE request	DISC, I QUEUE RC = 0 TX DISC P = 1 RESTART T200 6							
DL-DATA request	DATA INTO I QUEUE							
I FRAME IN QUEUE $V(S) < V(A) + k$	LEAVE I FRAME IN QUEUE							
I FRAME IN QUEUE $V(S) = V(A) + k$								
DL-UNIT DATA request	UNIT DATA INTO UI QUEUE							
UI FRAME IN QUEUE	TX UI P = 0							
MDL-ASSIGN request								
MDL-REMOVE request	DL-REL ind DISC. I and UI QUEUES STOP T200 1							
MDL-ERROR response								
PERSISTENT DEACTIVATION	DL-REL ind DISC. I and UI QUEUES STOP T200 4							

Table D.3/Q.921 (sheet 2 of 10) – State transition table: Receiving unnumbered frame with correct format

BASIC STATE	TIMER RECOVERY							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
SABME P = 1 V(S) = V(A)	MDL-ERR ind(F) V(S,R,A) = 0 TX UA F = 1 STOP T200 START T203 7.0							
SABME P = 1 V(S) ≠ V(A)	DL-EST ind MDL-ERR ind(F) DISC. I QUEUE V(S,R,A) = 0 TX UA F = 1 STOP T200 START T203 7.0							
SABME P = 0 V(S) = V(A)	MDL-ERR ind(F) V(S,R,A) = 0 TX UA F = 0 STOP T200 START T203 7.0							
SABME P = 0 V(S) ≠ V(A)	DL-EST ind MDL-ERR ind(F) DISC. I QUEUE V(S,R,A) = 0 TX UA F = 0 STOP T200 START T203 7.0							
DISC P = 1	DL-REL ind DISC I QUEUE TX UA F = 1 STOP T200 4							

Table D.3/Q.921 (sheet 2 of 10) – State transition table: Receiving unnumbered frame with correct format (concluded)

BASIC STATE	TIMER RECOVERY							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
DISC P = 0	DL-REL ind DISC I QUEUE TX UA F = 0 STOP T200 4							
UA F = 1	MDL-ERR ind(C)							
UA F = 0	MDL-ERR ind(D)							
DM F = 1	MDL-ERR ind(B) RC = 0 TX SABME P = 1 RESTART T200 5.1							
DM F = 0	MDL-ERR ind(E) RC = 0 TX SABME P = 1 RESTART T200 5.1							
UI command	DL-UNIT DATA ind							

Table D.3/Q.921 (sheet 3 of 10) – State transition table: Receiving FRMR unnumbered frame with correct format

BASIC STATE	TIMER RECOVERY							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
FRMR response rejecting SABME	/	/	/	/	/	/	/	/
FRMR response rejecting DISC	/	/	/	/	/	/	/	/
FRMR response rejecting UA	/	/	/	/	/	/	/	/
FRMR response rejecting DM	/	/	/	/	/	/	/	/
FRMR response rejecting I command	MDL-ERR ind(K) RC = 0 TX SABME P = 1 RESTART T200 5.1							
FRMR response rejecting S frame								
FRMR response rejecting FRMR	/	/	/	/	/	/	/	/

Table D.3/Q.921 (sheet 4 of 10) – State transition table: Receiving RR supervisory frame with correct format, clearance of timer recovery if there is F = 1 only

BASIC STATE	TIMER RECOVERY							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
RR command P = 1 V(A) ≤ N(R) ≤ V(S)	TX RR F = 1 V(A) = N(R)		TX RNR F = 1 V(A) = N(R)		TX RR F = 1 V(A) = N(R) 8.0	TX RR F = 1 V(A) = N(R) 8.1	TX RNR F = 1 V(A) = N(R) 8.2	TX RNR F = 1 V(A) = N(R) 8.3
RR command P = 0 V(A) ≤ N(R) ≤ V(S)	V(A) = N(R)				V(A) = N(R) 8.0	V(A) = N(R) 8.1	V(A) = N(R) 8.2	V(A) = N(R) 8.3
RR response F = 0 V(A) ≤ N(R) ≤ V(S)								
RR response F = 1 V(A) ≤ N(R) ≤ V(S)	V(S) = N(R) STOP T200 START T203 V(A) = N(R) 7.0	V(S) = N(R) STOP T200 START T203 V(A) = N(R) 7.1	V(S) = N(R) STOP T200 START T203 V(A) = N(R) 7.2	V(S) = N(R) STOP T200 START T203 V(A) = N(R) 7.3	V(S) = N(R) STOP T200 START T203 V(A) = N(R) 7.0	V(S) = N(R) STOP T200 START T203 V(A) = N(R) 7.1	V(S) = N(R) STOP T200 START T203 V(A) = N(R) 7.2	V(S) = N(R) STOP T200 START T203 V(A) = N(R) 7.3
RR command N(R) error P = 1	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	
RR command N(R) error P = 0	MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1							
RR response N(R) error F = 0								
RR response N(R) error F = 1								

Table D.3/Q.921 (sheet 5 of 10) – State transition table: Receiving REJ supervisory frame with correct format, clearance of timer recovery if there is F = 1 only

BASIC STATE	TIMER RECOVERY							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
REJ command P = 1 V(A) ≤ N(R) ≤ V(S)	TX RR F = 1 V(A) = N(R)		TX RNR F = 1 V(A) = N(R)		TX RR F = 1 V(A) = N(R) 8.0	TX RR F = 1 V(A) = N(R) 8.1	TX RNR F = 1 V(A) = N(R) 8.2	TX RNR F = 1 V(A) = N(R) 8.3
REJ command P = 0 V(A) ≤ N(R) ≤ V(S)	V(A) = N(R)				V(A) = N(R) 8.0	V(A) = N(R) 8.1	V(A) = N(R) 8.2	V(A) = N(R) 8.3
REJ response F = 0 V(A) ≤ N(R) ≤ V(S)								
REJ response F = 1 V(A) ≤ N(R) ≤ V(S)	V(S) = V(A) = N(R) STOP T200 START T203 7.0	V(S) = V(A) = N(R) STOP T200 START T203 7.1	V(S) = V(A) = N(R) STOP T200 START T203 7.2	V(S) = V(A) = N(R) STOP T200 START T203 7.3	V(S) = V(A) = N(R) STOP T200 START T203 7.0	V(S) = V(A) = N(R) STOP T200 START T203 7.1	V(S) = V(A) = N(R) STOP T200 START T203 7.2	V(S) = V(A) = N(R) STOP T200 START T203 7.3
REJ command P = 1 N(R) error	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	
REJ command P = 0 N(R) error	MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1							
REJ response F = 0 N(R) error								
REJ response F = 1 N(R) error								

Table D.3/Q.921 (sheet 6 of 10) – State transition table: Receiving RNR supervisory frame with correct format, clearance of timer recovery if there is F = 1 only

BASIC STATE	TIMER RECOVERY							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
RNR command P = 1 V(A) ≤ N(R) ≤ V(S)	TX RR F = 1 V(A) = N(R) 8.4	TX RR F = 1 V(A) = N(R) 8.5	TX RNR F = 1 V(A) = N(R) 8.6	TX RNR F = 1 V(A) = N(R) 8.7	TX RR F = 1 V(A) = N(R)		TX RNR F = 1 V(A) = N(R)	
RNR command P = 0 V(A) ≤ N(R) ≤ V(S)	V(A) = N(R) 8.4	V(A) = N(R) 8.5	V(A) = N(R) 8.6	V(A) = N(R) 8.7	V(A) = N(R)			
RNR response F = 0 V(A) ≤ N(R) ≤ V(S)								
RNR response F = 1 V(A) ≤ N(R) ≤ V(S)	V(S) = N(R) RESTART T200 V(A) = N(R) 7.4	V(S) = N(R) RESTART T200 V(A) = N(R) 7.5	V(S) = N(R) RESTART T200 V(A) = N(R) 7.6	V(S) = N(R) RESTART T200 V(A) = N(R) 7.7	V(S) = N(R) RESTART T200 V(A) = N(R) 7.4	V(S) = N(R) RESTART T200 V(A) = N(R) 7.5	V(S) = N(R) RESTART T200 V(A) = N(R) 7.6	V(S) = N(R) RESTART T200 V(A) = N(R) 7.7
RNR command P = 1 N(R) error	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	
RNR command P = 0 N(R) error	MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1							
RNR response F = 0 N(R) error								
RNR response F = 1 N(R) error								

Table D.3/Q.921 (sheet 7 of 10) – State transition table: Receiving I command frame with correct format acknowledging all outstanding I frames or containing an N(R) which satisfies $V(A) < N(R) < V(S)$; no clearance of timer recovery

BASIC STATE	TIMER RECOVERY							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
I command P = 1 N(S) = V(R) N(R) = V(S)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	
		8.0				8.4		
I command P = 0 N(S) = V(R) N(R) = V(S)	V(R) = V(R) + 1 DL-DATA ind TX ACK V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX ACK V(A) = N(R)	"DISCARD" V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)	
		8.0				8.4		
I command P = 1 N(S) ≠ V(R) N(R) = V(S)	"DISCARD" TX REJ F = 1 V(A) = N(R)	"DISCARD" TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)		"DISCARD" TX REJ F = 1 V(A) = N(R)	"DISCARD" TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	
	8.1				8.5			
I command P = 0 N(S) ≠ V(R) N(R) = V(S)	"DISCARD" TX REJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)			"DISCARD" TX REJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)		
	8.1				8.5			
I command P = 1 N(S) = V(R) V(A) < N(R) < V(S)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	
		8.0				8.4		

Table D.3/Q.921 (sheet 7 of 10) – State transition table: Receiving I command frame with correct format acknowledging all outstanding I frames or containing an N(R) which satisfies $V(A) < N(R) < V(S)$; no clearance of timer recovery (*con cluded*)

BASIC STATE	TIMER RECOVERY							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
I command P = 0 N(S) = V(R) V(A) < N(R) < V(S)	V(R) = V(R) + 1 DL-DATA ind TX ACK V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX ACK V(A) = N(R) 8.0	"DISCARD" V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R) 8.4	"DISCARD" V(A) = N(R)	
I command P = 1 N(S) ≠ V(R) V(A) < N(R) < V(S)	"DISCARD" TX REJ F = 1 V(A) = N(R) 8.1	"DISCARD" TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)		"DISCARD" TX REJ F = 1 V(A) = N(R) 8.5	"DISCARD" TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	
I command P = 0 N(S) ≠ V(R) V(A) < N(R) < V(S)	"DISCARD" TX REJ F = 0 V(A) = N(R) 8.1	"DISCARD" V(A) = N(R)			"DISCARD" TX REJ F = 0 V(A) = N(R) 8.5	"DISCARD" V(A) = N(R)		

Table D.3/Q.921 (sheet 8 of 10) – State transition table: Receiving I command frame with correct format containing an N(R) which satisfies $V(A) = N(R) < V(S)$, or an N(R) error

BASIC STATE	TIMER RECOVERY							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
I command P = 1 N(S) = V(R) V(A) = N(R) < V(S)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 8.0	"DISCARD" TX RNR F = 1		V(R) = V(R) + 1 DL-DATA ind TX RR F = 1	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 8.4	"DISCARD" TX RNR F = 1	
I command P = 0 N(S) = V(R) V(A) = N(R) < V(S)	V(R) = V(R) + 1 DL-DATA ind TX ACK	V(R) = V(R) + 1 DL-DATA ind TX ACK 8.0	"DISCARD"		V(R) = V(R) + 1 DL-DATA ind TX RR F = 0	V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 8.4	"DISCARD"	
I command P = 1 N(S) ≠ V(R) V(A) = N(R) < V(S)	"DISCARD" TX REJ F = 1 8.1	"DISCARD" TX RR F = 1	"DISCARD" TX RNR F = 1		"DISCARD" TX REJ F = 1 8.5	"DISCARD" TX RR F = 1	"DISCARD" TX RNR F = 1	
I command P = 0 N(S) ≠ V(R) V(A) = N(R) < V(S)	"DISCARD" TX REJ F = 0 8.1	"DISCARD"			"DISCARD" TX REJ F = 0 8.5	"DISCARD"		
I command P = 1 N(S) = V(R) N(R) error	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	
I command P = 0 N(S) = V(R) N(R) error	V(R) = V(R) + 1 DL-DATA ind MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		"DISCARD" MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		V(R) = V(R) + 1 DL-DATA ind MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		"DISCARD" MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	

Table D.3/Q.921 (sheet 8 of 10) – State transition table: Receiving I command frame with correct format containing an N(R) which satisfies $V(A) = N(R) < V(S)$, or an N(R) error (concluded)

BASIC STATE	TIMER RECOVERY							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
I command N(S) ≠ V(R) N(R) error P = 1	"DISCARD" TX REJ F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	"DISCARD" TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		"DISCARD" TX REJ F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	"DISCARD" TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	
I command N(S) ≠ V(R) N(R) error P = 0	"DISCARD" TX REJ F = 0 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1			"DISCARD" TX REJ F = 0 MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SABME P = 1 RESTART T200 5.1		

Table D.3/Q.921 (sheet 9 of 10) – State transition table: Internal events (expiry of timers, receiver busy condition); initiation of a re-establishment procedure if the value of the retransmission count variable is equal to N200

BASIC STATE	TIMER RECOVERY							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
T200 TIME-OUT RC < N200 V(A) < V(S)	Either V(S) = V(S) – 1 TX I P = 1 V(S) = V(S) + 1 or TX RR P = 1 then RC = RC + 1 START T200		Either V(S) = V(S) – 1 TX I P = 1 V(S) = V(S) + 1 or TX RNR P = 1 then RC = RC + 1 START T200		TX RR P = 1 RC = RC + 1 START T200		TX RNR P = 1 RC = RC + 1 START T200	
T200 TIME-OUT RC < N200 V(A) = V(S)	TX RR P = 1 RC = RC + 1 START T200		TX RNR P = 1 RC = RC + 1 START T200					
T200 TIME-OUT RC = N200	MDL-ERR ind(I) RC = 0 TX SABME P = 1 START T200 5.1							
T203 TIME-OUT	/	/	/	/	/	/	/	/
SET OWN RECEIVER BUSY (Note)	TX RNR F = 0 8.2	TX RNR F = 0 8.3	–	–	TX RNR F = 0 8.6	TX RNR F = 0 8.7	–	–
CLEAR OWN RECEIVER BUSY (Note)	–	–	TX RR F = 0 8.0	TX RR F = 0 8.1	–	–	TX RR F = 0 8.4	TX RR F = 0 8.5

NOTE – These signals are generated outside the procedures specified in this state transition table, and may be generated by the connection management entity.

Table D.3/Q.921 (sheet 10 of 10) – State transition table: Receiving frame with incorrect format or frame not implemented

BASIC STATE	TIMER RECOVERY							
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	REJ recovery	Own rec busy	REJ and own rec busy	Normal	REJ recovery	Own rec busy	REJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
SABME incorrect length	MDL-ERR ind(N) RC = 0 TX SABME P = 1 RESTART T200 5.1							
DISC incorrect length								
UA incorrect length								
DM incorrect length								
FRMR incorrect length								
Supervisory frame RR, REJ, RNR incorrect length								
N201 error	MDL-ERR ind(O) RC = 0 TX SABME P = 1 RESTART T200 5.1							
Undefined command and response frames	MDL-ERR ind(L) RC = 0 TX SABME P = 1 RESTART T200 5.1							
I field not permitted	MDL-ERR ind(M) RC = 0 TX SABME P = 1 RESTART T200 5.1							

ANNEX E

Provision of Multi-Selective Reject Option

Subclauses E.1 to E.5 have a one-to-one correspondence with clauses 1 through 5 of this Recommendation. In this Annex, special rules apply in order to indicate **deviations** from text according to Q.921 clauses 1 through clauses 5, as follows:

- i) Added text in relation to Q.921 clauses 1 through clauses 5 is shown as double underlined;
- ii) Deleted text in relation to Q.921 clauses 1 through clauses 5 is shown as ~~strikethrough~~;
- iii) Replaced clauses are shown in *italics*;
- iv) Replaced clauses include those which describe procedures for REJ and SABME, and those which have no change or little change so that alignment of the clause numbering may be maintained.

E.1 General

This Annex defines the SET MODE (SM) command, the Selective Reject (SREJ) response, and the procedures used to support multi-selective reject option in LAPD. The multi-selective reject option of LAPD reduces the sensitivity of the throughput of the Data-link layer to degradations in the Bit Error Ratio of the underlying transmission media.

The procedures defined in this Annex are recommended for use on applications where there is a significant probability of having more than one unacknowledged I-frame. In the absence of such significant probability, the procedures of the main body of Recommendation Q.921 shall apply.

The procedures defined in this Annex are optional, and their use requires bilateral agreement between user and network. In the absence of provision of the procedures defined in this Annex, the procedures of the main body of ITU-T Recommendation Q.921 shall apply.

E.2 Frame structure for peer-to-peer communication

See clause 2.

E.3 Elements of procedures and formats of fields for data link layer peer-to-peer communication

E.3.1 General

See 3.1.

E.3.2 Address field format

See 3.2.

E.3.3 Address field variables

See 3.3.

E.3.4 Control field formats

See 3.4.

E.3.5 Control field parameters and associated state variables

See 3.5.

E.3.5.1 Poll/Final (P/F) bit

The use of P/F bit is described in E.5.

E.3.5.2 Multiple frame operation – variables and sequence numbers

E.3.5.2.1 Modulus

See 3.5.2.1.

E.3.5.2.2 Send state variable V(S)

See 3.5.2.2.

E.3.5.2.3 Acknowledge state variable V(A)

See 3.5.2.3.

E.3.5.2.4 Send sequence number N(S)

See 3.5.2.4.

E.3.5.2.5 Receive state variable V(R)

See 3.5.2.5.

E.3.5.2.6 Receive sequence number N(R)

See 3.5.2.6.

E.3.5.2.7 Poll sequence number variable V(P)

The data link layer entity shall maintain a poll sequence number variable V(P). V(P) shall be incremented after a frame with the P bit set to 1 is sent. V(P) together with V_S(P) is used to prevent duplicate retransmission.

E.3.5.2.8 Saved poll sequence number variable V_S(P)

Each I frame shall have an associated V_S(P). Whenever an I frame is transmitted or retransmitted, the current value of V(P) shall be saved in V_S(P). During retransmission, the value of V_S(P) of an I frame shall be compared with the current value of V(P). If the value of V_S(P) is equal to the current value of V(P), the I frame was transmitted after the last frame with the P bit set to 1 was sent, and shall not be retransmitted again.

E.3.5.3 Unacknowledged operation – variables and parameters

See 3.5.3

E.3.6 Frame types

E.3.6.1 Commands and responses

The following commands and responses are used by either the user or the network data link layer entities and are represented in Table 5/Q.924 E.5/Q.921. Each data link connection shall support the full set of commands and responses for each application implemented. The frame types associated with each of the two applications are identified in Table 5/Q.924 E.5/Q.921.

Frame types associated with an application not implemented shall be discarded and no action shall be taken as a result of that frame.

For purposes of the LAPD procedures in each application, those encodings not identified in Table 5/Q.921 E.5/Q.921 are identified as undefined command and response control fields. The actions to be taken are specified in E.5.8.5.

The commands and responses in Table 5/Q.921 E.5/Q.921 and defined in E.3.6.2 to E.3.6.12.

Table E.5/Q.921 – Commands and responses (modulo 128)

Application	Format	Commands	Responses	Encoding								
				8	7	6	5	4	3	2	1	
Unacknowledged and Multiple Frame Acknowledged Information Transfer	Information Transfer	I		N(S)							0	
				N(R)							P	
	Supervisory	RR	RR	0	0	0	0	0	0	0	1	
				N(R)							P/F	
		RNR	RNR	0	0	0	0	0	1	0	1	
				N(R)							P/F	
			SREJ	0	0	0	0	1	1	0	1	
		N(R)							P/F			
	Unnumbered	SM		1	1	0	P	0	0	1	1	
			DM	0	0	0	F	1	1	1	1	
		UI		0	0	0	P	0	0	1	1	
		DISC		0	1	0	P	0	0	1	1	
			UA		0	1	1	F	0	0	1	1
			FRMR		1	0	0	F	0	1	1	1
Connection Mgt.	XID	XID	1	0	1	P/F	1	1	1	1		
NOTE – Use of the XID frame other than for parameter negotiation procedures (see 5.4) is for further study.												

E.3.6.2 Information (I) command

See 3.6.2.

E.3.6.3 Set Mode (SM) command

The SM command is used to place the addressed user side or network side into multiple frame acknowledged operation and to invoke multi-selective reject error recovery procedure. The SM command is shown in Figure E.3.6.3-1/Q.921.

The SM command may contain an optional information field. The first octet of the optional information field is the format identifier and is encoded as "10000001".

The second octet of the optional information field is the Group Identifier (GI) and is encoded as "10000001" which is the value defined for Mode and Modulus Group.

The third octet of the optional information field is the Group Length (GL) and is set to 2 (binary value of "00000010").

The fourth octet of the optional information field is the Mode of Operation and is encoded as "00000001" which is the value defined for Asynchronous Balanced Mode.

The fifth octet of the information field is the Modulus of Operation and is encoded as "00000010" for Modulo 128 operation.

Octet	8	7	6	5	4	3	2	1	
1	Flag								
	0	1	1	1	1	1	1	0	
2	Address (high order octet)								
3	Address (low order octet)								
4	Control								Control = SM
	1	1	0	P	0	0	1	1	
5	Format identifier								
	1	0	0	0	0	0	0	1	
6	Group identifier								Group identifier = Mode and Modulus
	1	0	0	0	0	0	0	1	
7	Group length								Group length = 2
	0	0	0	0	0	0	1	0	
8	Mode of operation								Mode of operation = ABM
	0	0	0	0	0	0	0	1	
9	Modulus of operation								Modulus of operation = 128
	0	0	0	0	0	0	1	0	
10	FCS (first octet)								
11	FCS (second octet)								
12	Flag								
	0	1	1	1	1	1	1	0	

Figure E.3.6.3-1/Q.921 – Set Mode (SM) command

E.3.6.4 Disconnect (DISC) command

See 3.6.4.

E.3.6.5 Unnumbered information (UI) command

See 3.6.5.

E.3.6.6 Receive ready (RR) command/response

See 3.6.6.

E.3.6.7 Multi-selective reject (SREJ) response

The SREJ response is used to initiate a more efficient error recovery by selectively requesting the retransmission of one or more I frames following the detection of sequence errors.

E.3.6.7.1 SREJ response frame format

The SREJ frame format is shown in Figure E.3.6.7.1-1/Q.921.

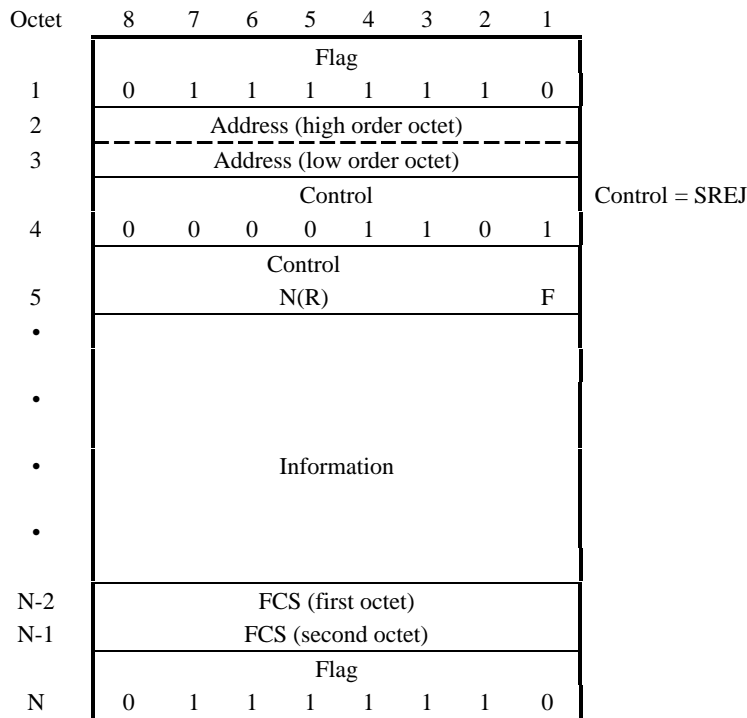


Figure E.3.6.7.1-1/Q.921 – Selective Reject (SREJ) response

The N(R) subfield of the control field of the SREJ response shall contain the sequence number of the oldest missing I frame, and the information field of the SREJ response shall contain the sequence numbers of the remaining missing I frames. If the list of sequence numbers is too large to fit in the information field of the SREJ response, then the list shall be truncated by including only the sequence numbers of the oldest missing I frames.

When the F bit of the SREJ response is set to "0", the N(R) subfield of the control field shall not be used for an acknowledgement function. When the F bit of the SREJ response is set to "1", the N(R) subfield of the control field shall be used for an acknowledgement function, i.e. the N(R) shall be used as an acknowledgement for all I frames that have been transmitted with an N(S) up to and including the received N(R)-1.

E.3.6.7.2 Encoding of the information field in the SREJ response frame

The identity of missing I frames shall be indicated by either:

- a) one octet for every I frame; or
- b) one octet for every stand-alone I frame plus a span list for every sequence of two or more contiguously numbered I frames.

The span list identifies the start and the end of a sequence of contiguously numbered I frames in need of retransmission.

The sequence number of the stand-alone I frames shall consist of a 7-bit N(R) value preceded by a "0" bit as shown in Figure E.3.6.7.2-1/Q.921. The sequence number of the start and end of a span list of I frames shall consist of a 7-bit N(R) value preceded by a "1" bit as shown in Figure E.3.6.7.2-2/Q.921.

Figure E.3.6.7.2-1/Q.921 shows a case where I frames 4, 6, 9 and 13 are in need of retransmission. The sequence number of I frame 4 is contained in the N(R) subfield of the SREJ control field. The

sequence number of I frames 6, 9 and 13 are encoded as stand-alone I frames and are indicated by having the bit that precedes the 7-bit N(R) value set to "0".

Figure E.3.6.7.2-2/Q.921 shows a case where I frames 4, 6, 9, 10, 11, 12 and 13 are in need of retransmission. The sequence number of I frame 4 is contained in the N(R) subfield of the SREJ control field. The sequence number of I frame 6 is encoded as a stand-alone I frame and is indicated by having the bit that precedes the 7-bit N(R) value set to "0". The sequence number of I frames 9, 10, 11, 12 and 13 are encoded as a span list of five contiguous I frames starting with I frame 9 and ending with I frame 13.

Octet	8	7	6	5	4	3	2	1	
	Flag								
1	0	1	1	1	1	1	1	0	
2	Address (high order octet)								
3	Address (low order octet)								
	Control								Control = SREJ
4	0	0	0	0	1	1	0	1	
	Control								N(R) = 4
5	0	0	0	0	1	0	0	F	
	Information								
6	0	0	0	0	0	1	1	0	
	Information								
7	0	0	0	0	1	0	0	1	
	Information								
8	0	0	0	0	1	1	0	1	
9	FCS (first octet)								
10	FCS (second octet)								
	Flag								
11	0	1	1	1	1	1	1	0	

Figure E.3.6.7.2-1/Q.921 – Stand-alone I frame encoding

Octet	8	7	6	5	4	3	2	1	
	Flag								
1	0	1	1	1	1	1	1	0	
2	Address (high order octet)								
3	Address (low order octet)								
	Control								Control = SREJ
4	0	0	0	0	1	1	0	1	
	Control								N(R) = 4
5	0	0	0	0	1	0	0	F	
	Information								
6	0	0	0	0	0	1	1	0	
	Information								
7	1	0	0	0	1	0	0	1	
	Information								
8	1	0	0	0	1	1	0	1	
9	FCS (first octet)								
10	FCS (second octet)								
	Flag								
11	0	1	1	1	1	1	1	0	

Figure E.3.6.7.2-2/Q.921 – Stand-alone I frame and span list encoding

E.3.6.8 Receive not ready (RNR) command/response

See 3.6.8.

E.3.6.9 Unnumbered acknowledgement (UA) response

The UA unnumbered response is used by a data layer entity to acknowledge the receipt and acceptance of the mode-setting commands (SM or DISC). An information field is optional in the UA response and may be present when the response acknowledges the receipt of the SM command.

The format of the information field of the UA response is the same as the format of the information field of the SM command as defined in E.3.6.3.

When the UA response has an information field containing the Mode and Modulus Group, at most one mode bit and one modulus bit can be set to 1.

The transmission of the UA response indicates the clearance of any busy condition that was reported by the earlier transmission of an RNR frame by that same data link layer entity.

E.3.6.10 Disconnected mode (DM) response

See 3.6.10.

E.3.6.11 Frame reject (FRMR) response

The FRMR unnumbered response may be received by a data link layer entity as a report of an error condition not recoverable by retransmission of the identical frame, i.e. at least one of the following error conditions resulting from the receipt of a valid frame:

- a) the receipt of a command or response control field that is undefined;
- b) the receipt of a supervisory or unnumbered frame with incorrect length;
- c) the receipt of an invalid N(R); or
- d) the receipt of an I frame with an information field which exceeds the maximum established length.

An undefined control field is any of the control field encodings that are not identified in Table ~~5/Q.924~~ E.5/Q.921.

A valid N(R) value is one that is in the range $V(A) \leq N(R) \leq V(S)$.

An information field which immediately follows the control field and consists of nine octets (modulo 128 operation) is returned with this response and provides the reason for the FRMR response. This information is given in Figure 6.

E.3.6.12 Exchange identification (XID) command/response

See 3.6.12.

E.4 Elements for layer-to-layer communication

See clause 4.

E.5 Definition of the peer-to-peer procedures of the data link layer

The procedures for use by the data link layer are specified in the following subclauses.

The elements of procedure (frame types) which apply are:

- a) for unacknowledged information transfer (see E.5.2):
UI-command;

b) for multiple frame acknowledged information transfer (E.5.5 to E.5.8):

~~SABMESM~~-command;

UA-response;

DM-response;

DISC-command;

RR-command/response;

RNR-command/response;

~~REJSREJ~~-command/response;

I-command;

FRMR-response (Note);

NOTE – An FRMR-response shall not be generated by a data link layer entity; however, on receipt of this frame actions according to E.5.8.6 shall be taken.

c) for connection management entity information transfer:

XID-command/response.

E.5.1 Procedure for the use of the P/F bit

E.5.1.1 Unacknowledged information transfer

See 5.1.1.

E.5.1.2 Acknowledged multiple frame information transfer

A data link layer entity receiving an ~~SABMESM~~, DISC, RR, RNR, ~~REJSREJ~~ or I frame, with the P bit set to 1, shall set the F bit to 1 in the next response frame it transmits, as defined in Table ~~7/Q.921~~ E.7/Q.921.

Table E.7/Q.921 – Immediate response operation of P/F bit

Command received with P bit = 1	Response transmitted with F bit = 1
SABMESM , DISC	UA, DM
I, RR, RNR, REJ	RR, RNR, REJSREJ (Note)
NOTE – A LAPB data link layer entity may transmit an FRMR or DM response with the F bit set to 1 in response to an I frame or supervisory command with the P bit set to 1.	

E.5.2 Procedures for unacknowledged information transfer

See 5.2.

E.5.3 Terminal Endpoint Identifier (TEI) management procedures

See 5.3.

E.5.4 Initialization of data link layer parameters

See 5.4.

E.5.5 Procedures for establishment and release of multiple frame operation

E.5.5.1 Establishment of multiple frame operation

The procedures for establishment of multiple frame operation defined in 5.5.1 apply with the following changes:

- a) *all references to "SABME" are replaced by "SM";*
- b) *all references to "set V(S), V(R) and V(A) to 0" are replaced by "set V(S), V(R), V(A) and V(P) to 0";*
- c) *all references to "5.7" are replaced by "E.5.7".*

E.5.5.2 Information transfer

Having either transmitted the UA response to a received ~~SABMESM~~ command or received the UA response to a transmitted ~~SABMESM~~ command, I frames and supervisory frames shall be transmitted and received according to the procedures described in E.5.6.

If an ~~SABMESM~~ command is received while in the multiple-frame-established state, the data link layer entity shall conform to the re-establishment procedure described in E.5.7.

On receipt of a UI command, the procedures defined in E.5.2 shall be followed.

E.5.5.3 Termination of multiple frame operation

Termination of multiple frame operation defined in 5.5.3 applies with all references to "SABME" replaced by "SM".

E.5.5.4 TEI-assigned state

The procedures used while in the TEI-assigned state as defined in 5.5.4 apply with all references to "SABME" replaced by "SM".

E.5.5.5 Collision of unnumbered commands and responses

The procedures used to handle collision of unnumbered commands and responses as defined in 5.5.5 apply with all references to "SABME" replaced by "SM".

E.5.6 Procedures for information transfer in multiple frame operation

The procedures which apply to the transmission of I frames are defined below.

NOTE – The term "transmission of an I frame" refers to the delivery of an I frame by the data link layer to the physical layer.

E.5.6.1 Transmitting I frames

Information received by the data link layer entity from layer 3 by means of a DL-DATA-REQUEST primitive shall be transmitted in an I frame. The control field parameters N(S) and N(R) shall be assigned the values of V(S) and V(R), respectively. V(S) shall be incremented by 1 at the end of the initial transmission of the I frame.

Retransmission of I frames is defined in E.5.6.4 and E.5.6.5.

Whenever an I frame is transmitted or retransmitted, the current value of V(P) shall be saved in $V_S(P)$ associated with the I frame.

If time T200 is not running at the time of transmission of an I frame, it shall be started. If timer T200 expires, the procedures defined in E.5.6.7 shall be followed.

If $V(S)$ is equal to $V(A)$ plus k (where k is the maximum number of outstanding I frames – see E.5.9.5), the data link layer entity shall not transmit any new I frames, but may retransmit an I frame as a result of the error recovery procedures as described in E.5.6.4 and E.5.6.7.

When the network side or user side is in the own receiver busy condition, it may still transmit I frames, provided that a peer receiver busy condition does not exist.

NOTE – Any DL-DATA-REQUEST primitives received whilst in the timer recovery condition shall be queued.

E.5.6.2 Receiving I frames

Independent of a timer recovery condition, when a data link layer entity is not in an own receiver busy condition and receives a valid I frame, the data link layer entity shall act as follows:

- if the data link layer entity is not in an $N(S)$ sequence error exception condition and the received I frame is an in-sequence I frame [i.e. $N(S)$ is equal to $V(R)$] then:
 - pass the information field of this frame to layer 3 using the DL-DATA-INDICATION primitive;
 - increment by 1 its $V(R)$ ~~and act as indicated below;~~
- if the data link layer entity is in an $N(S)$ sequence error exception condition and the received I frame is an in-sequence I frame then:
 - While there is in-sequence I frame in the receive queue:
 - pass the information field of this frame to layer 3 using the DL-DATA-INDICATION primitive;
 - increment by 1 its $V(R)$;
 - If there is no out-of-sequence I frame in receive queue then clear the $N(S)$ sequence error exception condition;
- if the received I frame is an out-of-sequence I frame [i.e. $N(S)$ is not equal to $V(R)$ and I frame numbered $N(S)-1$ has not been received] then:
 - hold the received I frame in the receive queue for later delivery [the data link layer entity shall deliver the I frame to the layer 3 only when it correctly receives all missing I frames numbered below $N(S)$];
 - create a list of sequence numbers of missing I frames ending at $N(S)-1$;
 - if the data link layer entity is not in an $N(S)$ sequence error exception condition then set the $N(S)$ sequence error exception condition.

E.5.6.2.1 P bit set to 1

If the P bit of the received I frame was set to 1, the data link layer entity shall respond to its peer in one of the following ways:

- if the data link layer entity receiving the I frame is still not in an own receiver busy condition, but is in an $N(S)$ sequence error exception condition, it shall send an SREJ frame with:
 - the F bit set to 1;
 - the $N(R)$ set to the current value of $V(R)$;
 - the information field set to the sequence number(s) of the remaining missing I frames as defined in E.3.6.7;

- if the data link layer entity receiving the I frame is still not in an own receiver busy condition and is not in an N(S) sequence error exception condition, it shall send an RR response with the F bit set to 1;
- if the data link layer entity receiving the I frame enters the own receiver busy condition upon receipt of the I frame, it shall send an RNR response with the F bit set to 1.

E.5.6.2.2 P bit set to 0

If the P bit of the received I frame was set to 0 and:

- a) if the data link layer entity is still not in an own receiver busy condition, it shall respond to its peer in one of the following ways:
 - If the data link layer entity is in an N(S) sequence error exception condition, it shall transmit an SREJ frame with:
 - the F bit set to 0;
 - the N(R) set to the sequence number of the oldest missing I frame in the list;
 - the information field set to the sequence numbers of the remaining missing I frames in the list as defined in E.3.6.7;
 - if the data link layer entity is not in an N(S) sequence error exception condition and, if no I frame is available for transmission or if an I frame is available for transmission but a peer receiver busy condition exists, the data link layer entity shall transmit an RR response with the F bit set to 0; or
 - if the data link layer entity is not in an N(S) sequence error exception condition and, if an I frame is available for transmission and no peer receiver busy condition exists, the data link layer entity shall transmit the I frame with the value of N(R) set to the current value of V(R) as defined in E.5.6.1; or
- b) if, on receipt of this I frame, the data link layer entity is now in an own receiver busy condition, it shall transmit an RNR response with the F bit set to 0.

When the data link layer entity is in an own receiver busy condition, it shall process any received I frame according to E.5.6.6.

E.5.6.3 Sending and receiving acknowledgements

E.5.6.3.1 Sending acknowledgements

Whenever a data link layer entity transmits an I frame, ~~or an RR, RNR supervisory frame or an SREJ frame with the F bit set to 1~~, N(R) shall be set equal to V(R).

E.5.6.3.2 Receiving acknowledgements

On receipt of a valid I frame ~~or supervisory frame (RR, RNR, or REJ)~~ RR, RNR frame or SREJ frame with the F bit set to 1, even in the own receiver busy, or timer recovery conditions, the data link layer entity shall treat the N(R) contained in this frame as an acknowledgement for all the I frames it has transmitted with an N(S) up to and including the received N(R)–1. V(A) shall be set to N(R). When not in the timer recovery condition the data link layer entity shall reset the timer T200 ~~either on receipt of a valid I frame or supervisory frame with the N(R) higher than V(A) (actually acknowledging some I frames) or an REJ frame with an N(R) equal to V(A).~~

- a valid I frame or RR, RNR frame with the N(R) higher than V(A) (actually acknowledging some I frames); or
- a valid SREJ frame with the F bit set to 1 and with the N(R) higher than V(A) (actually acknowledging some I frames); or

– a valid SREJ frame with the F bit set to 1 and with an N(R) equal to V(A).

NOTE 1 – If a supervisory frame or an I frame with the P bit set to 1 has been transmitted and not acknowledged by a supervisory frame response with the F bit set to 1, timer T200 shall not be reset.

NOTE 2 – Upon receipt of a valid I frame, timer T200 shall not be reset if the data link layer entity is in the peer receiver busy condition.

If timer T200 has been reset by the receipt of an I, RR, or RNR frame, and if there are outstanding I frames still unacknowledged, the data link layer entity shall restart timer T200. If timer T200 then expires, the data link layer shall follow the recovery procedure as defined in E.5.6.7 with respect to the unacknowledged I frames.

If timer T200 has been reset by the receipt of an SREJ frame with the F bit set to 1, the data link layer entity shall follow the retransmission procedures in E.5.6.4.

Upon receipt of a valid RR command frame with the P bit set to 1, the data link layer entity shall respond to its peer in one of the following ways at the earliest opportunity:

- if the data link layer entity is not in an own receiver busy condition, but is in an N(S) sequence error exception condition, the data link layer entity shall send an SREJ response frame with:
 - the F bit set to 1;
 - the N(R) set to V(R);
 - the information field set to the sequence numbers of the remaining missing I frames as defined in E.3.6.7;
- if the data link layer entity is not in an own receiver busy condition and is not in an N(S) sequence error exception condition, it shall respond with an RR response with the F bit set to 1;
- if the data link layer entity is in an own receiver busy condition, it shall respond with an RNR response with the F bit set to 1.

E.5.6.4 Receiving SREJ frames

- a) *On receipt of a valid SREJ response with the F bit set to 0, the data link layer entity shall act as follows:*
- *clear an existing peer receiver busy condition;*
 - *retransmit all requested I frames, taking into account the items 1) to 3) below.*
 - *Set the N(S) of each retransmitted I frame to the appropriate sequence number contained in the SREJ response and set the N(R) of each retransmitted I frame to the current value of V(R);*
 - *save the current value of V(P) in $VS_s(P)$ associated with each retransmitted I frame;*
 - *if it is not in the timer recovery condition:*
 - *send a poll either by sending an RR command (or RNR command if the data link layer entity is in own receiver busy condition) with the P bit set to 1 or by setting the P bit to 1 in the last retransmitted I frame;*
 - *increment V(P) by 1;*
 - *restart timer T200;*
 - *enter the timer recovery condition.*

- b) *On receipt of a valid SREJ response with the F bit set to 1, the data link layer entity shall act as follows:*
- *clear an existing peer receiver busy condition;*
 - *set its V(A) to the value of the N(R) contained in the SREJ response control field;*
 - *stop timer T200;*
 - *start timer T203, if implemented;*
 - *if it is not in the timer recovery condition then notify a protocol violation to the connection management entity by means of the MDL-ERROR indication primitive;*
 - *if it is in the timer recovery condition then enter the multiple-frame-established state;*
 - *retransmit all requested I frames whose $V_s(P)$ is not equal to the current value of V(P), taking into account the items 1) to 3) below;*
 - *if any I frames are retransmitted then:*
 - *set the N(S) of each retransmitted I frame to the appropriate sequence number contained in the SREJ response and set the N(R) of each retransmitted I frame to the current value of V(R);*
 - *save the current value of V(P) in $V_s(P)$ associated with each retransmitted I frames;*
 - *send a poll either by sending an RR command (or RNR command if the data link layer entity is in the own receiver busy condition) with the P bit set to 1 or by setting the P bit to 1 in the last retransmitted I frame;*
 - *increment V(P) by 1;*
 - *restart timer T200;*
 - *enter the timer recovery condition.*

Transmission of I frames shall take account of the following:

- 1) if the data link layer entity is transmitting a supervisory frame when it receives the REJSREJ response, it shall complete that transmission before commencing transmission of the requested I frame;
- 2) if the data link layer entity is transmitting an SABMESM command, a DISC command, a UA response or a DM response when it receives the REJSREJ response, it shall ignore the request for retransmission; and
- 3) if the data link layer entity is not transmitting a frame when the REJSREJ response is received, it shall immediately commence transmission of the requested frames.

E.5.6.5 Receiving RNR frames

After receiving a valid RNR command or response, if the data link layer entity is not engaged in a mode-setting operation, it shall set a peer receiver busy condition and then:

- ~~if it was an RNR command with the P bit set to 1, it shall respond with an RR response with the F bit set to 1 if the data link layer entity is not in an own receiver busy condition, and shall respond with an RNR response with the F bit set to 1 if the data link layer entity is an own receiver busy condition; and the data link layer entity shall respond to its peer in one of the following ways:~~
 - if the data link layer entity is not in an own receiver busy condition, but is in an N(S) sequence error exception condition, the data link layer entity shall send an SREJ frame with:
 - the F bit set to 1;

- the N(R) set to V(R);
- the information field set to the sequence numbers of the remaining missing I frames as defined in E.3.6.7;
- if the data link layer entity is not in an own receiver busy condition and is not in an N(S) sequence error exception condition, it shall respond with an RR response with the F bit set to 1;
- if the data link layer entity is in an own receiver busy condition, it shall respond with an RNR response with the F bit set to 1;
- if it was an RNR response with the F bit set to 1, an existing timer recovery condition shall be cleared and the N(R) contained in this RNR response shall be used to update ~~V(S)~~ V(A).

The data link layer entity shall take note of the peer receiver busy condition and not transmit any I frames to the peer which has indicated the busy condition.

NOTE 1 – The N(R) in an RR or RNR command frame (irrespective of the setting of the P bit) will not be used to update the V(S).

The data link layer entity shall then:

- treat the N(R) contained in the received RNR frame as an acknowledgement for all the I frames that have been (re)transmitted with an N(S) up to and including N(R)–1, and set its V(A) to the value of the N(R) contained in the RNR frame; and
- restart timer T200 unless a supervisory response frame with the F bit set to 1 is still expected.

If timer T200 expires, the data link layer entity shall:

- if it is not yet in a timer recovery condition, enter the timer recovery condition and reset the retransmission count variable; or
- if it is already in a timer recovery condition, continue as indicated below.

The data link layer entity shall then:

- a) if the value of the retransmission count variable is less than N200:
 - transmit an ~~appropriate supervisory~~ RR command (see ~~NOTE 2~~) with a P bit set to 1;
 - restart timer T200;
 - add one to its retransmission count variable; and
- b) if the value of the retransmission count variable is equal to N200, initiate a re-establishment procedure as defined in E.5.7, and indicate this by means of the MDL-ERROR indication primitive to the connection management entity.

The data link layer entity receiving the ~~supervisory~~ RR command frame with the P bit set to 1 shall respond at the earliest opportunity, with a supervisory response frame (see ~~NOTE 2~~ E.5.6.3.2) with the F bit set to 1 to indicate whether or not its own receiver busy condition still exists.

Upon receipt of the supervisory response with the F bit set to 1, the data link layer entity shall reset timer T200, and:

- ~~if the response is an RR or REJ response, the peer receiver busy condition is cleared and the data link layer entity may transmit new I frames or retransmit I frames as defined in §§ 5.6.1 or 5.6.4, respectively; or the data link layer entity receiving the response shall proceed according to E.5.6.7;~~
- if the response is an SREJ response, the data link layer entity receiving the response shall proceed according to E.5.6.4 b);

- if the response is an RNR response, the data link layer entity receiving the response shall proceed according to this E.5.6.5, first paragraph.

If a supervisory command (RR, or RNR, ~~or REJ~~) with the P bit set to 0 or 1, or a supervisory response frame (RR, RNR, or SREJ) with the F bit set to 0 is received during the enquiry process, the data link layer entity shall:

- ~~— if the supervisory frame is an RR or REJ command frame or an RR or REJ response frame with the F bit set to 0, clear the peer receiver busy condition and if the supervisory frame received was a command with the P bit set to 1, transmit the appropriate supervisory response frame (see NOTE 2) with the F bit set to 1. However, the transmission or retransmission of I frames shall not be undertaken until the appropriate supervisory response frame with the F bit set to 1 is received or until expiry of timer T200; or~~
- ~~— if the supervisory frame is an RNR command frame or an RNR response frame with the F bit set to 0, retain the peer receiver busy condition and if the supervisory frame received was an RNR command with the P bit set to 1, transmit the appropriate supervisory response frame (see NOTE 2) with the F bit set to 1;~~
- if the supervisory frame is an SREJ response frame with the F bit set to 0, proceed according to E.5.6.4 a);
- if the supervisory frame is an RR response frame with the F bit set to 0, clear the peer receiver busy condition. New I frames shall be transmitted as they are available according to E.5.6.1;
- if the supervisory frame is an RR command with the P bit set to 1, clear the peer receiver busy condition and transmit the appropriate supervisory response frame with the F bit set to 1 (see E.5.6.3.2);
- if the supervisory frame is an RNR command frame or an RNR response frame with the F bit set to 0, retain the peer receiver busy condition and if the supervisory was an RNR command with P bit set to 1, proceed according to this E.5.6.5, first paragraph.

Upon receipt of an SABMESM command, the data link layer entity shall clear the peer receiver busy condition.

NOTE 2 – NOTE 2 is deleted.

E.5.6.6 Data link layer own receiver busy condition

When the data link layer entity enters an own receiver busy condition, it shall transmit an RNR frame at the earliest opportunity.

The RNR frame may be either:

- an RNR response with the F bit set to 0; or
- if this condition is entered on receiving a command frame with the P bit set to 1, an RNR response with the F bit set to 1; or
- if this condition is entered on expiry of timer T200, and RNR command with the P bit set to 1.

All received I frames with the P bit set to 0 shall be discarded, after updating V(A). An SREJ response with the F bit set to 0 shall not be transmitted until the own receiver busy condition is clear.

All supervisory frames with the P/F bit set to 0 shall be processed.

All received I frames with the P bit set to 1 shall be discarded, after updating V(A). However, an RNR response frame with the F bit set to 1 shall be transmitted.

All received supervisory frames with the P bit set to 1 shall be processed including updating V(A). An RNR response with the F bit set to 1 shall be transmitted. An SREJ response with the F bit set to 1 shall not be transmitted until the own receiver busy condition is clear.

To indicate to the peer data link layer entity the clearance of the own receiver busy condition, the data link layer entity shall transmit an RR frame or, if a previously detected N(S) sequence error has not yet been reported, an ~~REJSREJ~~ response frame with the N(R) set to the current value of V(R) and with the information field set to the sequence numbers of the remaining missing I frames as defined in E.3.6.7.

The transmission of an ~~SABMESM~~ command or a UA response (in reply to an ~~SABMESM~~ command) also indicates to the peer data link layer entity the clearance of the own receiver busy condition.

E.5.6.7 Waiting acknowledgement

The data link layer entity shall maintain an internal retransmission count variable.

If timer T200 expires, the data link layer entity shall:

- if it is not yet in a timer recovery condition, enter the timer recovery condition and reset the retransmission count variable; or
- if it is already in a timer recovery condition, continue as indicated below.

The data link layer entity shall then:

- a) if the value of the retransmission count variable is less than N200:
 - add one to its retransmission count variable; and
 - restart timer T200; and ~~either~~
 - transmit an ~~appropriate supervisory~~ RR command ~~(see NOTE 2 in § 5.6.5)~~ with a P bit set to 1; or
 - ~~retransmit the last transmitted I frame [V(S) – 1] with the P bit set to 1; or~~
- b) if the value of the retransmission count variable is equal to N200, initiate a re-establishment procedure as defined in E.5.7, and indicate this by means of the MDL-ERROR indication primitive to the connection management entity.

The following paragraphs ~~applies~~ apply only for a data link layer which is in the timer recovery condition since the case of receiving acknowledgement in the multiframe established state is described in E.5.6.3.2.

The timer recovery condition is cleared only if the data link layer entity receives a valid supervisory frame response with the F bit set to 1. ~~If the N(R) of this received supervisory frame is within the range from its current V(A) to its current V(S) inclusive, it shall set its V(S) to the received N(R).~~ Timer T200 shall be reset if the received supervisory frame response is an RR or ~~REJSREJ~~ response with the F bit set to 1. ~~The data link layer entity shall resume then with I frame transmission or retransmission, as appropriate.~~ Timer T200 shall be reset and restarted if the received supervisory response is an RNR response with the F bit set to 1, to proceed with the enquiry process according to E.5.6.5.

If the response is an RR response, the data link layer entity receiving the response shall act as follows:

- clear an existing peer receiver busy condition;
- set its V(A) to the value of the N(R) contained in the RR response control field;
- stop timer T200;

- = start timer T203, if implemented;
- = enter the multiple-frame-established state;
- = if there are no outstanding I frames still unacknowledged, the data link layer entity may transmit new I frames as defined in E.5.6.1;
- = if there are outstanding I frames still unacknowledged, the data link layer entity shall retransmit all unacknowledged I frames whose $V_S(P)$ is not equal to the current value of $V(P)$, taking into account the items 1) to 3) of E.5.6.4;
- = if any I frames are retransmitted then:
 - set the $N(S)$ of each retransmitted I frame to the sequence number of an appropriate unacknowledged I frame and set the $N(R)$ of each retransmitted I frame to the current value of $V(R)$;
 - save the current value of $V(P)$ in $V_S(P)$ associated with each retransmitted I frames;
 - send a poll either by sending an RR command (or RNR command if the data link layer entity is in the own receiver busy condition) with the P bit set to 1 or by setting the P bit to 1 in the last retransmitted I frame;
 - increment $V(P)$ by 1;
 - restart timer T200;
 - enter the timer recovery condition.

If the response is an SREJ response, the data link layer entity receiving the response shall proceed according to E.5.6.4 b).

If the response is an RNR response, the data link layer entity receiving the response shall proceed according to this E.5.6.5, first paragraph.

While in the timer recovery condition, the data link layer entity shall process the SREJ response with the F bit set to 0 according to E.5.6.4 a). While in the timer recovery condition, if no peer receiver busy condition exists, the data link layer entity shall transmit new I frames as these new I frames are available.

E.5.7 Re-establishment of multiple frame operation

The criteria for re-establishing the multiple frame mode of operation defined in 5.7 apply with all references to "SABME" replaced by "SM".

E.5.8 Exception condition reporting and recovery

See 5.8.

E.5.8.1 N(S) sequence error

An N(S) sequence error exception condition occurs in the receiver when a valid I frame is received which contains an N(S) value which is not equal to the V(R) at the receiver. The information field of all frames whose N(S) does not equal V(R) shall be discarded held in the receive queue for later delivery. The I frame shall be delivered to the upper layer only when all I frames numbered below N(S) are correctly received.

The receiver shall not acknowledge [nor increment its V(R)] the I frame causing the sequence error, nor any I frames which may follow, until an I frame with the correct N(S) is received.

A data link layer entity which receives one or more I frames having sequence errors but otherwise error-free, or subsequent supervisory frames (RR, RNR, and ~~REJ~~ SREJ with the P bit set to 1), shall use the control field information contained in the N(R) field and the P or F bit to perform data link

control functions; for example, to receive acknowledgement of previously transmitted I frames and to cause the data link layer entity to respond if the P bit is set to 1. Therefore, the retransmitted I frame may contain an N(R) field value and P bit that are updated from, and therefore different from, the ones contained in the originally transmitted I frame.

The ~~REJSREJ~~ response frame is used by a receiving data link layer entity to initiate an exception condition recovery (retransmission) following the detection of an N(S) sequence error. The procedures for using SREJ are defined in E.5.6.

~~Only one REJ exception condition for a given direction of information transfer shall be established at a time.~~

A data link layer entity receiving an ~~REJSREJ~~ command or response shall initiate sequential transmission (retransmission) of I frames starting with the I frame indicated by the N(R) contained in the control field of the REJSREJ frame, followed by the I frames whose sequence numbers are indicated in the information field of the SREJ frame.

An ~~REJSREJ~~ exception condition is cleared when ~~the~~ all requested I frames is are received or when an SABME or DISC command is received.

~~An optional procedure for the retransmission of an REJ response is described in Appendix I.~~

E.5.8.2 N(R) sequence error

N(R) sequence error defined in 5.8.2 applies with all references to "SABME" replaced by "SM".

E.5.8.3 Timer recovery condition

If a data link layer entity, due to a transmission error, does not receive a single I frame or the last I frame(s) in a sequence of I frames, it will not detect an out-of-sequence exception condition and therefore will not transmit an ~~REJSREJ~~ frame.

The data link layer entity which transmitted the unacknowledged I frame(s) shall, on the expiry of timer T200, take appropriate recovery action as defined in E.5.6.7 to determine at which I frame retransmission must begin.

E.5.8.4 Invalid frame condition

See 5.8.4.

E.5.8.5 Frame rejection condition

A frame rejection condition results from one of the following conditions:

- the receipt of an undefined frame (see E.3.6.1, third paragraph);
- the receipt of a supervisory or unnumbered frame with incorrect length;
- the receipt of an invalid N(R); or
- the receipt of a frame with an information field which exceeds the maximum established length.

Upon occurrence of a frame rejection condition whilst in the multiple frame operation, the data link layer entity shall:

- issue an MDL-ERROR indication primitive; and
- initiate re-establishment (see E.5.7.2).

Upon occurrence of a frame rejection condition during establishment or release from multiple frame operation, or whilst a data link is not established, the data link layer entity shall:

- issue an MDL-ERROR indication primitive; and

- discard the frame.

NOTE – For satisfactory operation it is essential that a receiver is able to discriminate between invalid frames, as defined in E.2.9, and frames with an I-field which exceeds the maximum established length [see d) of E.3.6.11]. An unbounded frame may be assumed, and thus discarded, if two times the longest permissible frame plus two octets are received without a flag detection.

E.5.8.6 Receipt of an FRMR response frame

Receipt of an FRMR response frame defined in 5.8.6 applies with all references to "SABME" replaced by "SM".

E.5.8.7 Unsolicited response frames

Unsolicited response frames defined in 5.8.7 applies with all references to "SABME" replaced by "SM".

E.5.8.8 Duplicate assignment of a TEI value

See 5.8.8.

E.5.9 List of System parameters

See 5.9.

E.5.10 Data link layer monitor function

See 5.10.

E.6 SDL for point-to-point procedures

E.6.1 General

See B.1 of Annex B.

E.6.2 An overview of the states of the point-to-point data link layer entity

See B.2 of Annex B.

E.6.3 Cover notes

See B.3 of Annex B.

E.6.4 The use of queues

See B.4 of Annex B.

Two additional queues are added to support the SREJ procedure: a re-transmit queue (RE-TX QUEUE) and a receive queue (REC QUEUE).

RE-TX QUEUE is used to hold I frames to be re-transmitted. These I frames are requested by an SREJ response with $F = 0$ or $F = 1$, or an RR response with $F = 1$. A signal RE-TX I FRAME QUEUED UP is provided in order to cause the servicing of the RE-TX QUEUE.

REC QUEUE is used to temporarily hold out-of-sequence I frames received.

E.6.5 SDL representation

E.6.5.1 Figure E.B-3/Q.921 (sheet 1 of 3)

See Figure B-3/Q.921 (sheet 1 of 3) of Annex B.

E.6.5.2 Figure E.B-3/Q.921 (sheet 2 of 3)

See Figure B-3/Q.921 (sheet 2 of 3) of Annex B.

E.6.5.3 Figure E.B-3/Q.921 (sheet 3 of 3)

See Figure B-3/Q.921 (sheet 3 of 3) of Annex B.

E.6.5.4 Figure E.B-4/Q.921 (sheet 1 of 2)

Figure B.4/Q.921 (sheet 1 of 2) of Annex B applies with the following changes:

a) signal reception "SABME" is replaced by signal reception "SM";

b) process description "V(S) = 0

V(R) = 0

V(A) = 0"

is replaced by "V(S) = 0

V(R) = 0

V(A) = 0

V(P) = 0"

E.6.5.5 Figure E.B-4/Q.921 (sheet 2 of 2)

See Figure B.4/Q.921 (sheet 2 of 2) of Annex B.

E.6.5.6 Figure E.B-5/Q.921 (sheet 1 of 3)

Figure B.5/Q.921 (sheet 1 of 3) of Annex B applies with the following changes:

a) signal reception "SABME" is replaced by signal reception "SM";

b) process description "DISCARD I QUEUE"

is replaced by "DISCARD I, RE-TX and REC QUEUES"

E.6.5.7 Figure E.B-5/Q.921 (sheet 2 of 3)

Figure B.5/Q.921 (sheet 2 of 3) of Annex B applies with the following changes:

a) process description "V(S) = 0

V(R) = 0

V(A) = 0"

is replaced by "V(S) = 0

V(R) = 0

V(A) = 0

V(P) = 0"

b) process description "DISCARD I QUEUE"

is replaced by "DISCARD I, RE-TX and REC QUEUES"

E.6.5.8 Figure E.B-5/Q.921 (sheet 3 of 3)

Figure B.5/Q.921 (sheet 3 of 3) of Annex B applies with the following changes:

a) signal reception "SABME" is replaced by signal reception "SM";

b) process description "DISCARD I QUEUE"

is replaced by "DISCARD I, RE-TX and REC QUEUES"

E.6.5.9 Figure E.B-6/Q.921 (sheet 1 of 2)

Figure B.6/Q.921 (sheet 1 of 2) of Annex B applies with signal reception "SABME" replaced by signal reception "SM".

E.6.5.10 Figure E.B-6/Q.921 (sheet 2 of 2)

See Figure B.6/Q.921 (sheet 2 of 2) of Annex B.

E.6.5.11 Figure E.B-7/Q.921 (sheet 1 of 10)

See Figure E.B-7/Q.921 (sheet 1 of 10) of this Annex.

E.6.5.12 Figure E.B-7/Q.921 (sheet 2 of 10)

See Figure E.B-7/Q.921 (2 of 10) of this Annex.

E.6.5.13 Figure E.B-7/Q.921 (sheet 3 of 10)

Figure B.7/Q.921 (sheet 3 of 10) of Annex B applies with the following changes:

a) signal reception "SABME" is replaced by signal reception "SM";

b) process description "V(S) = 0

V(R) = 0

V(A) = 0"

is replaced by "V(S) = 0

V(R) = 0

V(A) = 0

V(P) = 0"

c) process description "DISCARD I QUEUE"

is replaced by "DISCARD I, RE-TX and REC QUEUES"

E.6.5.14 Figure E.B-7/Q.921 (sheet 4 of 10)

See Figure E.B-7/Q.921 (sheet 4 of 10) of this Annex.

E.6.5.15 Figure E.B-7/Q.921 (sheet 5 of 10)

See Figure E.B-7/Q.921 (sheet 5 of 10) of this Annex.

E.6.5.16 Figure E.B-7/Q.921 (sheet 6 of 10)

See Figure E.B-7/Q.921 (sheet 6 of 10) of this Annex.

E.6.5.17 Figure E.B-7/Q.921 (sheet 7 of 10)

See Figure B.7/Q.921 (sheet 7 of 10) of Annex B.

E.6.5.18 Figure E.B-7/Q.921 (sheet 8 of 10)

See Figure E.B-7/Q.921 (sheet 8 of 10) of this Annex.

E.6.5.19 Figure E.B-7/Q.921 (sheet 9 of 10)

See Figure B.7/Q.921 (sheet 9 of 10) of Annex B.

E.6.5.20 Figure E.B-7/Q.921 (sheet 10 of 10)

See Figure B.7/Q.921 (sheet 10 of 10) of Annex B.

E.6.5.21 Figure E.B-8/Q.921 (sheet 1 of 9)

See Figure E.B-8/Q.921 (sheet 1 of 9) of this Annex.

E.6.5.22 Figure E.B-8/Q.921 (sheet 2 of 9)

See Figure E.B-8/Q.921 (sheet 2 of 9) of this Annex.

E.6.5.23 Figure E.B-8/Q.921 (sheet 3 of 9)

Figure B.8/Q.921 (sheet 3 of 9) of Annex B applies with the following changes:

- a) signal reception "SABME" is replaced by signal reception "SM";
- b) process description $V(S) = 0$
 $V(R) = 0$
 $V(A) = 0$
is replaced by $V(S) = 0$
 $V(R) = 0$
 $V(A) = 0$
 $V(P) = 0$
- c) process description "DISCARD I QUEUE"
is replaced by "DISCARD I, RE-TX and REC QUEUES"

E.6.5.24 Figure E.B-8/Q.921 (sheet 4 of 9)

See Figure E.B-8/Q.921 (sheet 4 of 9) of this Annex.

E.6.5.25 Figure E.B-8/Q.921 (sheet 5 of 9)

See Figure E.B-8/Q.921 (sheet 5 of 9) of this Annex.

E.6.5.26 Figure E.B-8/Q.921 (sheet 6 of 9)

See Figure E.B-8/Q.921 (sheet 6 of 9) of this Annex.

E.6.5.27 Figure E.B-8/Q.921 (sheet 7 of 9)

See Figure E.B-8/Q.921 (sheet 7 of 9) of this Annex.

E.6.5.28 Figure E.B-8/Q.921 (sheet 8 of 9)

See Figure B.8/Q.921 (sheet 8 of 9) of Annex B.

E.6.5.29 Figure E.B-8/Q.921 (sheet 9 of 9)

See Figure B.8/Q.921 (sheet 9 of 9) of Annex B.

E.6.5.30 Figure E.B-9/Q.921 (sheet 1 of 5)

See Figure B.9/Q.921 (sheet 1 of 5) of Annex B.

E.6.5.31 Figure E.B-9/Q.921 (sheet 2 of 5)

See Figure B.9/Q.921 (sheet 2 of 5) of Annex B.

E.6.5.32 Figure E.B-9/Q.921 (sheet 3 of 5)

See Figure B.9/Q.921 (sheet 3 of 5) of Annex B.

E.6.5.33 Figure E.B-9/Q.921 (sheet 4 of 5)

See Figure B.9/Q.921 (sheet 4 of 5) of Annex B.

E.6.5.34 Figure E.B-9/Q.921 (sheet 5 of 5)

See Figure E.B-9/Q.921 (sheet 5 of 5) of this Annex.

E.7 State transition Tables

E.7.1 General

Subclause D.1 applies with all references to REJ and SABME replaced by SREJ and SM, respectively.

E.7.2 Key to the state transition Table

Subclause D.2 applies with all references to SABME replaced by SM.

E.7.3 State Transition Table

E.7.3.1 Table E.D-1/Q.921 (sheet 1 of 10)

Table D.1/Q.921 (sheet 1 of 10) of Annex D applies with all references to "SABME" replaced by "SM".

E.7.3.2 Table E.D-1/Q.921 (sheet 2 of 10)

Table D.1/Q.921 (sheet 2 of 10) of Annex D applies with the following changes:

- a) References to "SABME" are replaced by "SM";*
- b) References to "V(S,R,A)" are replaced by "V(S,R,A,P)";*
- c) References to "DISC I QUEUE" are replaced by "DISC I, RE-TX and REC QUEUES".*

E.7.3.3 Table E.D-1/Q.921 (sheet 3 of 10)

Table D.1/Q.921 (sheet 3 of 10) of Annex D applies.

E.7.3.4 Table E.D-1/Q.921 (sheet 4 of 10)

Table D.1/Q.921 (sheet 4 of 10) of Annex D applies.

E.7.3.5 Table E.D-1/Q.921 (sheet 5 of 10)

Table D.1/Q.921 (sheet 5 of 10) of Annex D applies with the following changes:

- a) References to "REJ" are replaced by "SREJ";*
- b) Conditions of "SREJ command P = 1" and "SREJ command P = 0" are not supported.*

E.7.3.6 Table E.D-1/Q.921 (sheet 6 of 10)

Table D.1/Q.921 (sheet 6 of 10) of Annex D applies.

E.7.3.7 Table E.D-1/Q.921 (sheet 7 of 10)

Table D.1/Q.921 (sheet 7 of 10) of Annex D applies.

E.7.3.8 Table E.D-1/Q.921 (sheet 8 of 10)

Table D.1/Q.921 (sheet 8 of 10) of Annex D applies.

E.7.3.9 Table E.D-1/Q.921 (sheet 9 of 10)

Table D.1/Q.921 (sheet 9 of 10) of Annex D applies with the following changes:

- a) References to "SABME" are replaced by "SM";*
- b) References to "DISC I QUEUE" are replaced by "DISC I, RE-TX and REC QUEUES".*

E.7.3.10 Table E.D-1/Q.921 (sheet 10 of 10)

Table D.1/Q.921 (sheet 10 of 10) of Annex D applies with all references to "SABME" replaced by "SM".

E.7.3.11 Table E.D-2/Q.921 (sheet 1 of 10)

See Table E.D-2/Q.921 (sheet 1 of 10) of this Annex.

E.7.3.12 Table E.D-2/Q.921 (sheet 2 of 10)

Table D.2/Q.921 (sheet 2 of 10) of Annex D applies with the following changes:

- a) References to "REJ" are replaced by "SREJ";*
- b) References to "SABME" are replaced by "SM";*
- c) References to "V(S,R,A)" are replaced by "V(S,R,A,P)";*
- d) References to "DISC I QUEUE" are replaced by "DISC I, RE-TX and REC QUEUES".*

E.7.3.13 Table E.D-2/Q.921 (sheet 3 of 10)

Table D.2/Q.921 (sheet 3 of 10) of Annex D applies with the following changes:

- a) References to "REJ" are replaced by "SREJ";*
- b) References to "SABME" are replaced by "SM".*

E.7.3.14 Table E.D-2/Q.921 (sheet 4 of 10)

See Table E.D-2/Q.921 (sheet 4 of 10) of this Annex.

E.7.3.15 Table E.D-2/Q.921 (sheet 5 of 10)

See Table E.D-2/Q.921 (sheet 5 of 10) of this Annex.

E.7.3.16 Table E.D-2/Q.921 (sheet 6 of 10)

See Table E.D-2/Q.921 (sheet 6 of 10) of this Annex.

E.7.3.17 Table E.D-2/Q.921 (sheet 7 of 10)

See Table E.D-2/Q.921 (sheet 7 of 10) of this Annex.

E.7.3.18 Table E.D-2/Q.921 (sheet 8 of 10)

See Table E.D-2/Q.921 (8 of 10) of this Annex.

E.7.3.19 Table E.D-2/Q.921 (sheet 9 of 10)

See Table E.D-2/Q.921 (sheet 9 of 10) of this Annex.

E.7.3.20 Table E.D-2/Q.921 (sheet 10 of 10)

Table D.2/Q.921 (sheet 10 of 10) of Annex D applies with the following changes:

- a) *References to "REJ" are replaced by "SREJ";*
- b) *References to "SABME" are replaced by "SM".*

E.7.3.21 Table E.D-3/Q.921 (sheet 1 of 10)

See Table E.D-3/Q.921 (sheet 1 of 10) of this Annex.

E.7.3.22 Table E.D-3/Q.921 (sheet 2 of 10)

Table D.3/Q.921 (sheet 2 of 10) of Annex D applies with the following changes:

- a) *References to "REJ" are replaced by "SREJ";*
- b) *References to "SABME" are replaced by "SM";*
- c) *References to "V(S,R,A)" are replaced by "V(S,R,A,P)";*
- d) *References to "DISC I QUEUE" are replaced by "DISC I, RE-TX and REC QUEUES".*

E.7.3.23 Table E.D-3/Q.921 (sheet 3 of 10)

Table D.3/Q.921 (sheet 3 of 10) of Annex D applies with the following changes:

- a) *References to "REJ" are replaced by "SREJ";*
- b) *References to "SABME" are replaced by "SM".*

E.7.3.24 Table E.D-3/Q.921 (sheet 4 of 10)

See Table E.D-3/Q.921 (sheet 4 of 10) of this Annex.

E.7.3.25 Table E.D-3/Q.921 (sheet 5 of 10)

See Table E.D-3/Q.921 (sheet 5 of 10) of this Annex.

E.7.3.26 Table E.D-3/Q.921 (sheet 6 of 10)

See Table E.D-3/Q.921 (sheet 6 of 10) of this Annex.

E.7.3.27 Table E.D-3/Q.921 (sheet 7 of 10)

See Table E.D-3/Q.921 (sheet 7 of 10) of this Annex.

E.7.3.28 Table E.D-3/Q.921 (sheet 8 of 10)

See Table E.D-3/Q.921 (sheet 8 of 10) of this Annex.

E.7.3.29 Table E.D-3/Q.921 (sheet 9 of 10)

See Table E.D-3/Q.921 (sheet 9 of 10) of this Annex.

E.7.3.30 Table E.D-3/Q.921 (sheet 10 of 10)

Table D.3/Q.921 (sheet 10 of 10) of Annex D applies with the following changes:

- a) *References to "REJ" are replaced by "SREJ";*
- b) *References to "SABME" are replaced by "SM".*

E.8 Examples of the use of multi-selective reject option

This subclause shows examples of the use of multi-selective reject option as described in the Annex C of ISO/IEC 4335.

Figure E.8-1 shows the frame exchange between two communicating data link layer entities when I frames are lost and recovered by retransmissions using the SREJ frame with the F bit set to "0".

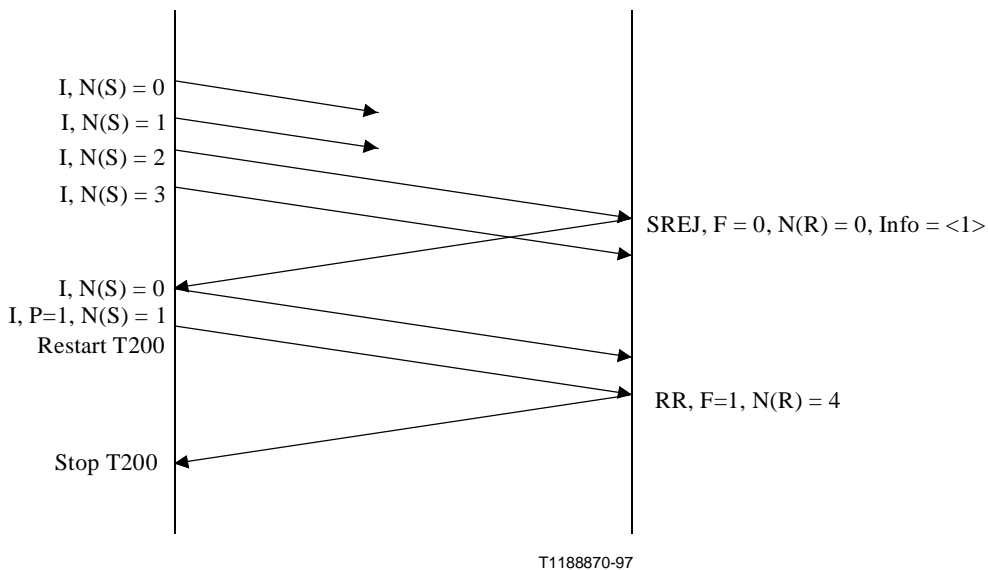


Figure E.8-1/Q.921 – I frame recovery due to SREJ frame with F bit set to "0"

Figure E.8-2 shows the frame exchange between two communicating data link layer entities when I frames are lost and the resulting SREJ frame with the F bit set to "0" is also lost.

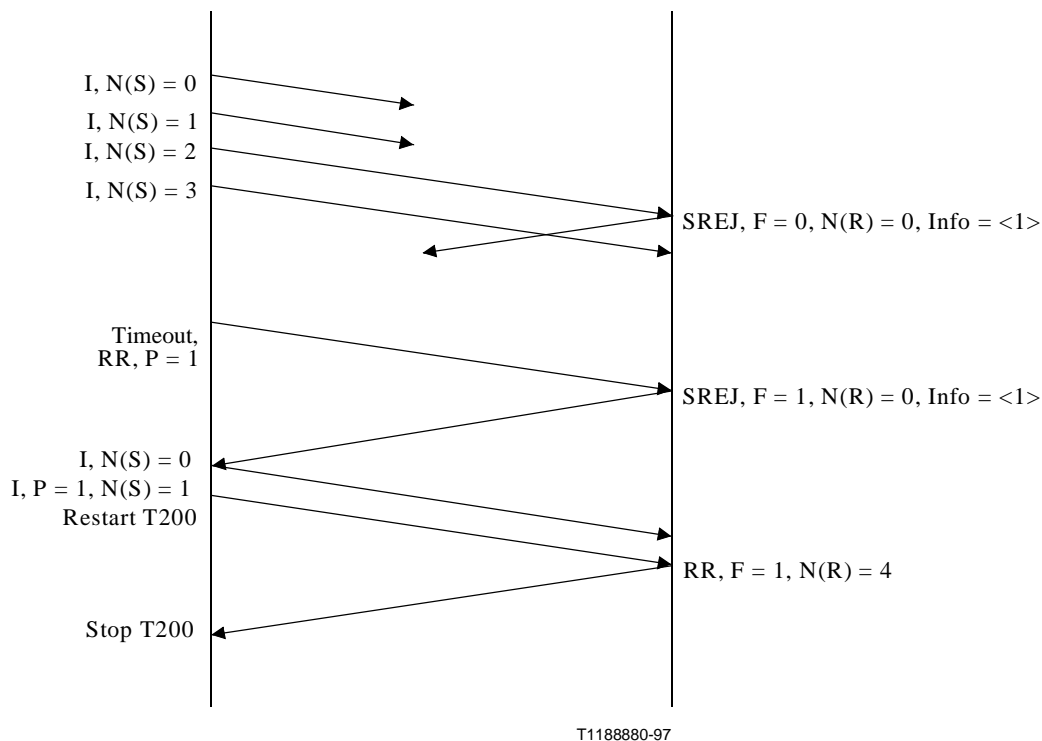


Figure E.8-2/Q.921 – I frame recovery, when SREJ frame with F bit set to "0" is lost

Figure E.8-3 shows the frame exchange between two communicating data link layer entities when the last few I frames in a sequence of I frames are lost.

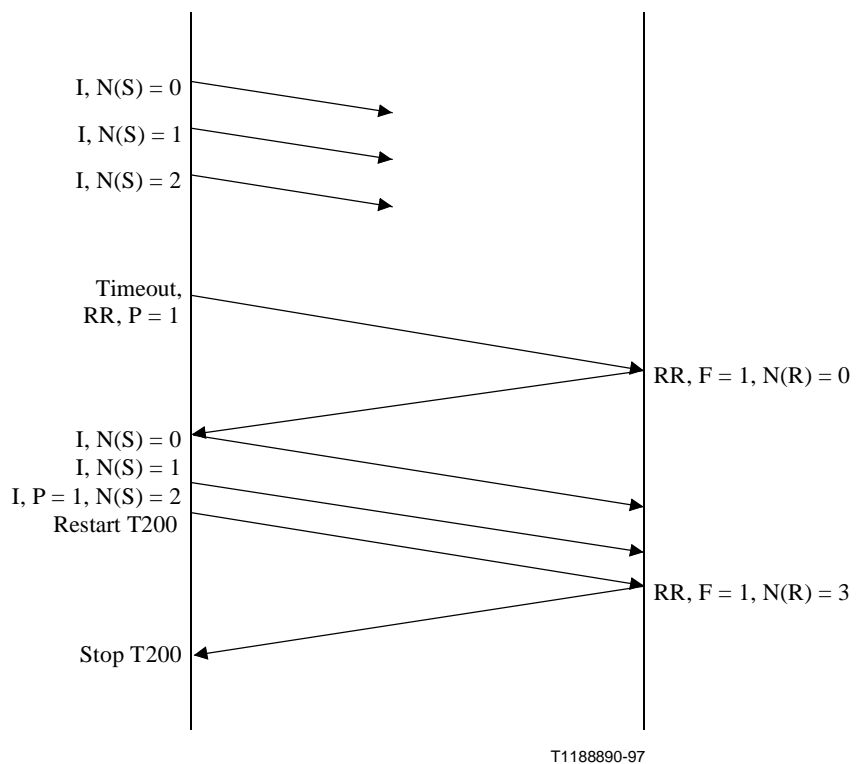


Figure E.8-3/Q.921 – I frame recovery, when last few I frames in sequence of I frames are lost

Figure E.8-4 shows a more complex exchange of frames between two communicating data link layer entities where retransmitted I frames are lost.

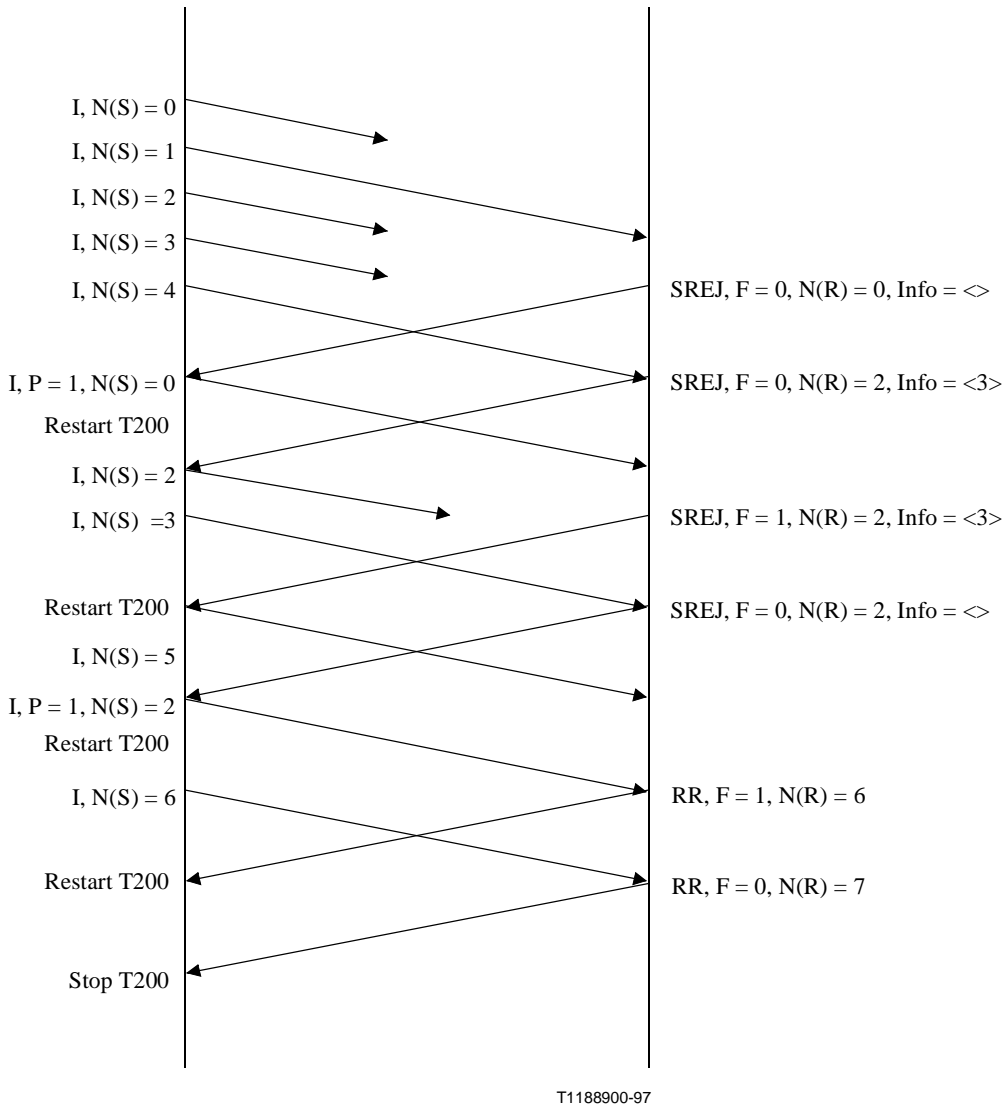


Figure E.8-4/Q.921 – I frame recovery, when retransmitted I frames are lost

Figure E.8-5 shows a more complex exchange of frames between two communicating data link layer entities where the multiple I frames, including the last I frame in a sequence of I frames, and SREJ frames are lost.

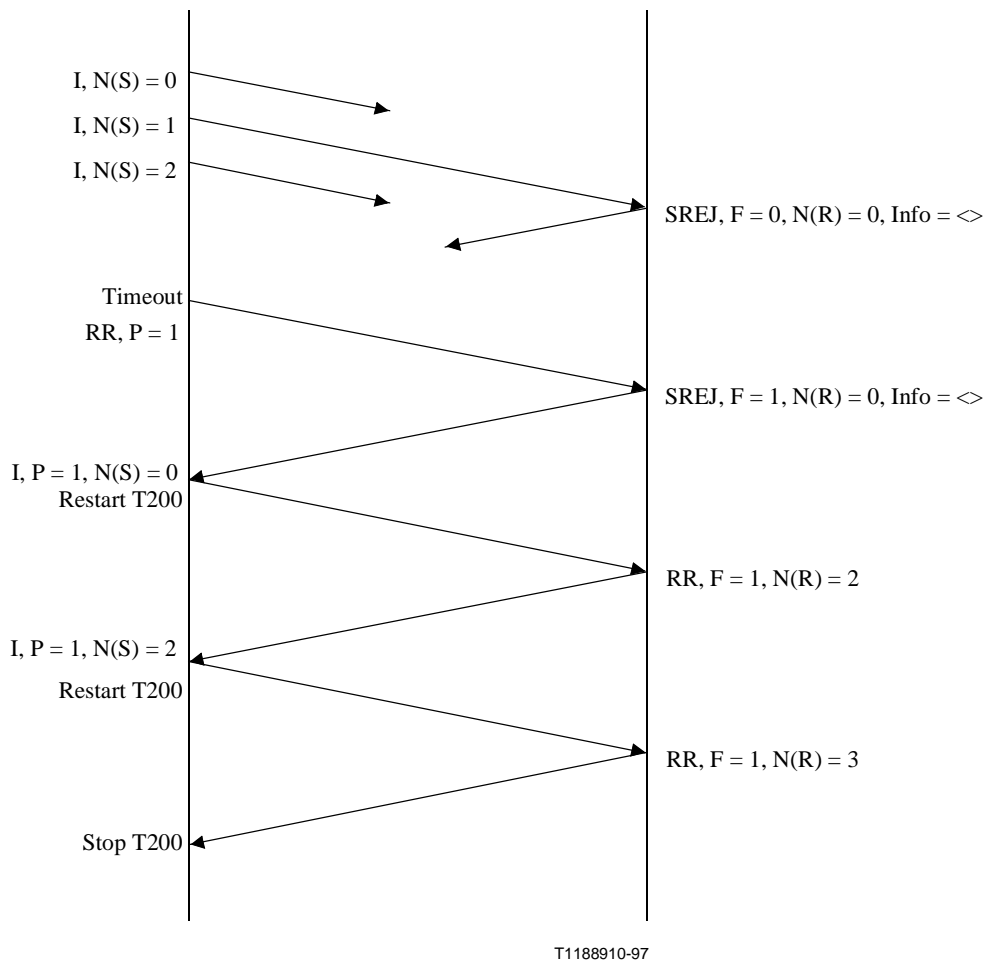
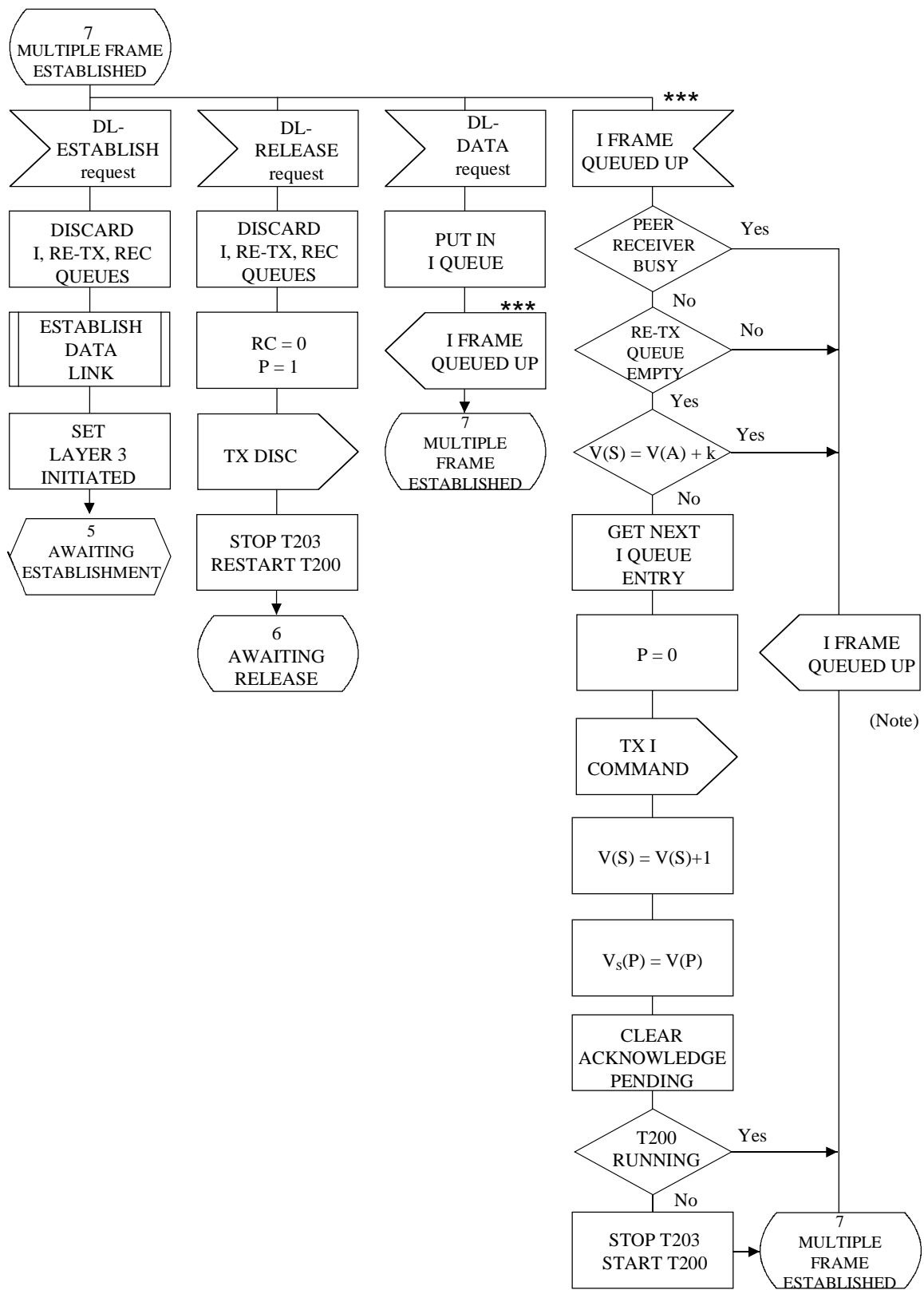


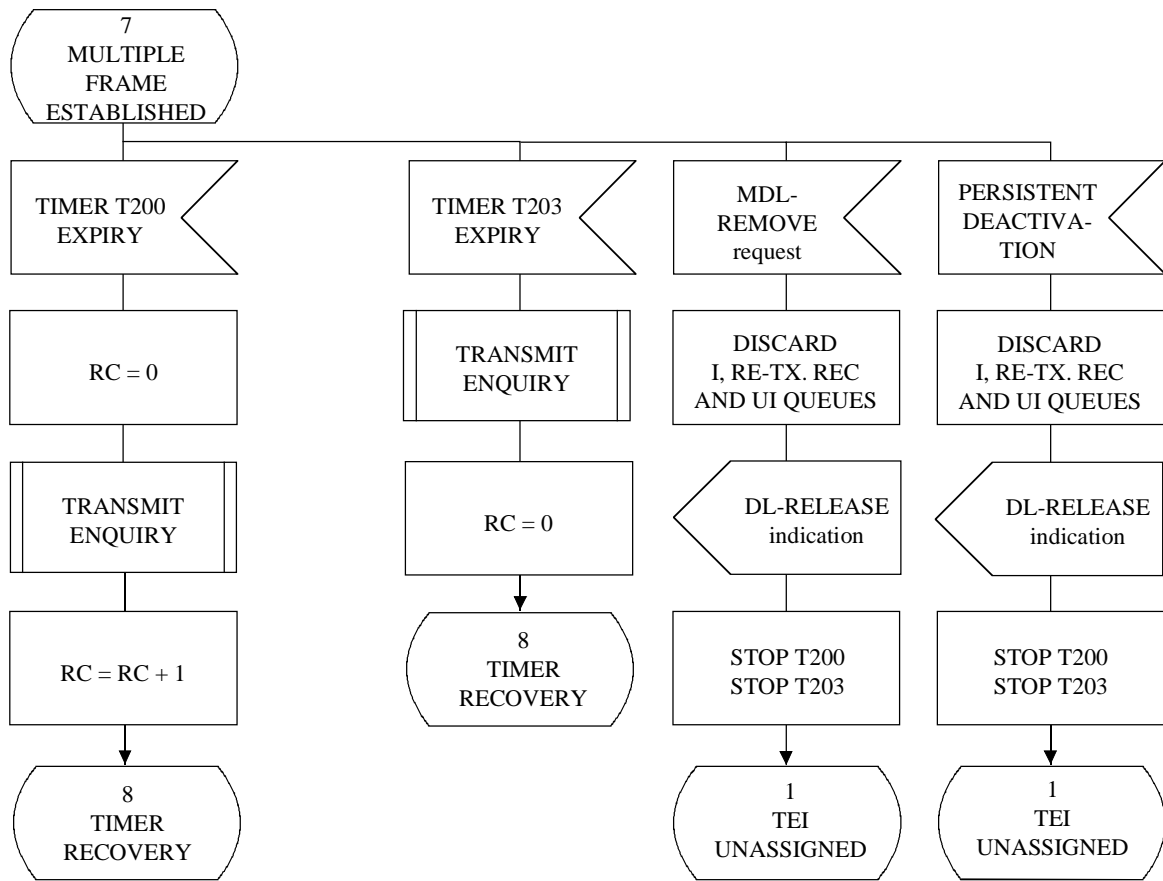
Figure E.8-5/Q.921 – I frame recovery, when multiple I frames, last I frame and SREJ frames are lost



T1188920-97

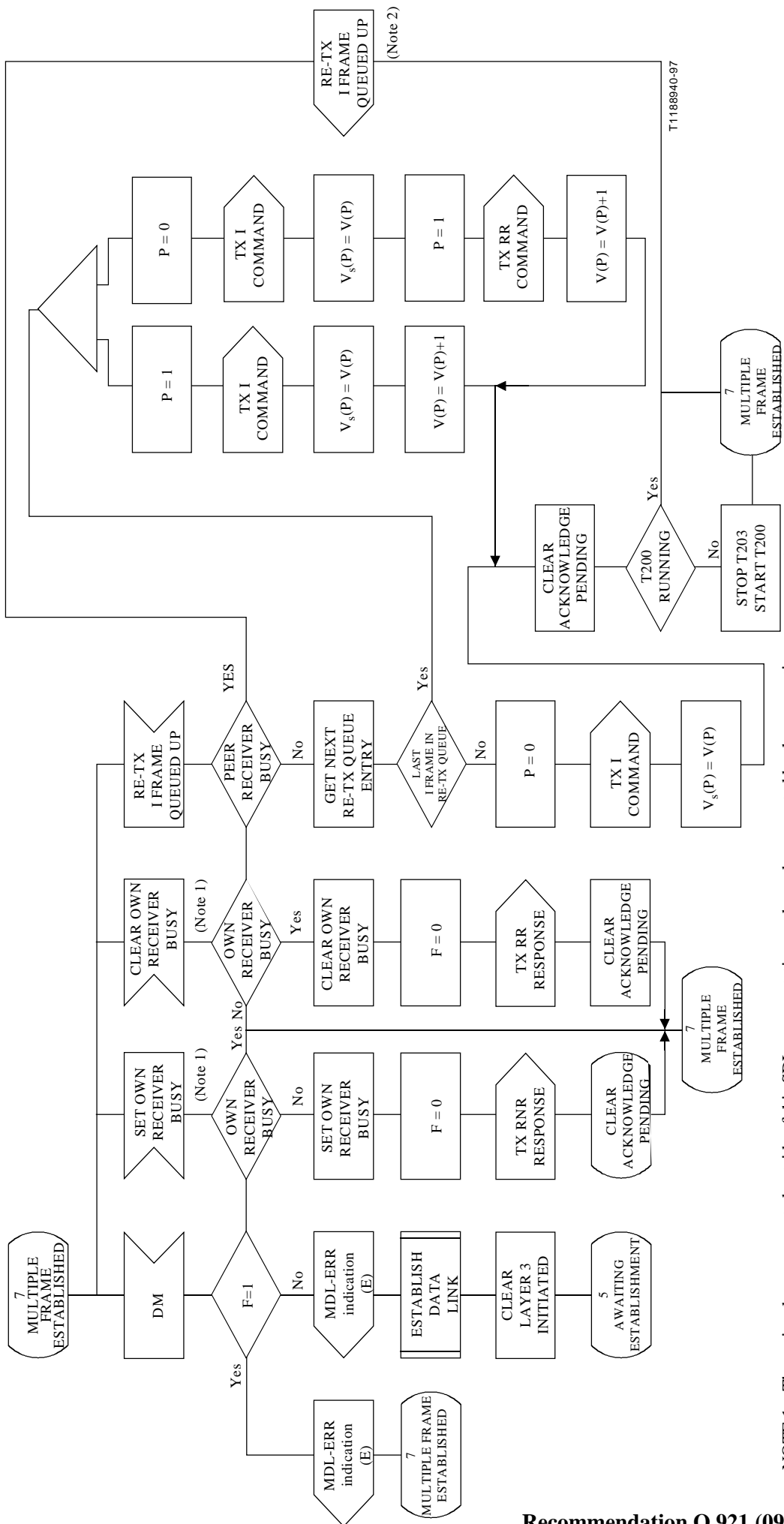
NOTE – The regeneration of this signal does not affect the sequence integrity of the queue.

Figure E.B-7/Q.921 (sheet 1 of 10)



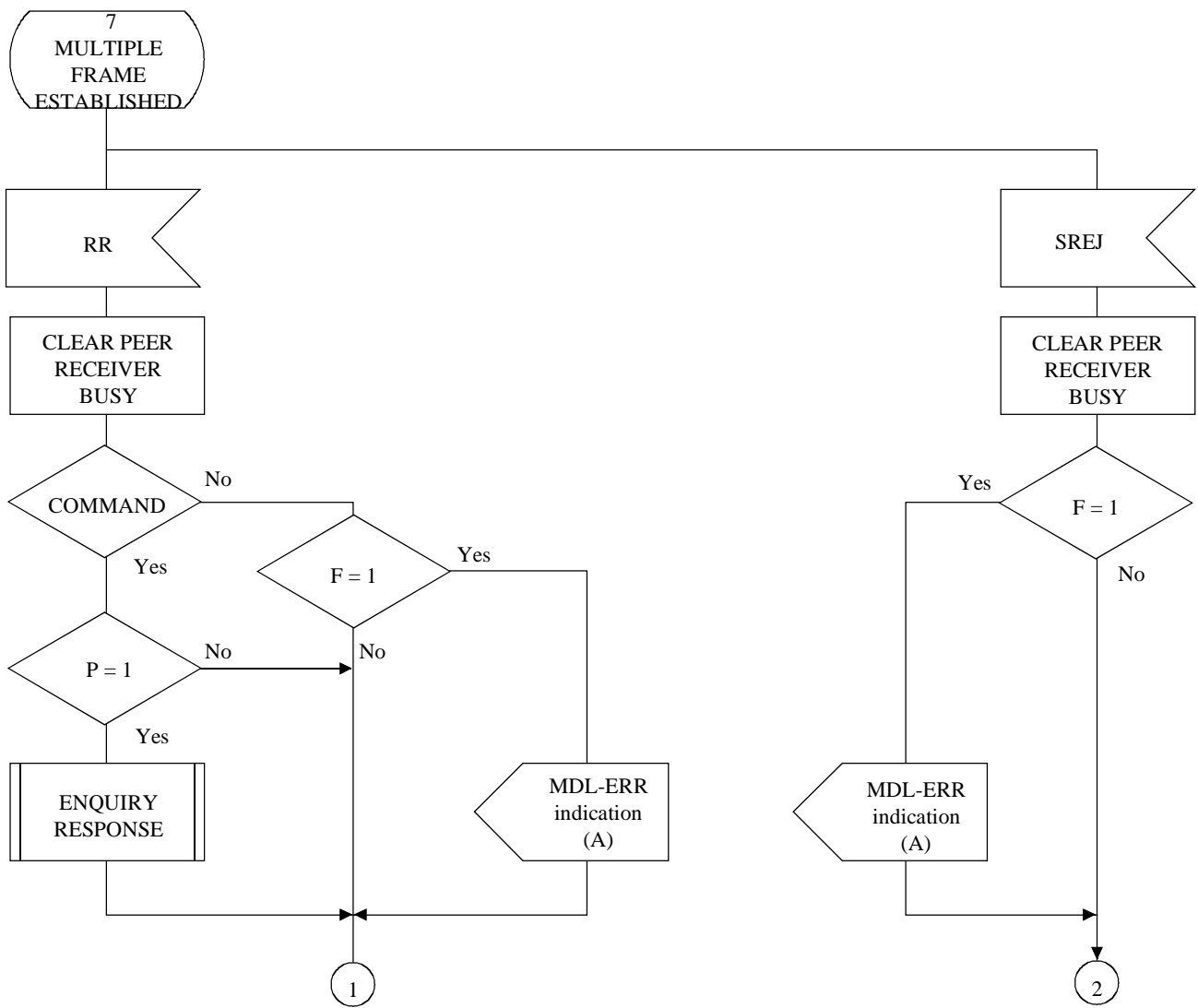
T1188930-97

Figure E.B-7/Q.921 (sheet 2 of 10)



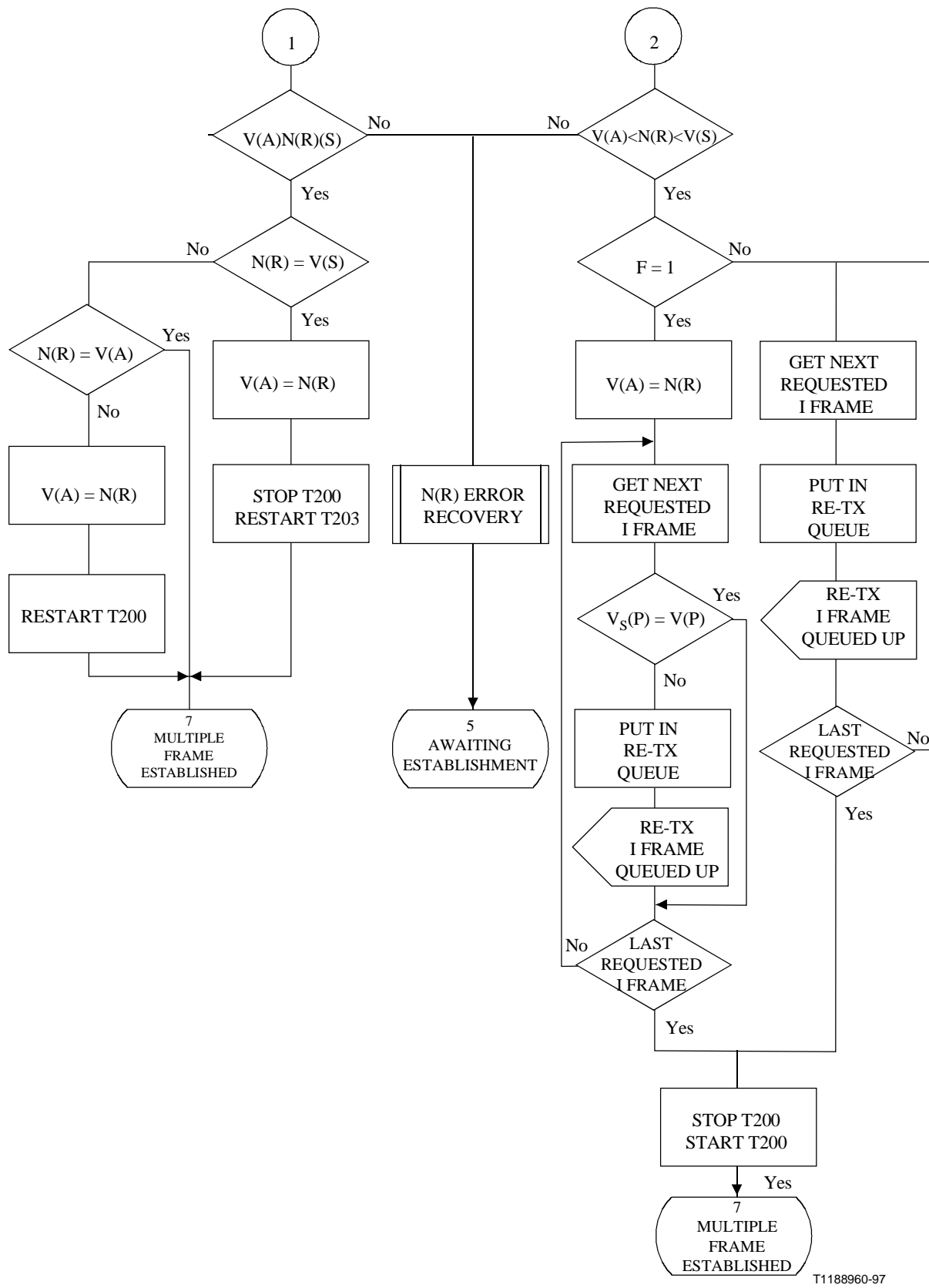
NOTE 1 – These signals are generated outside of this SDL representation, and may be generated by the connection management.
 NOTE 2 – The regeneration of this signal does not affect the sequence integrity of the queue.

Figure E.B-7/Q.921 (sheet 4 of 10)



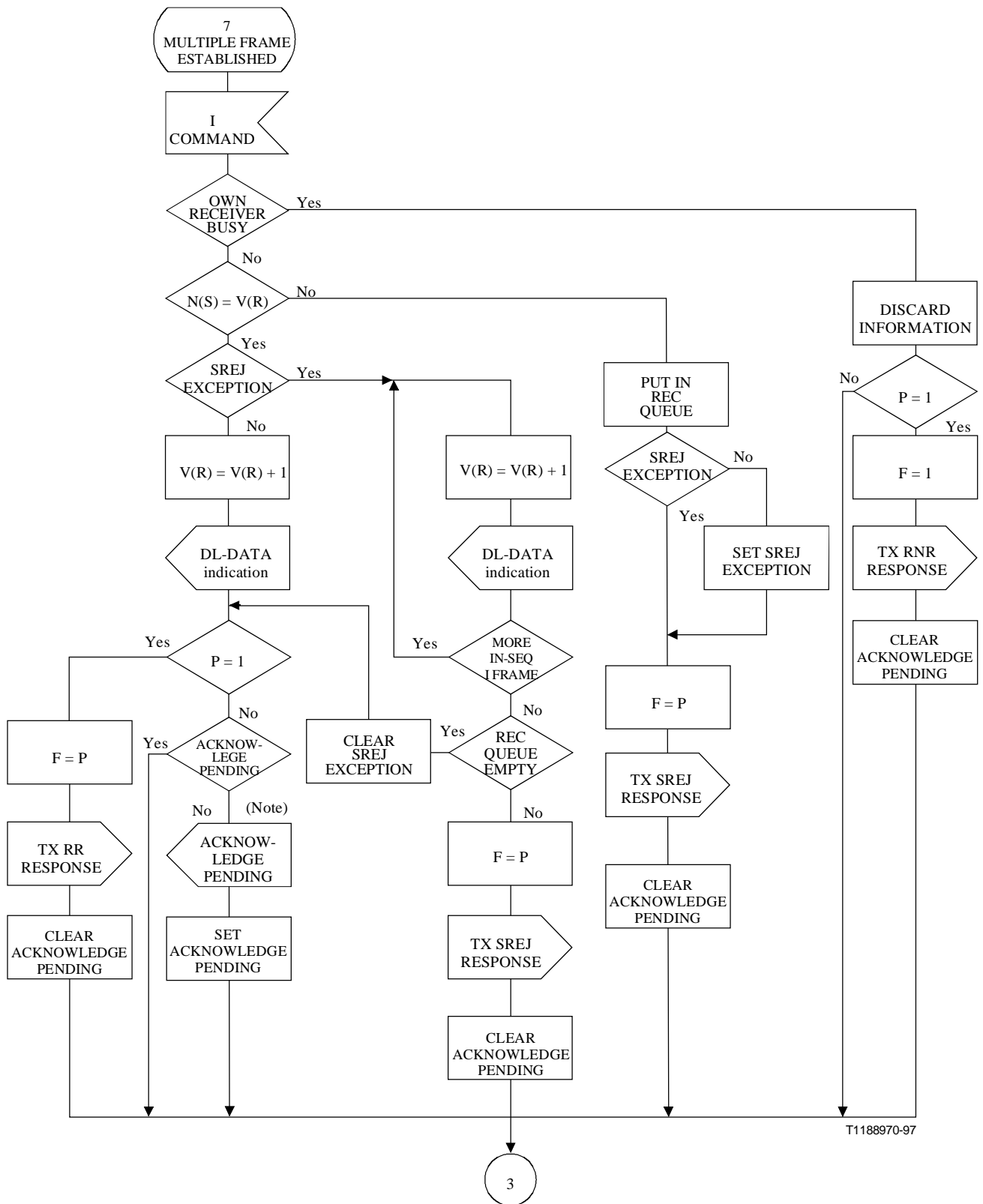
T1188950-97

Figure E.B-7/Q.921 (sheet 5 of 10)



T1188960-97

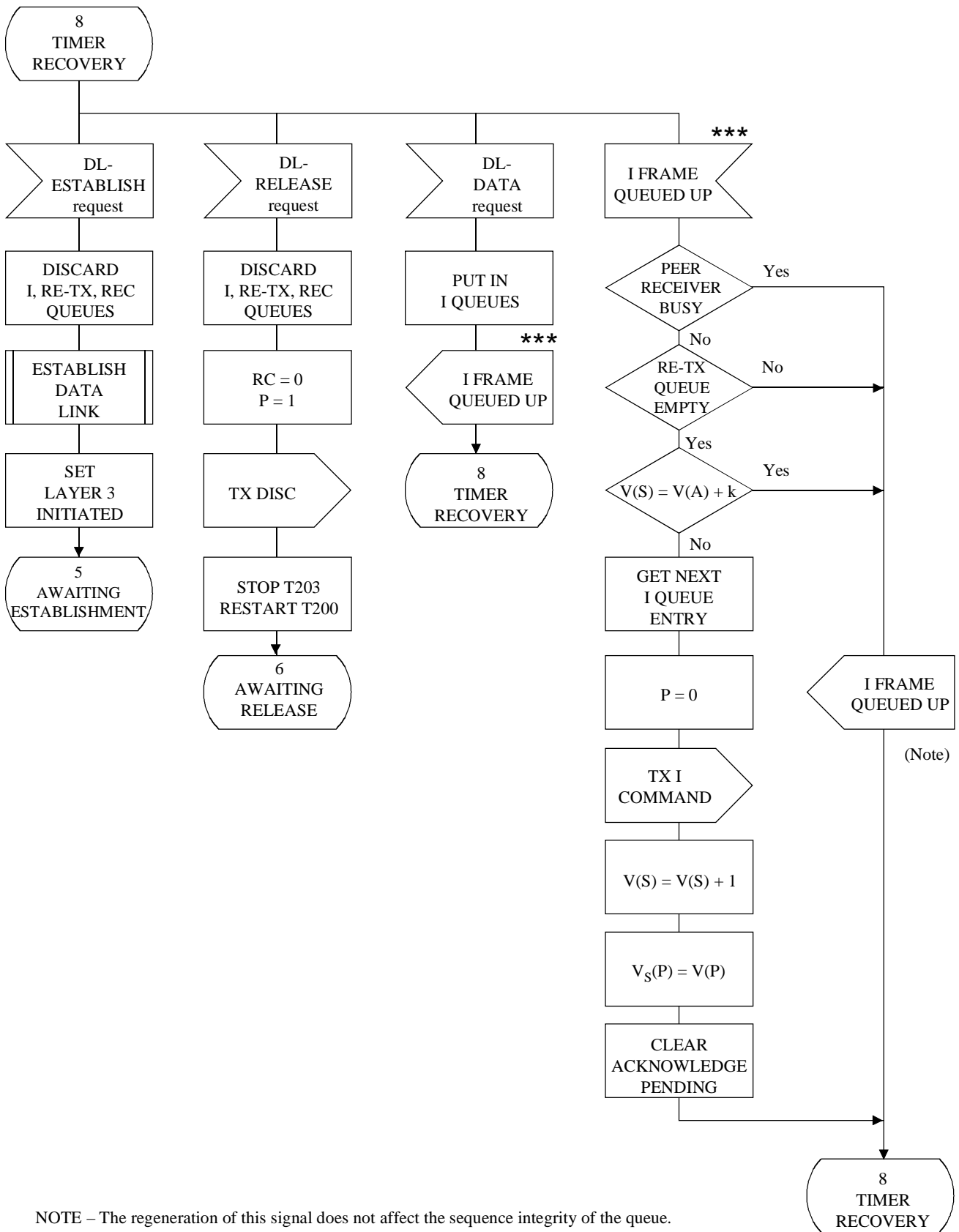
Figure E.B-7/Q.921 (sheet 6 of 10)



T1188970-97

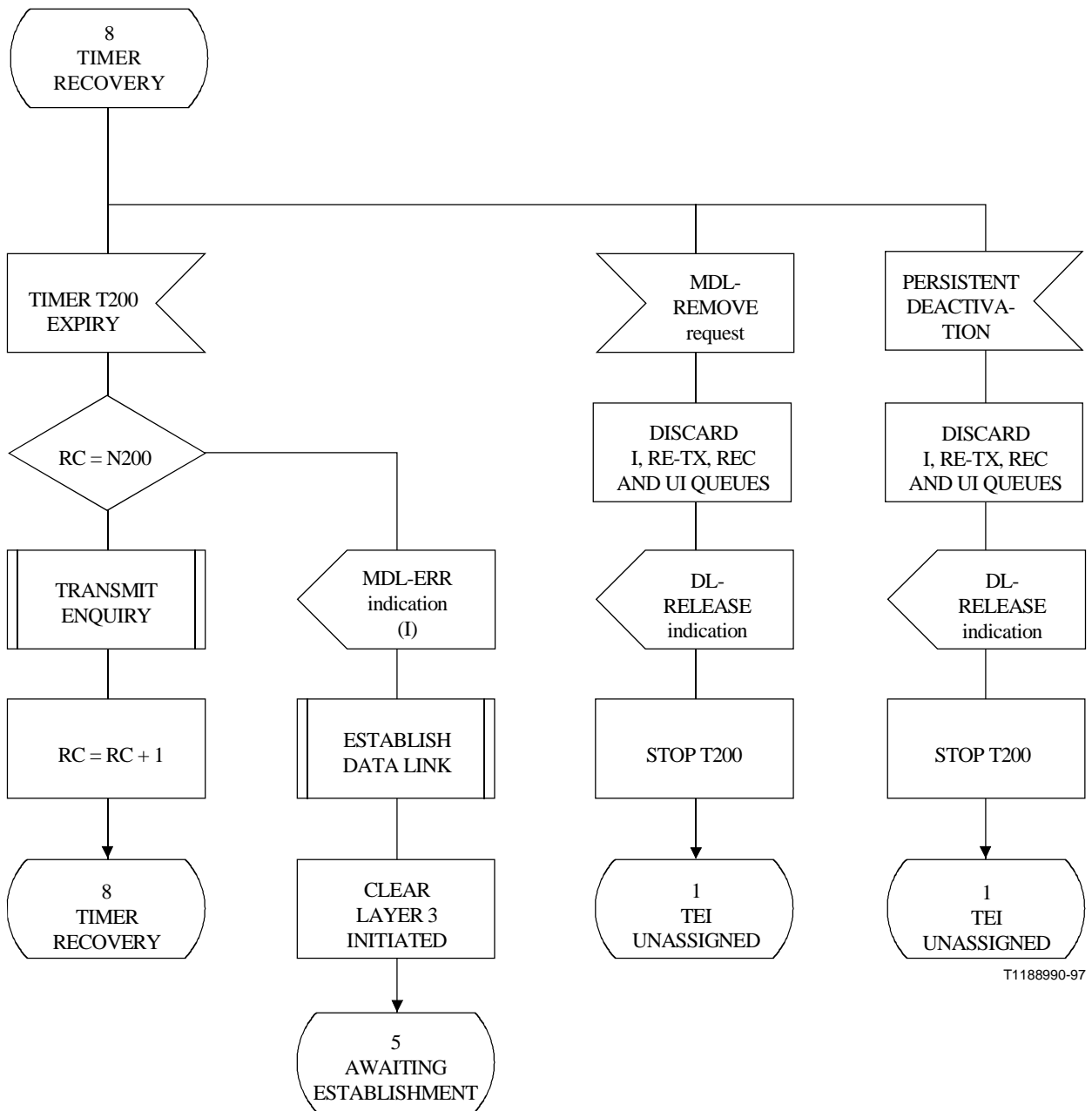
NOTE – Processing of acknowledge pending is described on sheet 10 of Figure B.7/Q.921.

Figure E.B-7/Q.921 (sheet 8 of 10)



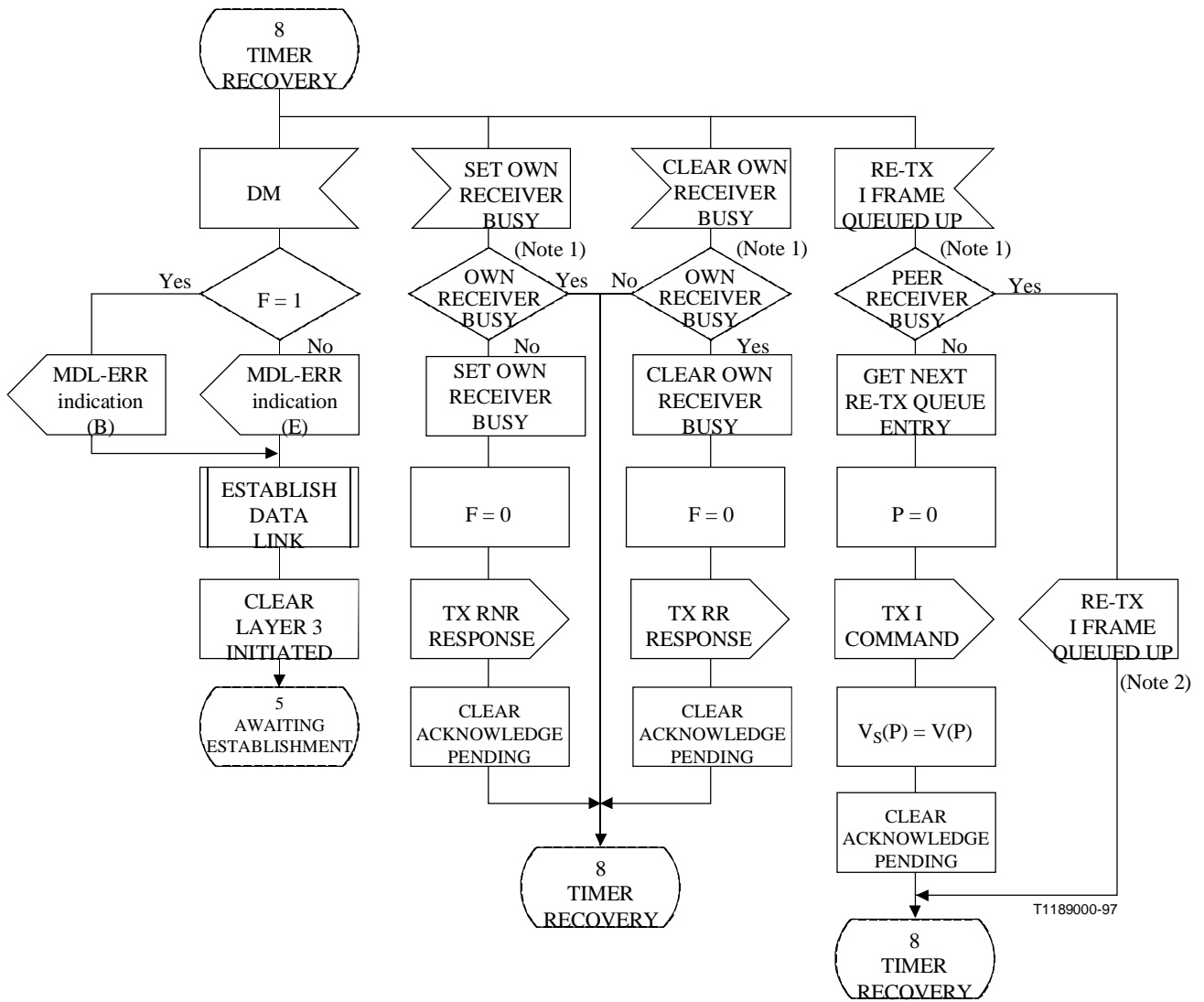
T1188980-97

Figure E.B-8/Q.921 (sheet 1 of 9)



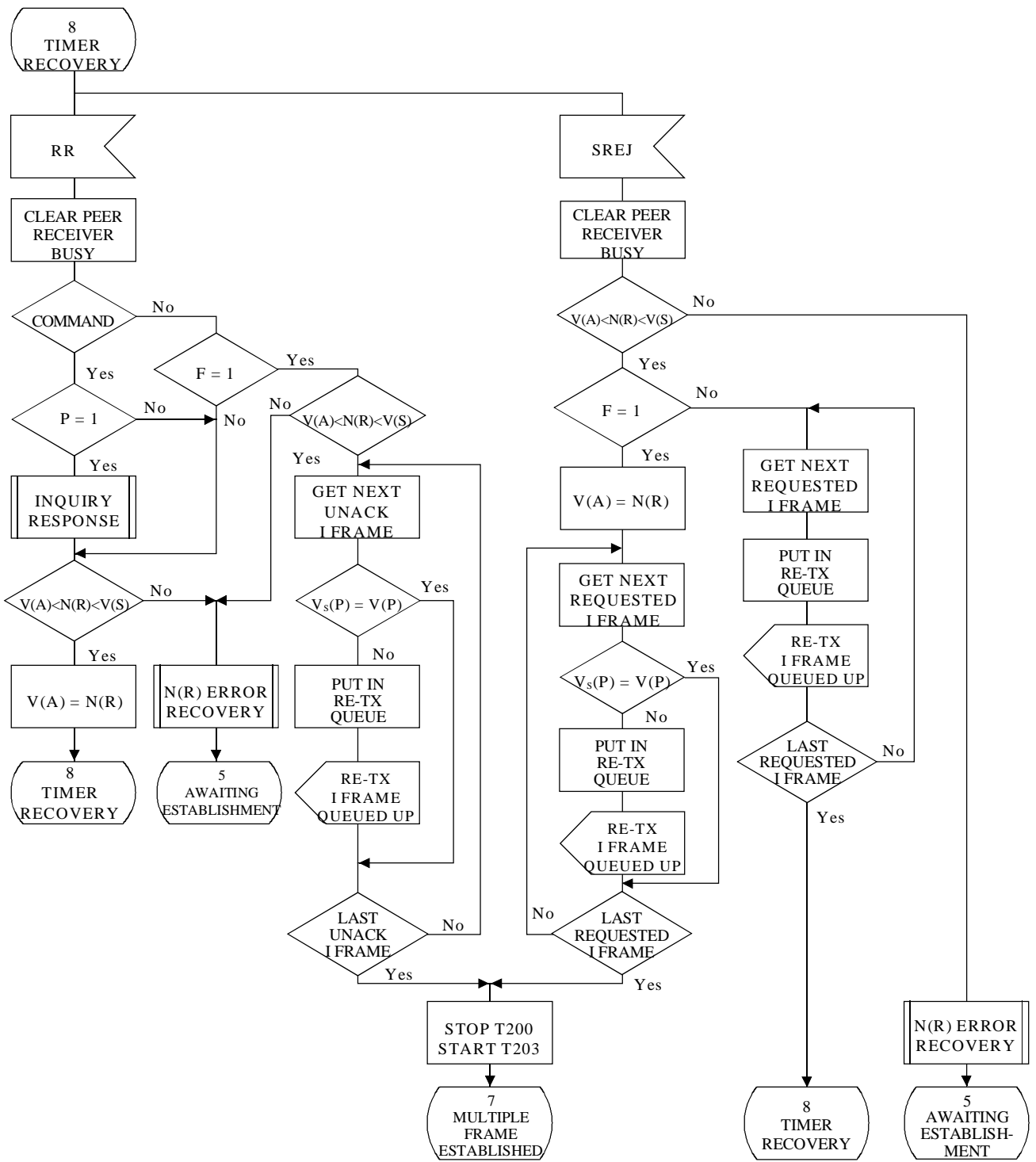
T1188990-97

Figure E.B-8/Q.921 (sheet 2 of 9)



NOTE 1 – These signals are generated outside of this SDL representation, and may be generated by the connection management.
 NOTE 2 – The regeneration of this signal does not affect the sequence integrity of the queue.

Figure E.B-8/Q.921 (sheet 4 of 9)



T1189010-97

Figure E.B-8/Q.921 (sheet 5 of 9)

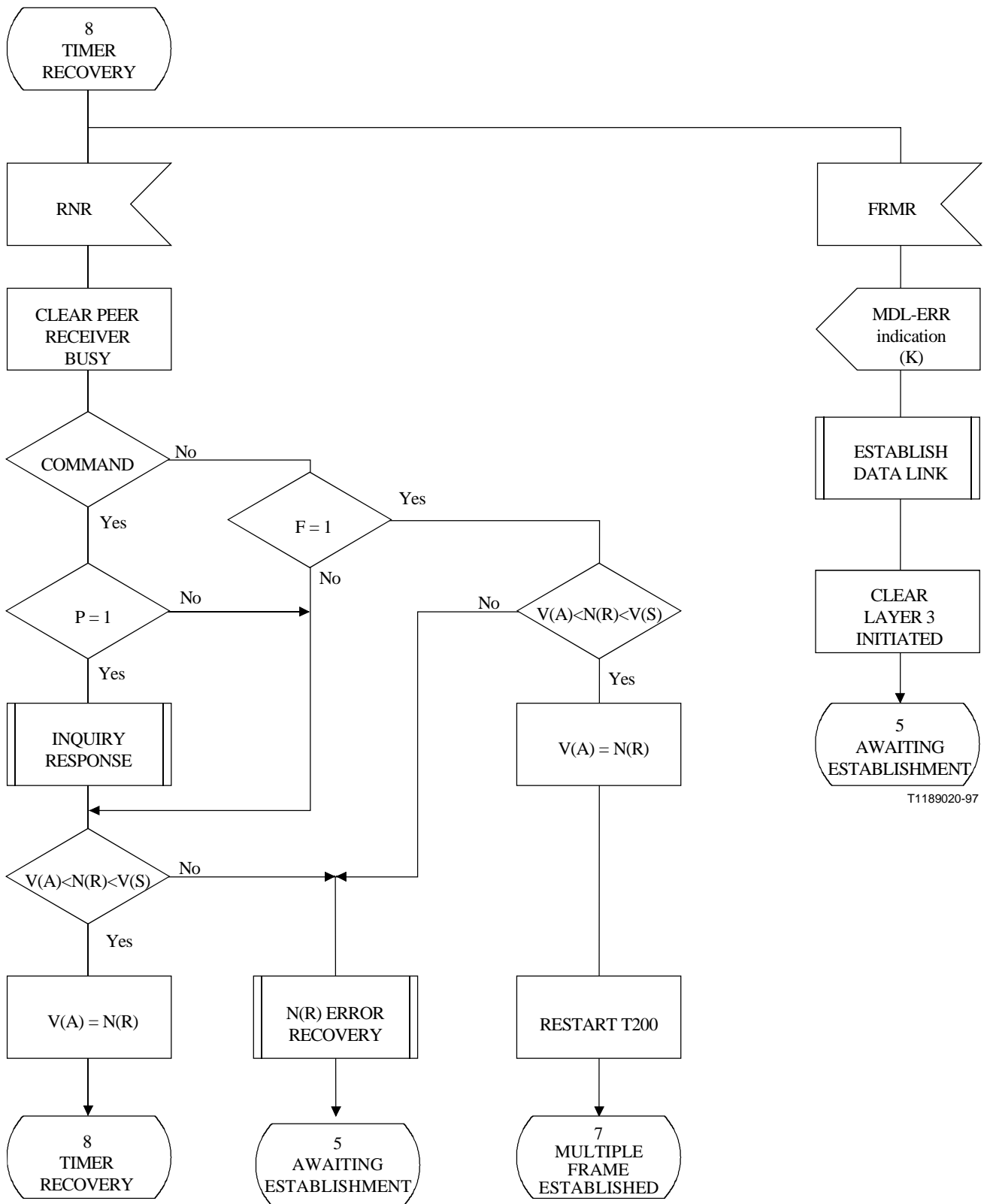
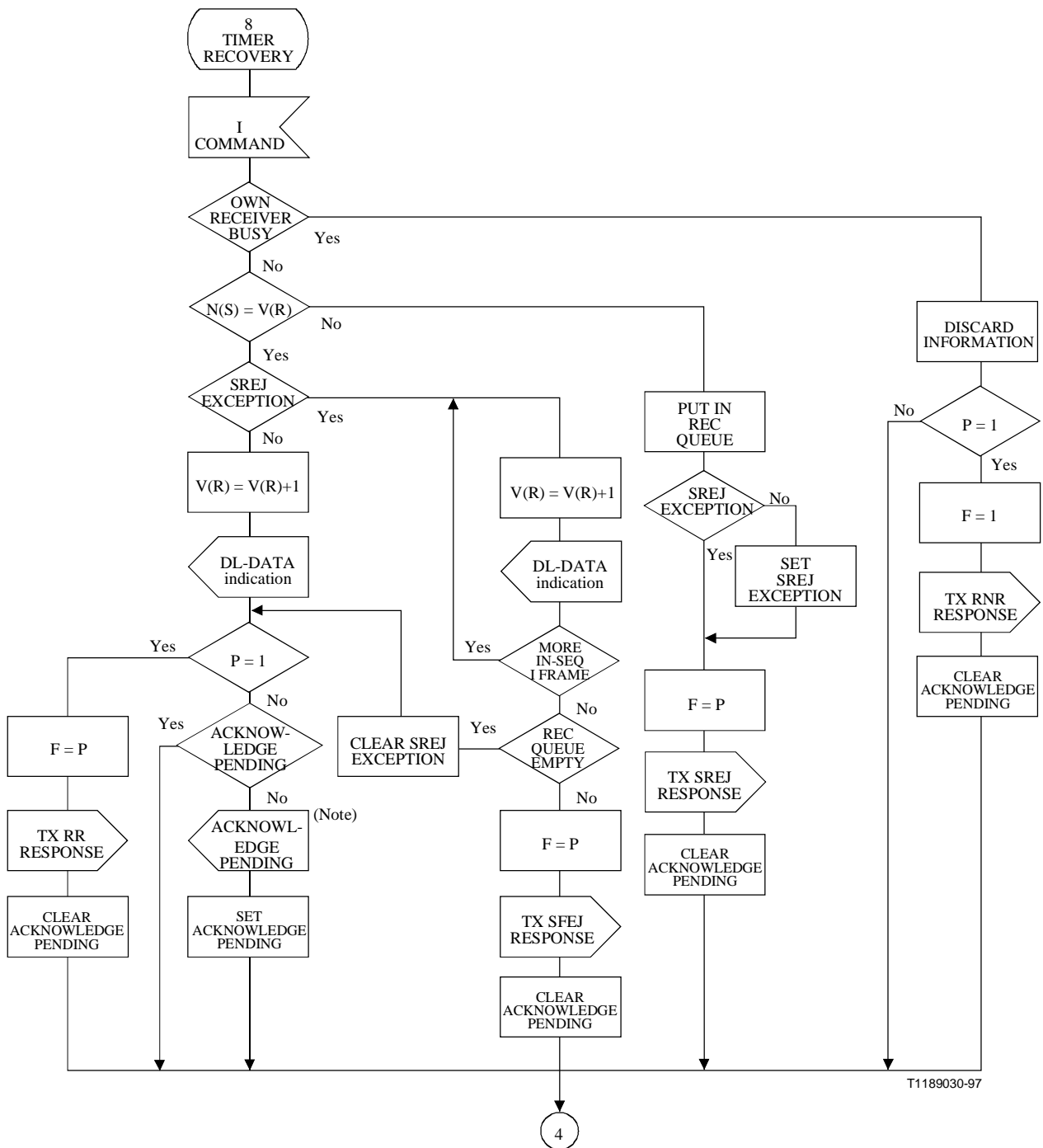


Figure E.B-8/Q.921 (sheet 6 of 9)



NOTE – Processing of acknowledge pending is described on sheet 9 of Figure B-8/Q.921.

Figure E.B-8/Q.921 (sheet 7 of 9)

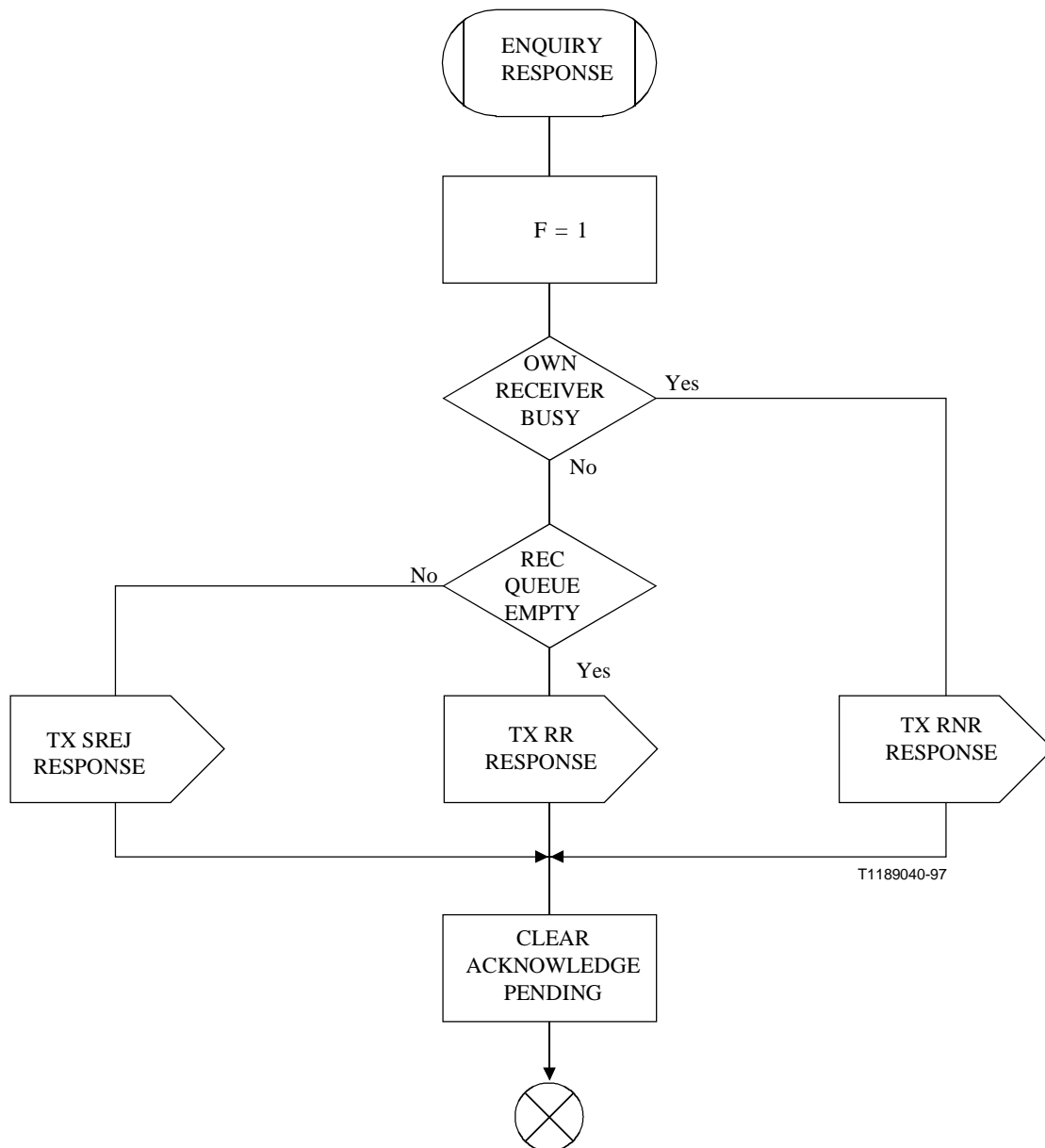


Figure E.B-9/Q.921 (sheet 5 of 5)

Table E.D-2/Q.921 (sheet 1 of 10) – State transition table: receiving primitive

BASIC STATE	MULTIPLE FRAME ESTABLISHED									
	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	<i>SREJ</i> recovery	Own rec busy	<i>SREJ</i> and Own rec Busy	Normal	<i>SREJ</i> recovery	<i>SREJ</i> recovery	<i>SREJ</i> recovery	Own rec busy	<i>SREJ</i> and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7		
DL-ESTABLISH request	DISC I RE-TX and <i>REC QUEUES</i> RC = 0 TX SM P = 1 Stop T203 Restart T200 5.0									
DL-RELEASE request	DISC I RE-TX and <i>REC QUEUES</i> RC = 0 TX DISC P = 1 Stop T203 Restart T200 6									
DL-DATA request	DATA INTO I QUEUE									
I FRAME IN QUEUE $V(S) < V(A) + k$ RE-TX QUEUE EMPTY	TX IP = 0 $V(S) = V(S) + 1$ $VS(P) = V(P)$ Stop T203 TIMER T200				LEAVE I FRAME IN QUEUE					
I FRAME IN QUEUE $V(S) < V(A) + k$ RE-TX QUEUE NOT EMPTY	LEAVE I FRAME IN QUEUE									
I FRAME IN QUEUE $V(S) = V(A) + k$										

Table E.D-2/Q.921 (sheet 1 of 10) – State transition table: receiving primitive (concluded)

BASIC STATE		MULTIPLE FRAME ESTABLISHED									
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	<i>SREJ</i> recovery	Own rec busy	<i>SREJ</i> and Own rec Busy	Normal	Normal	<i>SREJ</i> recovery	Own rec busy	<i>SREJ</i> and own rec busy	Own rec busy	<i>SREJ</i> and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7			
DL-UNIT DATA request	UNIT DATA INTO UI QUEUE										
UI FRAME IN QUEUE	TX UI P = 0										
MDL-ASSIGN request											
MDL-REMOVE request	DL-REL.ind DISC I, RE-TX, REC and UI QUEUES Stop T200 Stop T203 1										
MDL-ERROR response											
PERSISTENT DEACTIVATION	DL-REL.ind DISC I, RE-TX, REC and UI QUEUES Stop T200 Stop T203 4										

Table E.D-2/Q.921 (sheet 4 of 10) – State transition table: receiving RR supervisory frame with correct format

BASIC STATE		MULTIPLE FRAME ESTABLISHED										
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Peer rec Busy	Peer rec Busy	Peer rec Busy	Peer rec Busy	Peer rec Busy	Peer rec Busy	
RECEIVER CONDITION	Normal	SREJ recovery	Own rec Busy	SREJ and own rec Busy	Normal	Peer rec Busy	SREJ recovery	Own rec Busy	SREJ and own rec Busy	Peer rec Busy	SREJ and own rec Busy	
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7	7.8	7.9	7.10	
RR cmd P = 1 N(R) = V(S)	TX RR F = 1 Stop T200 Restart T203 V(A) = N(R)	TX SREJ F = 1 Stop T200 Restart T203 V(A) = N(R)	TX RNR F = 1 Stop T200 Restart T203 V(A) = N(R)		TX RR F = 1 Stop T200 Start T203 V(A) = N(R)	TX SREJ F = 1 Stop T200 Start T203 V(A) = N(R)	TX RNR F = 1 Stop T200 Start T203 V(A) = N(R)	TX RNR F = 1 Stop T200 Start T203 V(A) = N(R)	TX RNR F = 1 Stop T200 Start T203 V(A) = N(R)	TX RNR F = 1 Stop T200 Start T203 V(A) = N(R)	TX RNR F = 1 Stop T200 Start T203 V(A) = N(R)	TX RNR F = 1 Stop T200 Start T203 V(A) = N(R)
RR cmd P = 0 N(R) = V(S)	Stop T200 Restart T203 V(A) = N(R)				Stop T200 Start T203 V(A) = N(R)	Stop T200 Start T203 V(A) = N(R)	Stop T200 Start T203 V(A) = N(R)	Stop T200 Start T203 V(A) = N(R)	Stop T200 Start T203 V(A) = N(R)	Stop T200 Start T203 V(A) = N(R)	Stop T200 Start T203 V(A) = N(R)	Stop T200 Start T203 V(A) = N(R)
RR resp F = 1 N(R) = V(S)	MDL-ERR ind(A) Stop T200 Restart T203 V(A) = N(R)				MDL-ERR ind(A) Stop T200 Start T203 V(A) = N(R)	MDL-ERR ind(A) Stop T200 Start T203 V(A) = N(R)	MDL-ERR ind(A) Stop T200 Start T203 V(A) = N(R)	MDL-ERR ind(A) Stop T200 Start T203 V(A) = N(R)	MDL-ERR ind(A) Stop T200 Start T203 V(A) = N(R)	MDL-ERR ind(A) Stop T200 Start T203 V(A) = N(R)	MDL-ERR ind(A) Stop T200 Start T203 V(A) = N(R)	MDL-ERR ind(A) Stop T200 Start T203 V(A) = N(R)
RR cmd P = 1 V(A) < N(R) < V(S)	TX RR F = 1 Restart T200 V(A) = N(R)	TX SREJ F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)		TX RR F = 1 Restart T200 V(A) = N(R)	TX SREJ F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)
RR cmd P = 0 V(A) < N(R) < V(S)	Restart T200 V(A) = N(R)				Restart T200 V(A) = N(R)	Restart T200 V(A) = N(R)	Restart T200 V(A) = N(R)	Restart T200 V(A) = N(R)	Restart T200 V(A) = N(R)	Restart T200 V(A) = N(R)	Restart T200 V(A) = N(R)	Restart T200 V(A) = N(R)
RR resp F = 0 V(A) < N(R) < V(S)												
RR resp F = 1 V(A) < N(R) < V(S)	MDL-ERR ind(A) Restart T200 V(A) = N(R)				MDL-ERR ind(A) Restart T200 V(A) = N(R)	MDL-ERR ind(A) Restart T200 V(A) = N(R)	MDL-ERR ind(A) Restart T200 V(A) = N(R)	MDL-ERR ind(A) Restart T200 V(A) = N(R)	MDL-ERR ind(A) Restart T200 V(A) = N(R)	MDL-ERR ind(A) Restart T200 V(A) = N(R)	MDL-ERR ind(A) Restart T200 V(A) = N(R)	MDL-ERR ind(A) Restart T200 V(A) = N(R)
RR cmd P = 1 V(A) = N(R) < V(S)	TX RR F = 1	TX SREJ F = 1	TX RNR F = 1		TX RR F = 1	TX SREJ F = 1	TX RNR F = 1	TX RNR F = 1	TX RNR F = 1	TX RNR F = 1	TX RNR F = 1	TX RNR F = 1

Table E.D-2/Q.921 (sheet 4 of 10) – State transition table: receiving RR supervisory frame with correct format (concluded)

BASIC STATE	MULTIPLE FRAME ESTABLISHED									
	Normal	Normal	Normal	Normal	Peer rec Busy	Peer rec Busy	Peer rec Busy	Peer rec Busy	Peer rec Busy	Peer rec Busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec Busy	Peer rec Busy	Peer rec Busy	Peer rec Busy	Peer rec Busy	Peer rec Busy
RECEIVER CONDITION	Normal	Own rec Busy	Own rec Busy	SREJ and own rec Busy	Normal	SREJ recovery	SREJ recovery	Own rec Busy	Own rec Busy	SREJ and own rec Busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.5	7.6	7.6	7.7
RR cmd P = 0 V(A) = N(R) < V(S)	-	-	-	-	-	7.0	7.0	7.2	7.2	7.3
RR resp F = 0 V(A) = N(R) < V(S)	-	-	-	-	-	-	-	-	-	-
RR resp F = 1 V(A) = N(R) < V(S)	MDL-ERR ind(A)						MDL-ERR ind(A) 7.0	MDL-ERR ind(A) 7.1	MDL-ERR ind(A) 7.2	MDL-ERR ind(A) 7.3
RR cmd P = 1 N(R) error	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1
RR cmd P = 0 N(R) error	MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1						MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1			
RR resp F = 0 N(R) error										
RR resp F = 1 N(R) error	MDL-ERR ind(A) MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1						MDL-ERR ind(A) MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1			

Table E.D-2/Q.921 (sheet 5 of 10) – State transition table: receiving SREJ supervisory frame with correct format

BASIC STATE	MULTIPLE FRAME ESTABLISHED									
	Normal	Normal	Normal	Normal	Peer rec Busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
	Normal	SREJ recovery	Own rec busy	Normal	Normal	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7		
SREJ resp F = 0 N(R) = V(S) (Note)	Stop T200 Restart T203				Stop T200 Start T203	Stop T200 Start T203	Stop T200 Start T203	Stop T200 Start T203	Stop T200 Start T203	Stop T200 Start T203
SREJ resp F = 1 N(R) = V(S) (Note)	MDL-ERR ind(A) V(A) = N(R) Stop T200 Restart T203				MDL-ERR ind(A) V(A) = N(R) Stop T200 Start T203	MDL-ERR ind(A) V(A) = N(R) Stop T200 Start T203	MDL-ERR ind(A) V(A) = N(R) Stop T200 Start T203	MDL-ERR ind(A) V(A) = N(R) Stop T200 Start T203	MDL-ERR ind(A) V(A) = N(R) Stop T200 Start T203	MDL-ERR ind(A) V(A) = N(R) Stop T200 Start T203
SREJ resp F = 0 V(A) ≤ N(R) < V(S)	REQUESTED I FRAMES INTO RE-TX QUEUE Stop T200 Restart T203				REQUESTED I FRAMES INTO RE-TX QUEUE Stop T200 Start T203	REQUESTED I FRAMES INTO RE-TX QUEUE Stop T200 Start T203	REQUESTED I FRAMES INTO RE-TX QUEUE Stop T200 Start T203	REQUESTED I FRAMES INTO RE-TX QUEUE Stop T200 Start T203	REQUESTED I FRAMES INTO RE-TX QUEUE Stop T200 Start T203	REQUESTED I FRAMES INTO RE-TX QUEUE Stop T200 Start T203
SREJ resp F = 1 V(A) ≤ N(R) < V(S)	MDL-ERR ind(A) REQUESTED I FRAMES WITH V _i (P) < V(P) INTO RE-TX QUEUE V(A) = N(R) Stop T200 Restart T203				MDL-ERR ind(A) REQUESTED I FRAMES WITH V _i (P) < V(P) INTO RE-TX QUEUE V(A) = N(R) Stop T200 Start T203	MDL-ERR ind(A) REQUESTED I FRAMES WITH V _i (P) < V(P) INTO RE-TX QUEUE V(A) = N(R) Stop T200 Start T203	MDL-ERR ind(A) REQUESTED I FRAMES WITH V _i (P) < V(P) INTO RE-TX QUEUE V(A) = N(R) Stop T200 Start T203	MDL-ERR ind(A) REQUESTED I FRAMES WITH V _i (P) < V(P) INTO RE-TX QUEUE V(A) = N(R) Stop T200 Start T203	MDL-ERR ind(A) REQUESTED I FRAMES WITH V _i (P) < V(P) INTO RE-TX QUEUE V(A) = N(R) Stop T200 Start T203	MDL-ERR ind(A) REQUESTED I FRAMES WITH V _i (P) < V(P) INTO RE-TX QUEUE V(A) = N(R) Stop T200 Start T203
SREJ resp F = 0 N(R) error	MDL-ERR ind(J) RC = 0 TX SMP = 1 Stop T203 Restart T200									
	5.1									

Table E.D-2/Q.921 (sheet 5 of 10) – State transition table: receiving SREJ supervisory frame with correct format (concluded)

BASIC STATE		MULTIPLE FRAME ESTABLISHED									
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Peer rec Busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	Normal	SREJ recovery	Own rec busy	SREJ recovery	Own rec busy	
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7	7.5	7.6	
SREJ resp F = 1 N(R) error	MDL-ERR ind(A) MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	7.1	7.2	7.3	7.4	7.5	7.6	7.7	7.5	7.6	
					MDL-ERR ind(A) MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1						

NOTE – This event is impossible by the definition of the peer-to-peer data link procedures. However, it would not harm the information transfer, if actions according to this table are taken.

Table E.D-2/Q.921 (sheet 6 of 10) – State transition table: receiving RNR supervisory frame with correct format

BASIC STATE		MULTIPLE FRAME ESTABLISHED									
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	Peer rec busy	SREJ recovery	Own rec busy	SREJ and own rec busy	Peer rec busy	Peer rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.4	7.5	7.6	7.7	7.5	7.6
RNR cmd P = 1 N(R) = V(S)	TX RR F = 1 Stop T203 Restart T200 V(A) = N(R)	TX SREJ F = 1 Stop T203 Restart T200 V(A) = N(R)	TX RNR F = 1 Stop T203 Restart T200 V(A) = N(R)	TX RNR F = 1 Stop T203 Restart T200 V(A) = N(R)	TX RR F = 1 Restart T200 V(A) = N(R)	TX SREJ F = 1 Restart T200 V(A) = N(R)	TX SREJ F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)	TX SREJ F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)
RNR cmd P = 0 N(R) = V(S)	Stop T203 Restart T200 V(A) = N(R)	Stop T203 Restart T200 V(A) = N(R)	Stop T203 Restart T200 V(A) = N(R)	Stop T203 Restart T200 V(A) = N(R)	Stop T203 Restart T200 V(A) = N(R)	Restart T200 V(A) = N(R)					
RNR resp F = 0 N(R) = V(S)											
RNR resp F = 1 N(R) = V(S)	MDL-ERR ind(A) Stop T203 Restart T200 V(A) = N(R)	MDL-ERR ind(A) Stop T203 Restart T200 V(A) = N(R)	MDL-ERR ind(A) Stop T203 Restart T200 V(A) = N(R)	MDL-ERR ind(A) Stop T203 Restart T200 V(A) = N(R)	MDL-ERR ind(A) Stop T203 Restart T200 V(A) = N(R)	MDL-ERR ind(A) Restart T200 V(A) = N(R)					
RNR cmd P = 1 V(A) ≤ N(R) < V(S)	TX RR F = 1 Restart T200 V(A) = N(R)	TX SREJ F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)	TX RR F = 1 Restart T200 V(A) = N(R)	TX SREJ F = 1 Restart T200 V(A) = N(R)	TX SREJ F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)	TX SREJ F = 1 Restart T200 V(A) = N(R)	TX RNR F = 1 Restart T200 V(A) = N(R)
RNR cmd P = 0 V(A) ≤ N(R) < V(S)	Restart T200 V(A) = N(R)	Restart T200 V(A) = N(R)	Restart T200 V(A) = N(R)	Restart T200 V(A) = N(R)	Restart T200 V(A) = N(R)	Restart T200 V(A) = N(R)					
RNR resp F = 0 V(A) ≤ N(R) < V(S)											

Table E.D-2/Q.921 (sheet 6 of 10) – State transition table: receiving RNR supervisory frame with correct format (concluded)

BASIC STATE	MULTIPLE FRAME ESTABLISHED									
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	SREJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.4	7.5	7.6	7.7	7.7
RNR resp F = 1 V(A) ≤ N(R) < V(S)	MDL-ERR ind(A) Restart T200 V(A) = N(R) 7.4	MDL-ERR ind(A) Restart T200 V(A) = N(R) 7.5	MDL-ERR ind(A) Restart T200 V(A) = N(R) 7.6	MDL-ERR ind(A) Restart T200 V(A) = N(R) 7.7	MDL-ERR ind(A) Restart T200 V(A) = N(R)	MDL-ERR ind(A) Restart T200 V(A) = N(R)				
RNR cmd P = 1 N(R) error	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1		TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1		
RNR cmd P = 0 N(R) error	MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1				MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1				
RNR resp F = 0 N(R) error										
RNR resp F = 1 N(R) error	MDL-ERR ind(A) MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1				MDL-ERR ind(A) MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	MDL-ERR ind(A) MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1				

Table E.D-2/Q.921 (sheet 7 of 10) – State transition table: receiving I command frame with correct format acknowledgedging all outstanding I frames or containing an N(R) which satisfies $V(A) < N(R) < V(S)$

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
I cmd P = 1 N(S) = V(R) N(R) = V(S) NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 Stop T200 Restart T203 V(A) = N(R)	UPDATE V(R) DL-DATA Inds TX RR F = 1 Stop T200 Restart T203 V(A) = N(R)	"DISCARD" TX RNR F = 1 Stop T200 Restart T203 V(A) = N(R)	SREJ and own rec busy	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	SREJ and own rec busy
I cmd P = 1 N(S) = V(R) N(R) = V(S) OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 Stop T200 Restart T203 V(A) = N(R)	UPDATE V(R) DL-DATA Inds TX SREJ F = 1 Stop T200 Restart T203 V(A) = N(R)	"DISCARD" TX RNR F = 1 Stop T200 Restart T203 V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX SREJ F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	
I cmd P = 0 N(S) = V(R) N(R) = V(S) NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX ACK Stop T200 Restart T203 V(A) = N(R)	UPDATE V(R) DL-DATA Inds TX ACK Stop T200 Restart T203 V(A) = N(R)	"DISCARD" Stop T200 Restart T203 V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX RR F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)	

**Table E.D-2/Q.921 (sheet 7 of 10) – State transition table: receiving I command frame with correct format acknowledged
all outstanding I frames or containing an N(R) which satisfies $V(A) < N(R) < V(S)$ (continued)**

BASIC STATE	MULTIPLE FRAME ESTABLISHED						
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6
I cmd P = 0 N(S) = V(R) N(R) = V(S) OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX ACK Stop T200 Restart T203 V(A) = N(R)	UPDATE V(R) DL-DATA Inds TX SREJ F = 0 Stop T200 Restart T203 V(A) = N(R)	"DISCARD" Stop T200 Restart T203 V(A) = N(R)	SREJ and own rec busy	V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX SREJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)
I cmd P = 1 N(S) ≠ V(R) N(R) = V(S)	I FRAME INTO REC QUEUE TX SREJ F = 1 Stop T200 Restart T203 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 1 Stop T200 Restart T203 V(A) = N(R)	"DISCARD" TX RNR F = 1 Stop T200 Restart T203 V(A) = N(R)	SREJ and own rec busy	I FRAME INTO REC QUEUE TX SREJ F = 1 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)
I cmd P = 0 N(S) ≠ V(R) N(R) = V(S)	I FRAME INTO REC QUEUE TX SREJ F = 0 Stop T200 Restart T203 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 0 Stop T200 Restart T203 V(A) = N(R)	"DISCARD" Stop T200 Restart T203 V(A) = N(R)		I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)
I cmd P = 1 N(S) = V(R) V(A) < N(R) < V(S) NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 Restart T200 V(A) = N(R)	UPDATE V(R) DL-DATA Inds TX RR F = 1 Restart T200 V(A) = N(R)	"DISCARD" TX RNR F = 1 Restart T200 V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)
I cmd P = 1 N(S) = V(R) V(A) < N(R) < V(S) OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 Restart T200 V(A) = N(R)	UPDATE V(R) DL-DATA Inds TX SREJ F = 1 Restart T200 V(A) = N(R)	"DISCARD" TX RNR F = 1 Restart T200 V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX SREJ F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)

**Table E.D-2/Q.921 (sheet 7 of 10) – State transition table: receiving I command frame with correct format acknowledged
all outstanding I frames or containing an N(R) which satisfies $V(A) < N(R) < V(S)$ (concluded)**

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
I cmd P = 0 N(S) = V(R) V(A) < N(R) < V(S) NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX ACK Restart T200 V(A) = N(R)	UPDATE V(R) DL-DATA Inds TX ACK Restart T200 V(A) = N(R) 7.0	"DISCARD" Restart T200 V(A) = N(R)	SREJ and own rec busy	V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX RR F = 0 V(A) = N(R) 7.4	"DISCARD" V(A) = N(R)	Peer rec busy
I cmd P = 0 N(S) = V(R) V(A) < N(R) < V(S) OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX ACK Restart T200 V(A) = N(R)	UPDATE V(R) DL-DATA Inds TX SREJ F = 0 Restart T200 V(A) = N(R)	"DISCARD" Restart T200 V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX SREJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)	
I cmd P = 1 N(S) ≠ V(R) V(A) < N(R) < V(S)	I FRAME INTO REC QUEUE TX SREJ F = 1 Restart T200 V(A) = N(R) 7.1	I FRAME INTO REC QUEUE TX SREJ F = 1 Restart T200 V(A) = N(R)	"DISCARD" TX RNR F = 1 Restart T200 V(A) = N(R)		I FRAME INTO REC QUEUE TX SREJ F = 1 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	
I cmd P = 0 N(S) ≠ V(R) V(A) < N(R) < V(S)	I FRAME INTO REC QUEUE TX SREJ F = 0 Restart T200 V(A) = N(R) 7.1	I FRAME INTO REC QUEUE TX SREJ F = 0 Restart T200 V(A) = N(R)	"DISCARD" Restart T203 V(A) = N(R)		I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)	

Table E.D-2/Q.921 (sheet 8 of 10) – State transition table: receiving I command frame with correct format containing an N(R) which satisfies $V(A) < N(R) < V(S)$, or an N(R) error

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	Normal	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
I cmd P = 1 N(S) = V(R) V(A) = N(R) < V(S) NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1	UPDATE V(R) DL-DATA Inds TX RR F = 1	"DISCARD" TX RNR F = 1	SREJ and own rec busy	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1	UPDATE V(R) DL-DATA inds TX RR F = 1	"DISCARD" TX RNR F = 1	SREJ and own rec busy
I cmd P = 1 N(S) = V(R) V(A) = N(R) < V(S) OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1	UPDATE V(R) DL-DATA Inds TX SREJ F = 1	"DISCARD" TX RNR F = 1		V(R) = V(R) + 1 DL-DATA ind TX RR F = 1	UPDATE V(R) DL-DATA inds TX SREJ F = 1	"DISCARD" TX RNR F = 1	
I cmd P = 0 N(S) = V(R) V(A) = N(R) < V(S) NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX ACK	UPDATE V(R) DL-DATA Inds TX ACK	"DISCARD"		V(R) = V(R) + 1 DL-DATA ind TX RR F = 0	UPDATE V(R) DL-DATA inds TX RR F = 0	"DISCARD"	
I cmd P = 0 N(S) = V(R) V(A) = N(R) < V(S) OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX ACK	UPDATE V(R) DL-DATA Inds TX SREJ F = 0	"DISCARD"		V(R) = V(R) + 1 DL-DATA ind TX RR F = 0	UPDATE V(R) DL-DATA inds TX SREJ F = 0	"DISCARD"	
I cmd P = 1 N(S) ≠ V(R) V(A) = N(R) < V(S)	I FRAME INTO REC QUEUE TX SREJ F = 1	I FRAME INTO REC QUEUE TX SREJ F = 1	"DISCARD" TX RNR F = 1		I FRAME INTO REC QUEUE TX SREJ F = 1	I FRAME INTO REC QUEUE TX SREJ F = 1	"DISCARD" TX RNR F = 1	
I cmd P = 0 N(S) ≠ V(R) V(A) = N(R) < V(S)	I FRAME INTO REC QUEUE TX SREJ F = 0	I FRAME INTO REC QUEUE TX SREJ F = 0	"DISCARD"		I FRAME INTO REC QUEUE TX SREJ F = 0	I FRAME INTO REC QUEUE TX SREJ F = 0	"DISCARD"	

Table E.D-2/Q.921 (sheet 8 of 10) – State transition table: receiving I command frame with correct format containing an N(R) which satisfies $V(A) < N(R) < V(S)$, or an N(R) error (continued)

BASIC STATE	MULTIPLE FRAME ESTABLISHED						
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6
I cmd P = 1 N(S) = V(R) N(R) error NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	UPDATE V(R) DL-DATA Inds TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	7.3	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	UPDATE V(R) DL-DATA inds TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1
I cmd P = 1 N(S) = V(R) N(R) error OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	UPDATE V(R) DL-DATA Inds TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	7.3	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	UPDATE V(R) DL-DATA inds TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1
I cmd P = 0 N(S) = V(R) N(R) error NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	UPDATE V(R) DL-DATA Inds MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	7.3	V(R) = V(R) + 1 DL-DATA ind MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	UPDATE V(R) DL-DATA inds MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1
I cmd P = 0 N(S) = V(R) N(R) error OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	UPDATE V(R) DL-DATA Inds MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1	7.3	V(R) = V(R) + 1 DL-DATA ind MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	UPDATE V(R) DL-DATA inds MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1

Table E.D-2/Q.921 (sheet 8 of 10) – State transition table: receiving I command frame with correct format containing an N(R) which satisfies $V(A) < N(R) < V(S)$, or an N(R) error (concluded)

BASIC STATE	MULTIPLE FRAME ESTABLISHED									
	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Normal	Normal	Normal	Normal	Normal
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	SREJ recovery	Normal	SREJ recovery	Own rec busy	SREJ recovery	SREJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7	7.5	7.6
I cmd P = 1 N(S)≠V(R) N(R) error	I FRAME INTO REC QUEUE TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Stop T203 Restart T200 5.1	I FRAME INTO REC QUEUE TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Stop T203 Restart T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Stop T203 Restart T200 5.1	I FRAME INTO REC QUEUE TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1	I FRAME INTO REC QUEUE TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1	I FRAME INTO REC QUEUE TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1	I FRAME INTO REC QUEUE TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1	I FRAME INTO REC QUEUE TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1
I cmd P = 0 N(S)≠V(R) N(R) error	I FRAME INTO REC QUEUE TX SREJ F = 0 MDL-ERR ind(J) RC = 0 TX SMP = 1 Stop T203 Restart T200 5.1	I FRAME INTO REC QUEUE TX SREJ F = 0 MDL-ERR ind(J) RC = 0 TX SMP = 1 Stop T203 Restart T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SMP = 1 Stop T203 Restart T200 5.1	I FRAME INTO REC QUEUE TX SREJ F = 0 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1	I FRAME INTO REC QUEUE TX SREJ F = 0 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1	I FRAME INTO REC QUEUE TX SREJ F = 0 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1	I FRAME INTO REC QUEUE TX SREJ F = 0 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1	I FRAME INTO REC QUEUE TX SREJ F = 0 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200 5.1

Table E.D-2/Q.921 (sheet 9 of 10) – State transition table: internal events (expiry of timers, receiver busy condition, I frame in re-tx queue)

BASIC STATE	MULTIPLE FRAME ESTABLISHED									
	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Normal	Normal	Normal	Normal	Normal
RECEIVER CONDITION	Normal	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	SREJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7	7.8	7.9
T200 TIME-OUT RC < N200	RC = 0 TX RR P = 1 RC = RC + 1 Start T200 8.0	RC = 0 TX RR P = 1 RC = RC + 1 Start T200 8.1	RC = 0 TX RR P = 1 RC = RC + 1 Start T200 8.2	RC = 0 TX RNR P = 1 RC = RC + 1 Start T200 8.3	RC = 0 TX RR P = 1 RC = RC + 1 Start T200 8.4	RC = 0 TX RR P = 1 RC = RC + 1 Start T200 8.5	RC = 0 TX RNR P = 1 RC = RC + 1 Start T200 8.6	RC = 0 TX RR P = 1 RC = RC + 1 Start T200 8.7	RC = 0 TX RNR P = 1 RC = RC + 1 Start T200 8.8	RC = 0 TX RR P = 1 RC = RC + 1 Start T200 8.9
T200 TIME-OUT RC = N200	/	/	/	/	/	/	/	/	/	/
T203 TIME-OUT	RC = 0 TX RR P = 1 Start T200 8.0	RC = 0 TX RR P = 1 Start T200 8.1	RC = 0 TX RNR P = 1 Start T200 8.2	RC = 0 TX RNR P = 1 Start T200 8.3	RC = 0 TX RR P = 1 Start T200 8.4	RC = 0 TX RR P = 1 Start T200 8.5	RC = 0 TX RR P = 1 Start T200 8.6	RC = 0 TX RR P = 1 Start T200 8.7	RC = 0 TX RR P = 1 Start T200 8.8	RC = 0 TX RR P = 1 Start T200 8.9
SET OWN RECEIVER BUSY (Note)	TX RNR F = 0 7.2	TX RNR F = 0 7.3	-	-	TX RNR F = 0 7.6	TX RNR F = 0 7.7	-	-	-	-
CLEAR OWN RECEIVER BUSY (Note)	-	-	TX RR F = 0 7.0	TX RR F = 0 7.1	-	-	TX RR F = 0 7.4	TX RR F = 0 7.5	TX RR F = 0 7.6	TX RR F = 0 7.7
I FRAME IN RE-TX QUEUE (Note) NOT LAST I FRAME TO BE RE-TX	RE-TX I P = 0 V ₃ (P) = V(P) Stop T203 TIMER T200				LEAVE I FRAME IN RE-TX QUEUE					

Table E.D-2/Q.921 (sheet 9 of 10) – State transition table: internal events (expiry of timers, receiver busy condition, I frame in re-tx queue) (concluded)

BASIC STATE	MULTIPLE FRAME ESTABLISHED							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	7.0	7.1	7.2	7.3	7.4	7.5	7.6	7.7
I FRAME IN RE-TX QUEUE (Note) LAST I FRAME TO BE RE-TX	8.0 Either RE-TX IP = 1 or RE-TX IP = 0 TX RRP = 1 then $V(P) = V(P)$ $V(P) = V(P) + 1$ Stop T203 Start T200	8.1 Either RE-TX IP = 1 or RE-TX IP = 0 TX RRP = 1 then $V(P) = V(P)$ $V(P) = V(P) + 1$ Stop T203 Start T200	8.2 Either RE-TX IP = 1 or RE-TX IP = 0 TX RRP = 1 then $V(P) = V(P)$ $V(P) = V(P) + 1$ Stop T203 Start T200	8.3 Either RE-TX IP = 1 or RE-TX IP = 0 TX RRP = 1 then $V(P) = V(P)$ $V(P) = V(P) + 1$ Stop T203 Start T200	7.4 Normal	7.5 SREJ recovery	7.6 Own rec busy	7.7 SREJ and own rec busy

NOTE – These signals are generated outside the procedures specified in this state transition table, and may be generated by the connection management entity.

Table E.D-3/Q.921 (sheet 1 of 10) – State transition table: receiving primitive

BASIC STATE	TIMER RECOVERY							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
DL-ESTABLISH request	DISC I, RE-TX and REC QUEUES RC = 0 TX SMP = 1 Restart T200 5.0							
DL-RELEASE request	DISC I, RE-TX and REC QUEUES RC = 0 TX DISC P = 1 Restart T200 6							
DL-DATA request	DATA INTO I QUEUE							
I FRAME IN QUEUE $V(S) < V(A) + k$ RE-TX QUEUE EMPTY	TX I P = 0 $V(S) = V(S) + 1$ $V_d(P) = V(P)$				Normal	LEAVE I FRAME IN QUEUE		
I FRAME IN QUEUE $V(S) < V(A) + k$ RE-TX QUEUE NOT EMPTY	LEAVE I FRAME IN QUEUE							
I FRAME IN QUEUE $V(S) = V(A) + k$								
DL-UNIT DATA request	UNIT DATA INTO UI QUEUE							
UI FRAME IN QUEUE	TX UIP = 0							
MDL-ASSIGN request								

Table E.D-3/Q.921 (sheet 1 of 10) – State transition table: receiving primitive (concluded)

BASIC STATE	TIMER RECOVERY						
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6
MDL-REMOVE request	DL-REL ind DISC I, RE-TX, REC and UI QUEUES Stop T200 Stop T203 1						
MDL-ERROR response							
PERSISTENT DEACTIVATION	DL-REL ind DISC I, RE-TX, REC and UI QUEUES Stop T200 Stop T203 4						
							8.7

Table E.D-3/Q.921 (sheet 4 of 10) – State transition table: receiving RR supervisory frame with correct format, clearance of timer recovery if there is F = 1 only

BASIC STATE	TIMER RECOVERY									
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7	8.8	8.9
RR cmd P = 1 V(A) ≤ N(R) ≤ V(S)	TX RR F = 1 V(A) = N(R)	TX SREJ F = 1 V(A) = N(R)	TX RNR F = 1 V(A) = N(R)	TX RR F = 1 V(A) = N(R)	TX RR F = 1 V(A) = N(R)	TX SREJ F = 1 V(A) = N(R)	TX RNR F = 1 V(A) = N(R)	TX RNR F = 1 V(A) = N(R)	TX RNR F = 1 V(A) = N(R)	TX RNR F = 1 V(A) = N(R)
RR cmd P = 0 V(A) ≤ N(R) ≤ V(S)	V(A) = N(R)				V(A) = N(R)		V(A) = N(R)		V(A) = N(R)	V(A) = N(R)
RR resp F = 0 V(A) ≤ N(R) ≤ V(S)										
RR resp F = 1 N(R) = V(S)	Stop T200 Restart T203 V(A) = N(R)	Stop T200 Restart T203 V(A) = N(R)	Stop T200 Restart T203 V(A) = N(R)	Stop T200 Restart T203 V(A) = N(R)	Stop T200 Restart T203 V(A) = N(R)	Stop T200 Restart T203 V(A) = N(R)	Stop T200 Restart T203 V(A) = N(R)	Stop T200 Restart T203 V(A) = N(R)	Stop T200 Restart T203 V(A) = N(R)	Stop T200 Restart T203 V(A) = N(R)
RR resp F = 1 V(A) ≤ N(R) < V(S)	UNACK I FRAMES WITH V ₃ (P) < V(P) INTO RE-TX QUEUE Restart T200 V(A) = N(R)	UNACK I FRAMES WITH V ₃ (P) < V(P) INTO RE-TX QUEUE Restart T200 V(A) = N(R)	UNACK I FRAMES WITH V ₃ (P) < V(P) INTO RE-TX QUEUE Restart T200 V(A) = N(R)	UNACK I FRAMES WITH V ₃ (P) < V(P) INTO RE-TX QUEUE Restart T200 V(A) = N(R)	UNACK I FRAMES WITH V ₃ (P) < V(P) INTO RE-TX QUEUE Restart T200 V(A) = N(R)	UNACK I FRAMES WITH V ₃ (P) < V(P) INTO RE-TX QUEUE Restart T200 V(A) = N(R)	UNACK I FRAMES WITH V ₃ (P) < V(P) INTO RE-TX QUEUE Restart T200 V(A) = N(R)	UNACK I FRAMES WITH V ₃ (P) < V(P) INTO RE-TX QUEUE Restart T200 V(A) = N(R)	UNACK I FRAMES WITH V ₃ (P) < V(P) INTO RE-TX QUEUE Restart T200 V(A) = N(R)	UNACK I FRAMES WITH V ₃ (P) < V(P) INTO RE-TX QUEUE Restart T200 V(A) = N(R)
RR cmd P = 1 N(R) error	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200	TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200	TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SMP = 1 Restart T200

Table E.D-3/Q.921 (sheet 4 of 10) – State transition table: receiving RR supervisory frame with correct format, clearance of timer recovery if there is F = 1 only (concluded)

BASIC STATE	TIMER RECOVERY						
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6
RR cmd P = 0 N(R) error	MDL-ERR ind(J) RC = 0 TX SM P = 1 Stop T203 Restart T200 5.1				MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1		
RR resp F = 0 N(R) error							
RR resp F = 1 N(R) error							

Table E.D-3/Q.921 (sheet 5 of 10) – State transition table: receiving SREJ supervisory frame with correct format, clearance of timer recovery if there is F = 1 only

BASIC STATE	TIMER RECOVERY									
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	Own rec busy	SREJ and own rec busy	SREJ recovery	Normal	SREJ recovery	Own rec busy	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	8.0	8.2	8.3	8.1	8.4	8.5	8.6	8.5	8.6	8.7
SREJ resp F = 0 $V(A) \leq N(R) \leq V(S)$	REQUESTED I FRAMES INTO RE-TX QUEUE				REQUESTED I FRAMES INTO RE-TX QUEUE 8.0	REQUESTED I FRAMES INTO RE-TX QUEUE 8.1	REQUESTED I FRAMES INTO RE-TX QUEUE 8.2	REQUESTED I FRAMES INTO RE-TX QUEUE 8.3	REQUESTED I FRAMES INTO RE-TX QUEUE 8.4	REQUESTED I FRAMES INTO RE-TX QUEUE 8.5
SREJ resp F = 1 $V(A) \leq N(R) \leq V(S)$	REQUESTED I FRAMES WITH $V_3(P) < V(P)$ INTO RE-TX QUEUE Stop T200 Start T203 $V(A) = N(R)$	REQUESTED I FRAMES WITH $V_4(P) < V(P)$ INTO RE-TX QUEUE Stop T200 Start T203 $V(A) = N(R)$	REQUESTED I FRAMES WITH $V_5(P) < V(P)$ INTO RE-TX QUEUE Stop T200 Start T203 $V(A) = N(R)$	REQUESTED I FRAMES WITH $V_6(P) < V(P)$ INTO RE-TX QUEUE Stop T200 Start T203 $V(A) = N(R)$	REQUESTED I FRAMES WITH $V_7(P) < V(P)$ INTO RE-TX QUEUE Stop T200 Start T203 $V(A) = N(R)$	REQUESTED I FRAMES WITH $V_8(P) < V(P)$ INTO RE-TX QUEUE Stop T200 Start T203 $V(A) = N(R)$	REQUESTED I FRAMES WITH $V_9(P) < V(P)$ INTO RE-TX QUEUE Stop T200 Start T203 $V(A) = N(R)$	REQUESTED I FRAMES WITH $V_{10}(P) < V(P)$ INTO RE-TX QUEUE Stop T200 Start T203 $V(A) = N(R)$	REQUESTED I FRAMES WITH $V_{11}(P) < V(P)$ INTO RE-TX QUEUE Stop T200 Start T203 $V(A) = N(R)$	REQUESTED I FRAMES INTO RE-TX QUEUE 8.3
SREJ resp F = 0 N(R) error	MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1									
SREJ resp F = 1 N(R) error										

Table E.D-3/Q.921 (sheet 6 of 10) – State transition table: receiving RNR supervisory frame with correct format, clearance of timer recovery if there is F = 1 only

BASIC STATE	TIMER RECOVERY									
	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Normal	Normal	Normal	Normal	Normal
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.5	8.6	8.7
RNR cmd P = 1 V(A) ≤ N(R) ≤ V(S)	TX RR F = 1 V(A) = N(R) 8.4	TX SREJ F = 1 V(A) = N(R) 8.5	TX RNR F = 1 V(A) = N(R) 8.6	TX RNR F = 1 V(A) = N(R) 8.7	TX RR F = 1 V(A) = N(R) 8.7	TX SREJ F = 1 V(A) = N(R) 8.7	TX RNR F = 1 V(A) = N(R) 8.7	TX SREJ F = 1 V(A) = N(R) 8.7	TX RNR F = 1 V(A) = N(R) 8.7	TX RNR F = 1 V(A) = N(R) 8.7
RNR cmd P = 0 V(A) ≤ N(R) ≤ V(S)	V(A) = N(R) 8.4	V(A) = N(R) 8.5	V(A) = N(R) 8.6	V(A) = N(R) 8.7	V(A) = N(R) 8.7	V(A) = N(R) 8.7	V(A) = N(R) 8.7	V(A) = N(R) 8.7	V(A) = N(R) 8.7	V(A) = N(R) 8.7
RNR resp F = 0 V(A) ≤ N(R) ≤ V(S)										
RNR resp F = 1 V(A) ≤ N(R) ≤ V(S)	Restart T200 V(A) = N(R) 7.4	Restart T200 V(A) = N(R) 7.5	Restart T200 V(A) = N(R) 7.6	Restart T200 V(A) = N(R) 7.7	Restart T200 V(A) = N(R) 7.7	Restart T200 V(A) = N(R) 7.7	Restart T200 V(A) = N(R) 7.6	Restart T200 V(A) = N(R) 7.5	Restart T200 V(A) = N(R) 7.6	Restart T200 V(A) = N(R) 7.7
RNR cmd P = 1 N(R) error	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1
RNR cmd P = 0 N(R) error	MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1									
RNR resp F = 0 N(R) error										
RNR resp F = 1 N(R) error										

Table E.D-3/Q.921 (sheet 7 of 10) – State transition table: receiving I command frame with correct format acknowledging all outstanding I frames or containing an N(R) which satisfies $V(A) < N(R) < V(S)$; no clearance of timer recovery

BASIC STATE	TIMER RECOVERY							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	Own rec busy	SREJ and own rec busy	Normal	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
I cmd P = 1 N(S) = V(R) N(R) = V(S) NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	UPDATE V(R) DL-DATA Inds TX RR F = 1 V(A) = N(R) 8.0	"DISCARD" TX RNR F = 1 V(A) = N(R)	SREJ and own rec busy	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX RR F = 1 V(A) = N(R) 8.4	"DISCARD" TX RNR F = 1 V(A) = N(R)	SREJ and own rec busy
I cmd P = 1 N(S) = V(R) N(R) = V(S) OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	UPDATE V(R) DL-DATA Inds TX SREJ F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX SREJ F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	
I cmd P = 0 N(S) = V(R) N(R) = V(S) NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX ACK V(A) = N(R)	UPDATE V(R) DL-DATA Inds TX ACK V(A) = N(R) 8.0	"DISCARD" V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX RR F = 0 V(A) = N(R) 8.4	"DISCARD" V(A) = N(R)	
I cmd P = 0 N(S) = V(R) N(R) = V(S) OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX ACK V(A) = N(R)	UPDATE V(R) DL-DATA Inds TX SREJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)		V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX SREJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)	
I cmd P = 1 N(S) ≠ V(R) N(R) = V(S)	I FRAME INTO REC QUEUE TX SREJ F = 1 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)		I FRAME INTO REC QUEUE TX SREJ F = 1 V(A) = N(R) 8.5	I FRAME INTO REC QUEUE TX SREJ F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	

Table E.D-3/Q.921 (sheet 7 of 10) – State transition table: receiving I command frame with correct format acknowledging all outstanding I frames or containing an N(R) which satisfies $V(A) < N(R) < V(S)$; no clearance of timer recovery (continued)

BASIC STATE	TIMER RECOVERY							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	Own rec busy	SREJ and own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
I cmd P = 0 N(S) ≠ V(R) N(R) = V(S)	I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)	"DISCARD" V(A) = N(R)
I cmd P = 1 N(S) = V(R) V(A) < N(R) < V(S) NO OUT-OF-SEQ FRAMES IN REC QUEUE	8.1 V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	8.0 UPDATE V(R) DL-DATA Inds TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX RR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)
I cmd P = 1 N(S) = V(R) V(A) < N(R) < V(S) OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	8.0 UPDATE V(R) DL-DATA Inds TX SREJ F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX SREJ F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)
I cmd P = 0 N(S) = V(R) V(A) < N(R) < V(S) NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX ACK V(A) = N(R)	8.0 UPDATE V(R) DL-DATA Inds TX ACK V(A) = N(R)	"DISCARD" V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX RR F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)	"DISCARD" V(A) = N(R)
I cmd P = 0 N(S) = V(R) V(A) < N(R) < V(S) OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX ACK V(A) = N(R)	8.0 UPDATE V(R) DL-DATA Inds TX SREJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	V(R) = V(R) + 1 DL-DATA ind TX RR F = 0 V(A) = N(R)	UPDATE V(R) DL-DATA inds TX SREJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)	"DISCARD" V(A) = N(R)

Table E.D-3/Q.921 (sheet 7 of 10) – State transition table: receiving I command frame with correct format acknowledging all outstanding I frames or containing an N(R) which satisfies $V(A) < N(R) < V(S)$; no clearance of timer recovery (concluded)

BASIC STATE	TIMER RECOVERY							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
I cmd P = 1 N(S) ≠ V(R) V(A) < N(R) < V(S)	I FRAME INTO REC QUEUE TX SREJ F = 1 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 1 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 1 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 1 V(A) = N(R)	"DISCARD" TX RNR F = 1 V(A) = N(R)	
I cmd P = 0 N(S) ≠ V(R) V(A) < N(R) < V(S)	I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	I FRAME INTO REC QUEUE TX SREJ F = 0 V(A) = N(R)	"DISCARD" V(A) = N(R)	

Table E.D-3/Q.921 (sheet 8 of 10) – State transition table: receiving I command frame with correct format containing an N(R) which satisfies $V(A) < N(R) < V(S)$, or an N(R) error

BASIC STATE	TIMER RECOVERY									
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	Own rec busy	SREJ and own rec busy	Normal	Normal	SREJ recovery	Own rec busy	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.5	8.6	8.7
I cmd P = 1 N(S) = V(R) V(A) = N(R) < V(S) NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RRF = 1	UPDATE V(R) DL-DATA Inds TX RRF = 1	"DISCARD" TX RNR F = 1	SREJ and own rec busy	V(R) = V(R) + 1 DL-DATA ind TX RRF = 1	UPDATE V(R) DL-DATA inds TX RRF = 1	"DISCARD" TX RNR F = 1	UPDATE V(R) DL-DATA inds TX RRF = 1	"DISCARD" TX RNR F = 1	SREJ and own rec busy
I cmd P = 1 N(S) = V(R) V(A) = N(R) < V(S) OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RRF = 1	UPDATE V(R) DL-DATA Inds TX SREJ F = 1	"DISCARD" TX RNR F = 1		V(R) = V(R) + 1 DL-DATA ind TX RRF = 1	UPDATE V(R) DL-DATA inds TX SREJ F = 1	"DISCARD" TX RNR F = 1	UPDATE V(R) DL-DATA inds TX SREJ F = 1	"DISCARD" TX RNR F = 1	
I cmd P = 0 N(S) = V(R) V(A) = N(R) < V(S) NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX ACK	UPDATE V(R) DL-DATA Inds TX ACK	"DISCARD"		V(R) = V(R) + 1 DL-DATA ind TX RRF = 0	UPDATE V(R) DL-DATA inds TX RRF = 0	"DISCARD"	UPDATE V(R) DL-DATA inds TX RRF = 0	"DISCARD"	
I cmd P = 0 N(S) = V(R) V(A) = N(R) < V(S) OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX ACK	UPDATE V(R) DL-DATA Inds TX SREJ F = 0	"DISCARD"		V(R) = V(R) + 1 DL-DATA ind TX RRF = 0	UPDATE V(R) DL-DATA inds TX SREJ F = 0	"DISCARD"	UPDATE V(R) DL-DATA inds TX SREJ F = 0	"DISCARD"	
I cmd P = 1 N(S) ≠ V(R) V(A) = N(R) < V(S)	I FRAME INTO REC QUEUE TX SREJ F = 1	I FRAME INTO REC QUEUE TX SREJ F = 1	"DISCARD" TX RNR F = 1		I FRAME INTO REC QUEUE TX SREJ F = 1	I FRAME INTO REC QUEUE TX SREJ F = 1	"DISCARD" TX RNR F = 1	I FRAME INTO REC QUEUE TX SREJ F = 1	"DISCARD" TX RNR F = 1	
I cmd P = 0 N(S) ≠ V(R) V(A) = N(R) < V(S)	I FRAME INTO REC QUEUE TX SREJ F = 0	I FRAME INTO REC QUEUE TX SREJ F = 0	"DISCARD"		I FRAME INTO REC QUEUE TX SREJ F = 0	I FRAME INTO REC QUEUE TX SREJ F = 0	"DISCARD"	I FRAME INTO REC QUEUE TX SREJ F = 0	"DISCARD"	

Table E.D-3/Q.921 (sheet 8 of 10) – State transition table: receiving I command frame with correct format containing an N(R) which satisfies $V(A) < N(R) < V(S)$, or an N(R) error (continued)

BASIC STATE	TIMER RECOVERY						
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	Normal	Own rec busy	SREJ and own rec busy	Normal	Own rec busy	SREJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6
I cmd P = 1 N(S) = V(R) N(R) error NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	UPDATE V(R) DL-DATA Inds TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	8.3	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	UPDATE V(R) DL-DATA inds TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1
I cmd P = 1 N(S) = V(R) N(R) error OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	UPDATE V(R) DL-DATA Inds TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	8.3	V(R) = V(R) + 1 DL-DATA ind TX RR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	UPDATE V(R) DL-DATA inds TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1
I cmd P = 0 N(S) = V(R) N(R) error NO OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	UPDATE V(R) DL-DATA Inds MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	8.3	V(R) = V(R) + 1 DL-DATA ind MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	UPDATE V(R) DL-DATA inds MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1
I cmd P = 0 N(S) = V(R) N(R) error OUT-OF-SEQ FRAMES IN REC QUEUE	V(R) = V(R) + 1 DL-DATA ind MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	UPDATE V(R) DL-DATA Inds MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	8.3	V(R) = V(R) + 1 DL-DATA ind MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	UPDATE V(R) DL-DATA inds MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1	"DISCARD" MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200 5.1

Table E.D-3/Q.921 (sheet 8 of 10) – State transition table: receiving I command frame with correct format containing an N(R) which satisfies $V(A) < N(R) < V(S)$, or an N(R) error (continued)

BASIC STATE		TIMER RECOVERY						
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
I cmd P = 1 N(S)≠V(R) N(R) error	I FRAME INTO REC QUEUE TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200	I FRAME INTO REC QUEUE TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200	Normal SREJ and own rec busy	I FRAME INTO REC QUEUE TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200	I FRAME INTO REC QUEUE TX SREJ F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200	"DISCARD" TX RNR F = 1 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200	8.7
I cmd P = 0 N(S)≠V(R) N(R) error	I FRAME INTO REC QUEUE TX SREJ F = 0 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200	I FRAME INTO REC QUEUE TX SREJ F = 0 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200	"DISCARD" MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200	5.1	I FRAME INTO REC QUEUE TX SREJ F = 0 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200	I FRAME INTO REC QUEUE TX SREJ F = 0 MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200	"DISCARD" MDL-ERR ind(J) RC = 0 TX SM P = 1 Restart T200	5.1
T200 TIME-OUT RC < N200 V(A) < V(S)	TX RR P = 1 RC = RC + 1 Start T200	5.1	TX RNR P = 1 RC = RC + 1 Start T200	5.1	TX RR P = 1 RC = RC + 1 Start T200	5.1	TX RNR P = 1 RC = RC + 1 Start T200	5.1
T200 TIME-OUT RC < N200 V(A) = V(S)	MDL-ERR ind(I) RC = 0 TX SM P = 1 Start T200	5.1						

Table E.D-3/Q.921 (sheet 8 of 10) – State transition table: receiving I command frame with correct format containing an N(R) which satisfies $V(A) < N(R) < V(S)$, or an N(R) error (concluded)

BASIC STATE	TIMER RECOVERY							
	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.5	8.6	8.7
T203 TIME-OUT	/	/	/	/	/	/	/	/
SET OWN RECEIVER BUSY (Note)	TX RNR F = 0	TX RNR F = 0	-	-	TX RNR F = 0	TX RNR F = 0	-	-
CLEAR OWN RECEIVER BUSY (Note)	-	-	TX RR F = 0	TX RR F = 0	-	-	TX RR F = 0	TX RR F = 0
I FRAME IN RE-TX QUEUE (Note)	RE-TX I P = 0 $V_i(P) = V(P)$	8.3	8.0	8.1	LEAVE I FRAME IN RE-TX QUEUE	8.4	8.5	8.5

NOTE – These signals are generated outside the procedures specified in this state transition table, and may be generated by the connection management entity.

Table E.D-3/Q.921 (sheet 9 of 10) – State transition table: internal events (expiry of timers, receiver busy condition, I frame in re-tx queue); initiation of a re-establishment procedure if the value of the retransmission count variable is equal to N200

BASIC STATE	TIMER RECOVERY									
	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
TRANSMITTER CONDITION	Normal	Normal	Normal	Normal	Normal	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy	Peer rec busy
RECEIVER CONDITION	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	Normal	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	SREJ and own rec busy
STATE NUMBER	8.0	8.1	8.2	8.3	8.4	8.4	8.5	8.6	8.7	8.7
T200 TIME-OUT RC<N200 V(A)<V(S)	TX RR P=1 RC = RC + 1 Start T200	8.1	TX RNR P = 1 RC = RC + 1 Start T200	SREJ and own rec busy	TX RR P = 1 RC = RC + 1 Start T200	Normal	SREJ recovery	Own rec busy	SREJ and own rec busy	SREJ and own rec busy
T200 TIME-OUT RC<N200 V(A) = V(S)										
T200 TIME-OUT RC = N200	MDL-ERR ind(1) RC = 0 TX SM P = 1 Start T200	5.1								
T203 TIME-OUT	/	/	/	/	/	/	/	/	/	/
SET OWN RECEIVER BUSY (Note)	TX RNR P = 0	TX RNR F = 0			TX RNR F = 0	TX RNR F = 0	TX RNR F = 0			
CLEAR OWN RECEIVER BUSY (Note)	2.2	8.3	TX RR F = 0	8.1	8.6	8.6	8.7	TX RR F = 0	TX RR F = 0	TX RR F = 0
I FRAME IN RE-TX QUEUE (Note)	RE-TX I P = 0 V ₃ (P) = V(P)				LEAVE I FRAME IN RE-TX QUEUE					
NOTE – These signals are generated outside the procedures specified in this state transition table, and may be generated by the connection management entity.										

**Protocol Implementation Conformance Statement (PICS)
to Recommendation Q.921 for Basic Rate (User-side)**

NOTE – In Recommendation Q.921 issue March 1993, this Annex was Annex E.

F.1 General

The supplier of a protocol implementation claiming to conform to this Recommendation shall complete the following Protocol Implementation Conformance Statement (PICS) proforma and accompany it by the information necessary to identify fully both the supplier and the implementation. The PICS proforma applies to the basic rate user-side interface.

The PICS is a document specifying the capabilities and options which have been implemented, and any features which have been omitted, so that the implementation can be tested for conformance against relevant requirements, and against those requirements only.

This PICS has several uses, the most important are the static conformance review and test case selection in order to identify which conformance tests are applicable to this product.

The PICS proforma is a document, in the form of a questionnaire, normally designed by the protocol specifier or conformance test suite specifier which, when completed for an implementation or system, becomes the PICS.

F.2 Abbreviations and special symbols

APPX Appendix

CPE Customer Premises Equipment

DLCI Data Link Connection Identifier, DLCI = (SAPI, TEI)

DLE Data Link Entity

FR Prefix for the Index number of the Frames group

IUT Implementation Under Test

M Mandatory

N/A Not Applicable

O Optional

O.<n> Optional, but, if chosen, support is required for either at least one or only one of the options in the group labelled by the same numeral <n>

P Prohibited

PC Prefix for the Index number of the Protocol Capabilities group

PICS Protocol Implementation Conformance Statement

<r> receive (frame)

³ Copyright release for PICS proforma

Users of this Recommendation may freely reproduce the PICS proforma in this Annex so that it can be used for its intended purpose and may further publish the completed PICS.

<s>	send (frame)
SAPI	Service Access Point Identifier
SP	Prefix for the Index number of System Parameter group
TEI	Terminal End-point Identifier

F.3 Instructions for completing for PICS Proforma

The main part of the PICS proforma is a fixed-format questionnaire, divided into three sections. Answers to the questionnaire are to be provided in the rightmost column, either by simply marking an answer to indicate a restricted choice (such as Yes or No), or by entering a value or a set or range of values.

A supplier may also provide additional information categorized as either Exceptional Information or Supplementary Information (other than PIXIT). When present, each kind of additional information is to be provided as items labelled X.<i> or S.<i> respectively for cross-reference purposes, where <i> is any unambiguous identification for the item. An exception item should contain the appropriate rationale. The Supplementary Information is not mandatory and the PICS is complete without such information. The presence of optional supplementary or exceptional information should not affect test execution, and will in no way affect static conformance verification.

NOTE – Where an implementation is capable of being configured in more than one way, a single PICS may be able to describe all such configurations. However, the supplier has the choice of providing more than one PICS, each covering some subset of the implementation's configuration capabilities, in case this makes for easier or clearer presentation of the information.

In the case in which an IUT does not implement a condition listed, such as in PC 8, where a CPE may not support Layer three call procedures, the Support column of the PICS proforma table should be completed as: "Yes: No: X: X2". The entry of the exceptional information would read: "X2 This CPE does not support Layer 3 call procedures."

F.4 Global statement of conformance

Global statement: The implementation specified in this PICS meets all the mandatory requirements of the referenced standards:

Yes/No

NOTE – Answering "No" to this question indicates non-conformance to this Recommendation. Non-supported mandatory capabilities are to be listed in the PICS below, with an explanation for the abnormal status of the implementation.

The client will have fully complied with the requirements for a statement of conformance by completing the statement contained in this section. However, the client may find it helpful to continue to complete the detailed tabulations in the sections which follow.

F.5 Protocol Capabilities (PC)

Index	Protocol feature	Status	Reference	Support
PC 1.1	Is the CPE of the non-automatic TEI assignment category?	O.1	3.3.4.2	Yes:___ No:___ X:___
PC 1.2	Is the CPE of the automatic TEI assignment category?	O.1	3.3.4.2	Yes:___ No:___ X:___
PC 1.3	Does the CPE support point-to-point procedures?	O.13	Annex A	Yes:___ No:___ X:___
PC 1.4	Does the CPE support point-to-multipoint procedures?	O.13	Annex A	Yes:___ No:___ X:___
PC 2	If the CPE supports point-to-multipoint configuration, does the CPE support the broadcast data link?	M	5.2	Yes:___ No:___ X:___
PC 2.1	If the CPE supports point-to-point configuration, does the CPE support the broadcast data link?	O	5.2.1, Annex A	Yes:___ No:___ X:___
PC 4	Does the CPE support data link monitor function?	O	5.10	Yes:___ No:___ X:___
PC 5	Does the CPE support reject retransmission procedure?	O	3.6.7, 5.8.1, APPX I	Yes:___ No:___ X:___
PC 6.1	Does the DLE support automatic negotiation of data link layer parameters?	O.2	APPX IV	Yes:___ No:___ X:___
PC 6.2	Does the DLE support internal parameter initialization?	O.2	5.4	Yes:___ No:___ X:___
PC 7	Does the CPE permit concurrent LAPB data link connection within the D-channel?	O	2.3	Yes:___ No:___ X:___
PC 7.1	Does the CPE support multi-selective reject?	O	Annex E	Yes:___ No:___ X:___
Service Access Point Identifier (SAPI)				
PC 8	If the CPE supports Layer 3 call control procedures, is SAPI = 0 supported?	M	3.3.3	Yes:___ No:___ X:___
PC 9	If the CPE supports X.25 Layer 3 packet procedures on D-channel, is SAPI = 16 supported?	M	3.3.3	Yes:___ No:___ X:___
PC 10	If the CPE supports point-to-multipoint is SAPI = 63 supported?	M	3.3.3	Yes:___ No:___ X:___
PC 10.1	If the CPE supports teleaction, is SAPI = 12 supported?	M	3.3.3	Yes:___ No:___ X:___
PC 10.2	If the CPE supports point-to-point, is SAPI = 0 supported?	M	5.2.1, 3.3.3, Annex A	Yes:___ No:___ X:___
PC 11.1	Does the implementation support the association of a given TEI with all SAPs which the CPE supports?	O	3.3.4, 5.3.1, (3.4.3/Q.920)	Yes:___ No:___ X:___
PC 11.2	If the CPE is an X.31 type of packet mode terminal equipment, is a given TEI for point-to-point data link connection (< 127) associated with all SAPs which the CPE supports?	M	3.3.4, 5.3.1, (3.4.3/Q.920)	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
PC11.3	If the CPE supports point-to-point, does the CPE support the association of TEI = 0 with SAPI = 0?	M	Annex A	Yes:___ No:___ X:___
PC 12	Does the implementation support modulus 128 for frames numbering?	M	3.5.2.1, 5.5.1	Yes:___ No:___ X:___
Peer-to-peer procedures				
PC 13	Unacknowledged Information Transfer If the CPE supports point-to-multipoint, does the CPE support UI-command?	M	5.2.2	Yes:___ No:___ X:___
PC 13.1	If the CPE supports point-to-point, does the CPE support UI-command?	O	5.2.1, Annex A	Yes:___ No:___ X:___
PC 14	If the CPE supports UI transfer, is the P/F bit set to 0?	M	5.1.1	Yes:___ No:___ X:___
TEI Management				
PC 15	If the CPE supports point-to-multipoint, does the CPE transmit management entity messages in UI frames with DLCI = (63,127)?	M	5.3.1	Yes:___ No:___ X:___
PC 15.1	If the CPE supports point-to-point, does the CPE transmit management entity messages in UI frames with DLCI = (63,127)?	O	5.3.1	Yes:___ No:___ X:___
TEI Assignment Procedures				
PC 16.1	Does the CPE initiate TEI assignment upon power-up?	O.3	5.3.1	Yes:___ No:___ X:___
PC 16.2	Does the CPE initiate TEI assignment at the time an incoming or an outgoing call is handled, if there is no TEI assigned?	O.3	5.3.1	Yes:___ No:___ X:___
PC 17	If the CPE is of the non-automatic category, does the CPE side management entity assign a TEI value?	M	5.3.2	Yes:___ No:___ X:___
PC 18	If the CPE is of the automatic category and supports point-to-multipoint configuration: Does the CPE side management entity initiate TEI assignment?	M	5.3.2	Yes:___ No:___ X:___
PC 19	Is the Ri randomly generated?	M	5.3.2	Yes:___ No:___ X:___
PC 20	Is the Ai value in an Identity Request message always equal to 127?	M	5.3.2	Yes:___ No:___ X:___
PC 21	Does the CPE retransmit an Identity Request message upon timer T202 expiry?	M	5.3.2.1	Yes:___ No:___ X:___
PC 22	Does the CPE use a new value of Ri in the above instance (PC 21)?	M	5.3.2.1	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
TEI Check Response/Removal/Identity Verify				
PC 23.1	If the CPE supports point-to-multipoint configuration: Does the CPE send a single Identity Check Response message, if the Ai value in the received Identity Check Request message is equal to 127?	O.4	5.3.3.2	Yes:___ No:___ X:___
PC 23.2	Does the CPE send an individual Identity Check Response message, for each TEI which is assigned to it, if the Ai value in the received Identity Check Request message is equal to 127?	O.4	5.3.3.2	Yes:___ No:___ X:___
PC 23.3	Does the CPE send any combination of (multiple) "single" and "individual" Identity Check Response messages in order to report all the TEIs assigned to it, if the Ai value in the received Identity Check Request message is equal to 127?	O.4	5.3.3.2	Yes:___ No:___ X:___
PC 24	Does the CPE support transmitting one Identity Check Response message in response to an Identity Check Request message with Ai < 127, if the TEI value being checked is in use?	M	5.3.3.2	Yes:___ No:___ X:___
PC 25	Does the DLE enter the TEI Unassigned state, upon removal of an automatic TEI?	M	5.3	Yes:___ No:___ X:___
PC 26	Does the CPE send an Identity Request message upon removal of an automatic TEI?	M	5.3.4	Yes:___ No:___ X:___
PC 27.1	If the CPE supports point-to-multipoint configuration, and if an Identity Request message is outstanding: Does the CPE remove the TEI from the DLE on receipt of an Identity Assigned message containing a TEI value which is already in use?	O.5	5.3.2, 5.3.4.2	Yes:___ No:___ X:___
PC 27.2	Does the CPE initiate TEI identity verify procedure on receipt of an Identity Assigned message containing a TEI value which is already in use?	O.5	5.3.2	Yes:___ No:___ X:___
PC 28	If the CPE is of the non-automatic TEI category: Does the CPE notify to the equipment user the need for corrective action after non-automatic TEI removal?	M	5.3.4, 5.3.4.2	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
PC 28.1	If the CPE is of the non-automatic TEI assignment category, does the CPE: i) remove the TEI from the data link layer entity; ii) discard it within the Layer management, and; iii) provide an indication of the removal to the user of the equipment, if one of the conditions for TEI removal applies?	O.14	5.3.4.2	Yes:___ No:___ X:___
PC 28.2	If the CPE is of the non-automatic TEI assignment category, does the CPE: i) remove the TEI from the data link layer entity; ii) retain it within the Layer management; and iii) notify the user of the equipment the need for some corrective action, if one of the conditions for TEI removal applies?	O.14	5.3.4	Yes:___ No:___ X:___
PC 29.1	If the CPE supports point-to-multipoint configuration, and if the CPE checks the TEI of all Identity Assign messages: Does the CPE remove TEI from the DLE on receipt of an Identity Assigned message containing a TEI value which is already in use?	O.6	5.3.2 5.3.4.2	Yes:___ No:___ X:___
PC 29.2	Does the CPE initiate TEI identity verify procedure on receipt of an Identity Assigned message containing a TEI value which is already in use?	O.6	5.3.2	Yes:___ No:___ X:___
PC 30	If the CPE supports point-to-multipoint configuration; and if the CPE initiates a TEI Identity Verify procedure, does the Ai contain the own TEI which has been assigned by ASP (automatic TEI) or entered (non-automatic TEI), respectively?	M	5.3.5.2	Yes:___ No:___ X:___
PC 31	If the CPE initiates the TEI identity verify procedure: Does the CPE remove the TEI from the DLE, if no Identity Check Request message with an Ai = 127 or an Ai value equal to Ai value in the Identity Verify Request message has been received when timer T202 (again) expired after retransmission of the Identity Verify Request message upon expiry of timer T202?	M	5.3.5.3	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
Establishment and Release of Multiple Frame Operation				
PC 32	Does the CPE support multiple frame operation?	M	5.5	Yes:___ No:___ X:___
PC 33.1	Does the DLE initiate multi-frame establishment: a) immediately after TEI assignment?	O.7	5.5	Yes:___ No:___ X:___
PC 33.2	b) when there is an incoming or an outgoing call?	O.7	5.5	Yes:___ No:___ X:___
PC 34.1	c) Does the DLE remain in TEI Assigned state when the multiple frame operation is released?	O.8	5.5.3	Yes:___ No:___ X:___
PC 34.2	d) Does the DLE initiate immediate re-establishment when the multiple frame operation is released?	O.8	5.5.3	Yes:___ No:___ X:___
Unsolicited Commands and Responses				
PC 35.1	If the CPE is of the automatic TEI assignment category: Does the CPE initiate TEI identity verify procedure on the receipt of an unsolicited UA response in the Multiple Frame Established State?	O.9	APPX II, 5.8.7	Yes:___ No:___ X:___
PC 35.2	Does the CPE remove the TEI from the DLE on the receipt of an unsolicited UA response in the Multiple Frame Established State?	O.9	APPX II, 5.8.7	Yes:___ No:___ X:___
PC 36.1	Does the CPE initiate TEI identity verify procedure on the receipt of an unsolicited UA response in the Timer Recovery State?	O.10	APPX II, 5.8.7	Yes:___ No:___ X:___
PC 36.2	Does the CPE remove the TEI from the DLE on the receipt of an unsolicited UA response in the Timer Recovery State?	O.10	APPX II, 5.8.7	Yes:___ No:___ X:___
PC 37.1	Does the CPE remove the TEI from the DLE, after N200 unsuccessful retransmissions of SABME?	O.11	APPX II	Yes:___ No:___ X:___
PC 37.2	Does the CPE initiate the TEI identity verify procedure, after N200 unsuccessful retransmissions of SABME?	O.11	APPX II	Yes:___ No:___ X:___
PC 38.1	Does the CPE remove the TEI from the DLE, after N200 unsuccessful retransmissions of DISC?	O.12	APPX II	Yes:___ No:___ X:___
PC 38.2	Does the CPE initiate the TEI identity verify procedure, after N200 unsuccessful retransmissions of DISC?	O.12	APPX II	Yes:___ No:___ X:___
Point-to-point procedures				
PC 39	If the CPE supports point-to-point configuration, does the CPE support only one TEI?	M	Annex A	Yes:___ No:___ X:___
PC 40	If the CPE supports point-to-point configuration, does the CPE not support peer-to-peer management procedures?	M	Annex A	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
PC 41	If the CPE supports point-to-point configuration, does the CPE support TEI = 0?	M	Annex A	Yes:___ No:___ X:___
PC 42	If the CPE supports point-to-point configuration, does the CPE use acknowledged information transfer service for peer-to-peer communication?	M	Annex A	Yes:___ No:___ X:___
Multi-Selective Reject				
PC 43	Does the implementation maintain a poll sequence number?	M	E.3.5.2.7	Yes:___ No:___ X:___
PC 44	Does the implementation increment the poll sequence number after a frame with the P bit set to 1 is sent?	M	E.3.5.2.7	Yes:___ No:___ X:___
PC 45	Does the N(R) subfield of the control field of the SREJ response contain the sequence number of the oldest missing I frame?	M	E.3.6.7.1	Yes:___ No:___ X:___
PC 46	Does the information field of the SREJ response contain the sequence numbers of the remaining missing I frames?	M	E.3.6.7.1	Yes:___ No:___ X:___
PC 47	Is the identity of the missing I frames indicated by one octet for every I frame?	O.15	E.3.6.7.1	Yes:___ No:___ X:___
PC 48	Is the identity of the missing I frames indicated by one octet for every stand-alone I frame plus a span list for every sequence of two or more contiguously numbered I frames?	O.15	E.3.6.7.1	Yes:___ No:___ X:___
O.1 = Support of at least one of these items is required. O.2 = Support of at least one of these items is required. O.3 = Support of at least one of these items is required. O.4 = Support of one, and only one, of these items is required. O.5 = Support of one, and only one, of these items is required. O.6 = Support of one, and only one, of these items is required. O.7 = Support of at least one of these items is required. O.8 = Support of at least one of these items is required. O.9 = Support of one, and only one, of these items is required. O.10 = Support of one, and only one, of these items is required. O.11 = Support of one, and only one, of these items is required. O.12 = Support of one, and only one, of these items is required. O.13 = Support of one, and only one, of these items is required. O.14 = Support of one, and only one, of these items is required. O.15 = Support of one, and only one, of these items is required.				

F.6 Frames – Protocol Data Units (FR)

Index	Protocol feature	Status	Reference	Support
Frame Format				
FR 1	Format A	M	2.1	Yes:___ No:___ X:___
FR 2	Format B	M	2.1	Yes:___ No:___ X:___
Flag Sequence				
FR 3	Opening flag	M	2.2	Yes:___ No:___ X:___
FR 4	Closing flag	M	2.2	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
Address Field				
FR 5	Two octets	M	2.3	Yes:___ No:___ X:___
FR 6	If the DLE permits concurrent LAPB data link connection within the D-channel, is the one octet address field recognized?	M	2.3	Yes:___ No:___ X:___
Control Field				
	Unacknowledged operation			
FR 7	Single octet	M	2.4	Yes:___ No:___ X:___
	Multiple frame operation			
FR 8	Two octets	M	2.4	Yes:___ No:___ X:___
FR 9	Single octet (unnumbered frame)	M	2.4	Yes:___ No:___ X:___
Order of Bit Transmission				
FR 10	Ascending numerical order	M	2.8.2	Yes:___ No:___ X:___
Field Mapping Convention				
FR 11	Lowest bit number = Lowest order value	M	2.8.3	Yes:___ No:___ X:___
	Do all transmitted frames contain the following fields?			
FR 12.1	– Flag	M	2.2	Yes:___ No:___ X:___
FR 12.2	– Address	M	2.3	Yes:___ No:___ X:___
FR 12.3	– Control	M	2.4	Yes:___ No:___ X:___
FR 12.4	– FCS	M	2.7	Yes:___ No:___ X:___
FR 13	Is the CPE capable of accepting the closing flag as the opening flag of the next frame?	M	2.2	Yes:___ No:___ X:___
FR 14	Does the CPE generate a single flag as above?	O	2.2	Yes:___ No:___ X:___
FR 15	Does the CPE ignore one flag, or two or more consecutive flags that do not delimit frames?	M	2.2	Yes:___ No:___ X:___
FR 16	Are all invalid frames discarded and no action taken?	M	2.9	Yes:___ No:___ X:___
FR 17	Are seven or more contiguous 1 bits interpreted as an abort and the associated frames ignored?	M	2.10	Yes:___ No:___ X:___
FR 18	If the CPE supports the automatic negotiation of data link layer parameters, does the CPE support XID frames?	M	APPX IV	Yes:___ No:___ X:___

F.7 System Parameters (SP)

Index	System parameters	Status	Reference	Support/Range
SP 1	If the DLE supports multiple frame operation: Retransmission time (T200)	M	5.9.1	Yes:___ No:___ Value:___
SP 2	Maximum number of retransmissions (N200) Maximum number of octets in information field (N201)	M	5.9.2	Yes:___ No:___ Value:___
SP 3	For SAP supporting signalling	M	5.9.3	Yes:___ No:___ Value:___

Index	System parameters	Status	Reference	Support/Range
SP 4	For SAP supporting packet on the D-channel	M	5.9.3	Yes:___ No:___ Value:___
SP 5	Maximum number of outstanding I frames (k)			
SP 5	For SAP supporting basic access signalling	M	5.9.5	Yes:___ No:___ Value:___
SP 6	For SAP supporting basic access packet on the D-channel	M	5.9.5	Yes:___ No:___ Value:___
	If the CPE is of the automatic TEI assignment category:			
SP 7	Maximum number of transmissions of TEI Identity Request message (N202)	M	5.9.4	Yes:___ No:___ Value:___
SP 8	Minimum time between the transmission of TEI Identity Request message (T202)	M	5.9.7	Yes:___ No:___ Value:___
	If the CPE supports the data link monitor function:			
SP 9	Maximum time allowed without frames being exchanged (T203)	M	5.9.8	Yes:___ No:___ Value:___
	If the CPE supports the automatic negotiation of data link parameters:			
SP 10	Retransmission time of XID frame (TM20)	M	IV.2	Yes:___ No:___ Value:___
SP 11	Maximum number of retransmissions of XID frame (NM20)	M	IV.2	Yes:___ No:___ Value:___

ANNEX G⁴

Protocol Implementation Conformance Statement (PICS) to Recommendation Q.921 for Basic Rate (Network-side)

G.1 General

The supplier of a protocol implementation claiming to conform to this Recommendation shall complete the following Protocol Implementation Conformance Statement (PICS) proforma and accompany it by the information necessary to identify fully both the supplier and the implementation. The PICS proforma applies to the basic rate user-side interface.

The PICS is a document specifying the capabilities and options which have been implemented, and any features which have been omitted, so that the implementation can be tested for conformance against relevant requirements, and against those requirements only.

This PICS has several uses, the most important are the static conformance review and test case selection in order to identify which conformance tests are applicable to this product.

The PICS proforma is a document, in the form of a questionnaire, normally designed by the protocol specifier or conformance test suite specifier which, when completed for an implementation or system, becomes the PICS.

⁴ **Copyright release for PICS proforma**

Users of this Recommendation may freely reproduce the PICS proforma in this Annex so that it can be used for its intended purpose and may further publish the completed PICS.

G.2 Abbreviations and special symbols

APPX	Appendix
ASP	Assignment Source Point
DLCI	Data Link Connection Identifier, DLCI = (SAPI, TEI)
DLE	Data Link Entity
FR	Prefix for the Index number of the Frames group
IUT	Implementation Under Test
M	Mandatory
N/A	Not Applicable
O	Optional
O.<n>	Optional, but, if chosen, support is required for either at least one or only one of the options in the group labelled by the same numeral <n>
P	Prohibited
PC	Prefix for the Index number of the Protocol Capabilities group
PICS	Protocol Implementation Conformance Statement
<r>	receive (frame)
<s>	send (frame)
SAPI	Service Access Point Identifier
SP	Prefix for the Index number of System Parameter group
TEI	Terminal End-point Identifier

G.3 Instructions for completing for PICS Proforma

The main part of the PICS proforma is a fixed-format questionnaire, divided into three sections. Answers to the questionnaire are to be provided in the rightmost column, either by simply marking an answer to indicate a restricted choice (such as Yes or No), or by entering a value or a set or range of values.

A supplier may also provide additional information categorized as either Exceptional Information or Supplementary Information (other than PIXIT). When present, each kind of additional information is to be provided as items labelled X.<i> or S.<i> respectively for cross-reference purposes, where <i> is any unambiguous identification for the item. An exception item should contain the appropriate rationale. The Supplementary Information is not mandatory and the PICS is complete without such information. The presence of optional supplementary or exceptional information should not affect test execution, and will in no way affect static conformance verification.

NOTE – Where an implementation is capable of being configured in more than one way, a single PICS may be able to describe all such configurations. However, the supplier has the choice of providing more than one PICS, each covering some subset of the implementation's configuration capabilities, in case this makes for easier or clearer presentation of the information.

In the case in which an IUT does not implement a condition listed, such as in PC 8, where a CPE may not support Layer 3 call procedures, the Support column of the PICS proforma table should be completed as: " Yes: ___ No: \surd X: X2". The entry of the exceptional information would read: "X2 This CPE does not support Layer 3 call procedures."

G.4 Global statement of conformance

Global statement: The implementation specified in this PICS meets all the mandatory requirements of the referenced standards:

Yes/No

NOTE – Answering "No" to this question indicates non-conformance to this Recommendation. Non-supported mandatory capabilities are to be listed in the PICS below, with an explanation for the abnormal status of the implementation.

The client will have fully complied with the requirements for a statement of conformance by completing the statement contained in this section. However, the client may find it helpful to continue to complete the detailed tabulations in the sections which follow.

G.5 Protocol Capabilities (PC)

Index	Protocol feature	Status	Reference	Support
PC 1.1	Does the implementation accept non-automatic TEI assignment?	M	3.3.4.2	Yes:___ No:___ X:___
PC 1.2	Does the implementation support automatic TEI assignment?	M	3.3.4.2	Yes:___ No:___ X:___
PC 1.3	Does the implementation support point-to-point procedures?	O.13	Annex A	Yes:___ No:___ X:___
PC 1.4	Does the implementation support point-to-multipoint procedures?	O.13	Annex A	Yes:___ No:___ X:___
PC 2	If the implementation supports point-to-multipoint configuration, does the implementation support the broadcast data link?	M	5.2	Yes:___ No:___ X:___
PC 2.1	If the implementation supports point-to-point configuration, does the implementation support the broadcast data link?	O	5.2.1, Annex A	Yes:___ No:___ X:___
PC 3	Does the implementation support the TEI Identity verify procedure?	O	5.3.5	Yes:___ No:___ X:___
PC 4	Does the implementation support data link monitor function?	O	5.10	Yes:___ No:___ X:___
PC 5	Does the implementation support reject retransmission procedure?	O	3.6.7, 5.8.1, APPX I	Yes:___ No:___ X:___
PC 6.1	Does the implementation support automatic negotiation of data link layer parameters?	O.2	APPX IV	Yes:___ No:___ X:___
PC 6.2	Does the implementation support internal parameter initialization?	O.2	5.4	Yes:___ No:___ X:___
PC 7	Does the implementation permit concurrent LAPB data link connection within the D-channel?	O	2.3	Yes:___ No:___ X:___
PC 7.1	Does the implementation support multi-selective reject?	O	Annex E	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
Service Access Point Identifier (SAPI)				
PC 8	Does the implementation support layer 3 call control procedures (SAPI = 0)?	M	3.3.3	Yes:___ No:___ X:___
PC 9	Does the implementation support X.25 Layer 3 packet procedures on D-channel (SAPI = 16)?	M	3.3.3	Yes:___ No:___ X:___
PC 10	Does the implementation support layer 2 management procedures on D-channel (SAPI = 63)?	M	3.3.3	Yes:___ No:___ X:___
PC 10.1	Does the implementation support teleaction communication on D-channel (SAPI = 12)?	M	3.3.3	Yes:___ No:___ X:___
PC 10.2	If the implementation supports point-to-point, is SAPI = 0 supported?	M	5.2.1, 3.3.3, Annex A	Yes:___ No:___ X:___
PC 11	Does the implementation give priority to SAPI = 0 information?	M	5.2/Q.920	Yes:___ No:___ X:___
PC 11.1	Does the implementation support the association of a given TEI with all SAPs which the implementation supports?	O	3.3.4, 5.3.1, (3.4.3/Q.920)	Yes:___ No:___ X:___
PC 11.2	If the implementation is an X.31 type of packet mode terminal equipment, is a given TEI for point-to-point data link connection (<127) associated with all SAPs which the implementation supports?	M	3.3.4, 5.3.1, (3.4.3/Q.920)	Yes:___ No:___ X:___
PC11.3	If the implementation supports point-to-point, does the implementation support the association of TEI = 0 with SAPI = 0?	M	Annex A	Yes:___ No:___ X:___
PC 12	Does the implementation support modulus 128 for frames numbering?	M	3.5.2.1, 5.5.1	Yes:___ No:___ X:___
Peer-to-peer procedures				
PC 13	Unacknowledged Information Transfer If the implementation supports point-to-multipoint, does the implementation support UI-command?	M	5.2.2	Yes:___ No:___ X:___
PC 13.1	If the implementation supports point-to-point, does the implementation support UI-command?	O	5.2.1, Annex A	Yes:___ No:___ X:___
PC 14	If the implementation supports UI transfer, is the P/F bit set to 0?	M	5.1.1	Yes:___ No:___ X:___
PC 15	Does the implementation discard all UI queues in the case of persistent layer 1 deactivation?	M	5.2.2	Yes:___ No:___ X:___
PC 16	Does the implementation complete all UI data transfer before initiating layer 1 deactivation procedures?	M	5.2.2	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
TEI Management				
PC 17	Does the ASP transmit management entity messages in UI frames with SAPI = 63 and TEI = 127?	M	5.3.1	Yes: __ No: __ X: __
PC 17.1	If the implementation supports point-to-point, does the ASP transmit management entity messages in UI frames with SAPI = 63 and TEI = 127?	O	5.3.1	Yes: __ No: __ X: __
PC 18	Does the ASP allocate, select and assign TEI values?	M	5.3.2	Yes: __ No: __ X: __
PC 19.1	Does the ASP support a map of the full range of automatic TEI values?	O.3	5.3.2	Yes: __ No: __ X: __
PC 19.2	Does the ASP support an updated list of all automatic TEI values available for Assignment, or a smaller subset?	O.3	5.3.2	Yes: __ No: __ X: __
TEI Assignment Procedures				
PC 20	Does the ASP ignore Identity Request messages containing identical Ri values?	M	5.3.2	Yes: __ No: __ X: __
PC 21	Does the ASP ignore Identity Request messages with Ai = 0 to 63?	M	5.3.2	Yes: __ No: __ X: __
PC 22	Does the ASP deny Identity Request messages with Ai = 64 to 126?	M	5.3.2	Yes: __ No: __ X: __
PC 23	Does the ASP initiate TEI check procedure if available TEI values are exhausted?	M	5.3.2	Yes: __ No: __ X: __
TEI Check Procedures				
PC 24	Does the ASP transmit an Identity Check Request message containing either the specific TEI value to be checked or the value 127 when all TEI values are to be checked?	M	5.3.3.2	Yes: __ No: __ X: __
PC 25	When the TEI check procedure is used to test whether a TEI value is in use, does the ASP retransmit an Identity Check Request message containing either the specific TEI value to once if no answer is received?	M	5.3.3.2	Yes: __ No: __ X: __
PC 26	Does the ASP accept a multiple Identity Check Response message in response to an Identity Check Request message with Ai = 127 ?	M	5.3.3.2	Yes: __ No: __ X: __
PC 27	Does the ASP assume that the TEI value under check is free if no response is received from the user?	M	5.3.3.2	Yes: __ No: __ X: __
PC 28	Does the ASP assume that the TEI value being checked is in use on receipt of one Identity Check Response message?	M	5.3.3.2	Yes: __ No: __ X: __
PC 29	Does the ASP assume duplicate TEI assignment on receipt of more than one Identity Check Response message containing the same TEI value?	M	5.3.3.2	Yes: __ No: __ X: __

Index	Protocol feature	Status	Reference	Support
TEI Removal/Identity Verify procedures				
PC 30	Does the ASP remove a non-automatic TEI value when duplicate TEI assignment has occurred?	M	5.3.4.2	Yes: __ No: __ X: __
PC 31	Does the ASP remove an automatic TEI value when either it is no longer in use or duplicate TEI assignment has occurred?	M	5.3.4.2	Yes: __ No: __ X: __
PC 32	Does the ASP transmit twice in succession an Identity Remove message containing either the specific TEI value to be removed or Ai = 127 when all TEI value are to be removed?	M	5.3.4	Yes: __ No: __ X: __
PC 33	Does the ASP respond with an Identity Check Request message, if the TEI Identity verify procedure is implemented and if an Identity Verify message is received from the user?	M	5.3.5	Yes: __ No: __ X: __
Establishment and Release of Multiple Frame Operation				
PC 34	Does the implementation support multiple frame operation?	M	5.5	Yes: __ No: __ X: __
PC 35.1	Does the implementation re-establish the multi-frame (MF) operation: a) on receiving SABME command while in the MF mode of operation?	M	5.7.1	Yes: __ No: __ X: __
PC 35.2	b) if N200 retransmission failures occur while in the Timer recovery condition?	M	5.7.1	Yes: __ No: __ X: __
PC 35.3	c) on receiving a supervisory or unnumbered frame with incorrect length?	M	3.6.11, 5.8.5	Yes: __ No: __ X: __
PC 35.4	d) on receiving an invalid sequence number N(R)?	M	3.6.11, 5.8.5	Yes: __ No: __ X: __
PC 35.5	e) on receiving a frame with an information field exceeding N201 (maximum number of octets)?	M	3.6.11, 5.8.5	Yes: __ No: __ X: __
PC 35.6	f) on receiving a FRMR response?	M	5.8.6	Yes: __ No: __ X: __
PC 35.7	g) on receiving an unsolicited DM (F = 0) response while in the MF mode of operation?	M	5.7.1	Yes: __ No: __ X: __
PC 35.8	h) on receiving an unsolicited DM (F = 1) response while in the Timer recovery condition?	M	5.7.1	Yes: __ No: __ X: __
Error Conditions				
PC 36	Does the implementation transmit a REJ frame in the event of a N(S) sequence error?	M	5.8.1	Yes: __ No: __ X: __
PC 37.1	Does the implementation issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the TEI Assigned state?	O.4	APPX II	Yes: __ No: __ X: __

Index	Protocol feature	Status	Reference	Support
PC 37.2	Does the implementation issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the TEI Assigned state?	O.4	APPX II	Yes:___ No:___ X:___
PC 38.1	Does the implementation issue an MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the Awaiting establishment state?	O.5	APPX II	Yes:___ No:___ X:___
PC 38.2	Does the implementation issue an MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the Awaiting establishment state?	O.5	APPX II	Yes:___ No:___ X:___
PC 39.1	Does the implementation issue an MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the Awaiting release state?	O.6	APPX II	Yes:___ No:___ X:___
PC 39.2	Does the implementation issue an MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the Awaiting release state?	O.6	APPX II	Yes:___ No:___ X:___
PC 40.1	Does the implementation issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the Multiple Frame Established state?	O.7	APPX II	Yes:___ No:___ X:___
PC 40.2	Does the implementation issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and remove the EI value on the receipt of an unsolicited UA response in the Multiple Frame Established state?	O.7	APPX II	Yes:___ No:___ X:___
PC 41.1	Does the implementation issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the Timer Recovery state?	O.8	APPX II	Yes:___ No:___ X:___
PC 41.2	Does the implementation issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the Timer Recovery state?	O.8	APPX II	Yes:___ No:___ X:___
PC 42.1	Does the implementation issue an MDL-ERROR.indication (G) and initiate TEI check procedure, after N200 unsuccessful retransmissions of SABME in the Awaiting establishment state?	O.9	APPX II	Yes:___ No:___ X:___
PC 43.1	Does the implementation issue an MDL-ERROR.indication (H) and initiate TEI check procedure, after N200 unsuccessful retransmissions of DISC in the Awaiting release state?	O.10	APPX II	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
Other network management actions				
PC 44.1	Does the implementation log the event on error code A?	O	APPX II	Yes:___ No:___ X:___
PC 44.2	Does the implementation log the event on error code B?	O	APPX II	Yes:___ No:___ X:___
PC 44.3	Does the implementation log the event on error code E?	O	APPX II	Yes:___ No:___ X:___
PC 44.4	Does the implementation log the event on error code F?	O	APPX II	Yes:___ No:___ X:___
PC 44.5	Does the implementation log the event on error code I?	O	APPX II	Yes:___ No:___ X:___
PC 44.6	Does the implementation log the event on error code J?	O	APPX II	Yes:___ No:___ X:___
PC 44.7	Does the implementation log the event on error code K?	O	APPX II	Yes:___ No:___ X:___
PC 44.8	Does the implementation log the event on error code L?	O	APPX II	Yes:___ No:___ X:___
PC 44.9	Does the implementation log the event on error code N?	O	APPX II	Yes:___ No:___ X:___
PC 44.10	Does the implementation log the event on error code O?	O	APPX II	Yes:___ No:___ X:___
Point-to-point procedures				
PC 45	If the implementation supports point-to-point configuration, does the implementation support only one TEI?	M	Annex A	Yes:___ No: X:___
PC 46	If the implementation supports point-to-point configuration, does the implementation not support peer-to-peer management procedures?	M	Annex A	Yes:___ No:___ X:___
PC 47	If the implementation supports point-to-point configuration, does the implementation support TEI = 0?	M	Annex A	Yes:___ No:___ X:___
PC 48	If the implementation supports point-to-point configuration, does the implementation use acknowledged information transfer service for peer-to-peer communication?	M	Annex A	Yes:___ No:___ X:___
Multi-Selective Reject				
PC 49	Does the implementation maintain a poll sequence number?	M	E.3.5.2.7	Yes:___ No:___ X:___
PC 50	Does the implementation increment the poll sequence number after a frame with the P bit set to 1 is sent?	M	E.3.5.2.7	Yes:___ No:___ X:___
PC 51	Does the N(R) subfield of the control field of the SREJ response contain the sequence number of the oldest missing I frame?	M	E.3.6.7.1	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
PC 52	Does the information field of the SREJ response contain the sequence numbers of the remaining missing I frames?	M	E.3.6.7.1	Yes:___ No:___ X:___
PC 53	Is the identity of the missing I frames indicated by one octet for every I frame?	O.14	E.3.6.7.1	Yes:___ No:___ X:___
PC 54	Is the identity of the missing I frames indicated by one octet for every stand-alone I frame plus a span list for every sequence of two or more contiguously numbered I frames?	O.14	E.3.6.7.1	Yes:___ No:___ X:___
O.2 = Support of at least one of these items is required. O.3 = Support of one, and only one, of these items is required. O.4 = Support of one, and only one, of these items is required. O.5 = Support of one, and only one, of these items is required. O.6 = Support of one, and only one, of these items is required. O.7 = Support of one, and only one, of these items is required. O.8 = Support of one, and only one, of these items is required. O.9 = This action is preferred. O.10 = This action is preferred. O.13 = Support of one, and only one, of these items is required. O.14 = Support of one, and only one, of these items is required.				

G.6 Frames – Protocol Data Units (FR)

Index	Protocol feature	Status	Reference	Support
Frame Format				
FR 1	Format A	M	2.1	Yes:___ No:___ X:___
FR 2	Format B	M	2.1	Yes:___ No:___ X:___
Flag Sequence				
FR 3	Opening flag	M	2.2	Yes:___ No:___ X:___
FR 4	Closing flag	M	2.2	Yes:___ No:___ X:___
Address Field				
FR 5	Two octets	M	2.3	Yes:___ No:___ X:___
FR 6	If the DLE permits concurrent LAPB data link connection within the D-channel, is the one octet address field recognized?	M	2.3	Yes:___ No:___ X:___
Control Field				
FR 7	Unacknowledged operation	M	2.4	Yes:___ No:___ X:___
	Single octet			
FR 8	Multiple frame operation	M	2.4	Yes:___ No:___ X:___
	Two octets			
FR 9	Single octet (unnumbered frame)	M	2.4	Yes:___ No:___ X:___
Order of Bit Transmission				
FR 10	Ascending numerical order	M	2.8.2	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
Field Mapping Convention				
FR 11	Lowest bit number = Lowest order value	M	2.8.3	Yes: __ No: __ X: __
	Do all transmitted frames contain the following fields?			
FR 12.1	– Flag	M	2.2	Yes: __ No: __ X: __
FR 12.2	– Address	M	2.3	Yes: __ No: __ X: __
FR 12.3	– Control	M	2.4	Yes: __ No: __ X: __
FR 12.4	– FCS	M	2.7	Yes: __ No: __ X: __
FR 13	Is the implementation capable of accepting the closing flag as the opening flag of the next frame?	M	2.2	Yes: __ No: __ X: __
FR 14	Does the implementation generate a single flag as above?	O	2.2	Yes: __ No: __ X: __
FR 15	Does the implementation ignore one flag, or two or more consecutive flags that do not delimit frames?	M	2.2	Yes: __ No: __ X: __
FR 16	Are all invalid frames discarded and no action taken?	M	2.9	Yes: __ No: __ X: __
FR 17	Are seven or more contiguous 1 bits interpreted as an abort and the associated frames ignored?	M	2.10	Yes: __ No: __ X: __
FR 18	Does the implementation discard frame types associated with an application (see Table 5/Q.921) not implemented?	M	3.6.1	Yes: __ No: __ X: __
FR 19	If the implementation supports the automatic negotiation of data link layer parameters, does the implementation support XID frames?	M	3.6.12 APPX IV	Yes: __ No: __ X: __
FR 20	Does the implementation discriminate invalid frames and frames with an information field exceeding N201 value?	M	5.8.5	Yes: __ No: __ X: __
FR 21	Does the implementation discard unbounded frames?	M	5.8.5	Yes: __ No: __ X: __

G.7 System Parameters (SP)

Index	System parameters	Status	Reference	Support/Range
SP 1	If the DLE supports multiple frame operation: Retransmission time (T200)	M	5.9.1	Yes: __ No: __ Value: __
SP 2	Maximum number of retransmissions (N200) Maximum number of octets in information field (N201)	M	5.9.2	Yes: __ No: __ Value: __
SP 3	For SAP supporting signalling	M	5.9.3	Yes: __ No: __ Value: __
SP 4	For SAP supporting packet on the D-channel	M	5.9.3	Yes: __ No: __ Value: __
SP 5	Maximum number of outstanding I frames (k) For SAP supporting basic access signalling	M	5.9.5	Yes: __ No: __ Value: __
SP 6	For SAP supporting basic access packet on the D-channel	M	5.9.5	Yes: __ No: __ Value: __

Index	System parameters	Status	Reference	Support/Range
SP 7	Maximum time between retransmission of TEI Identity Check Request messages (T201)			
SP 8	If the implementation supports the data link monitor function: Maximum time allowed without frames being exchanged (T203)	M	5.9.8	Yes:___ No:___ Value:___
SP 9	If the implementation supports the automatic negotiation of data link parameters: Retransmission time of XID frame (TM20)	M	IV.2	Yes:___ No:___ Value:___
SP 10	Maximum number of retransmissions of XID frame (NM20)	M	IV.2	Yes:___ No:___ Value:___

ANNEX H⁵

Protocol Implementation Conformance Statement (PICS) to Recommendation Q.921 for Primary Rate (User-side)

H.1 General

The supplier of a protocol implementation claiming to conform to this Recommendation shall complete the following Protocol Implementation Conformance Statement (PICS) proforma and accompany it by the information necessary to identify fully both the supplier and the implementation. The PICS proforma applies to the basic rate user-side interface.

The PICS is a document specifying the capabilities and options which have been implemented, and any features which have been omitted, so that the implementation can be tested for conformance against relevant requirements, and against those requirements only.

This PICS has several uses, the most important are the static conformance review and test case selection in order to identify which conformance tests are applicable to this product.

The PICS proforma is a document, in the form of a questionnaire, normally designed by the protocol specifier or conformance test suite specifier which, when completed for an implementation or system, becomes the PICS.

H.2 Abbreviations and special symbols

APPX Appendix

CPE Customer Premises Equipment

DLCI Data Link Connection Identifier, DLCI= (SAPI, TEI)

DLE Data Link Entity

FR Prefix for the Index number of the Frames group

IUT Implementation Under Test

⁵ **Copyright release for PICS proforma**

Users of this Recommendation may freely reproduce the PICS proforma in this Annex so that it can be used for its intended purpose and may further publish the completed PICS.

M	Mandatory
N/A	Not Applicable
O	Optional
O.<n>	Optional, but, if chosen, support is required for either at least one or only one of the options in the group labelled by the same numeral <n>
P	Prohibited
PC	Prefix for the Index number of the Protocol Capabilities group
PICS	Protocol Implementation Conformance Statement
<r>	receive (frame)
<s>	send (frame)
SAPI	Service Access Point Identifier
SP	Prefix for the Index number of System Parameter group
TEI	Terminal End-point Identifier

H.3 Instructions for completing for PICS Proforma

The main part of the PICS proforma is a fixed-format questionnaire, divided into three sections. Answers to the questionnaire are to be provided in the rightmost column, either by simply marking an answer to indicate a restricted choice (such as Yes or No), or by entering a value or a set or range of values.

A supplier may also provide additional information categorized as either Exceptional Information or Supplementary Information (other than PIXIT). When present, each kind of additional information is to be provided as items labelled X.<i> or S.<i> respectively for cross-reference purposes, where <i> is any unambiguous identification for the item. An exception item should contain the appropriate rationale. The Supplementary Information is not mandatory and the PICS is complete without such information. The presence of optional supplementary or exceptional information should not affect test execution, and will in no way affect static conformance verification.

NOTE – Where an implementation is capable of being configured in more than one way, a single PICS may be able to describe all such configurations. However, the supplier has the choice of providing more than one PICS, each covering some subset of the implementation's configuration capabilities, in case this makes for easier or clearer presentation of the information.

In the case in which an IUT does not implement a condition listed, such as in PC 8, where a CPE may not support Layer 3 call procedures, the Support column of the PICS proforma table should be completed as: " Yes: No: X: X2". The entry of the exceptional information would read: "X2 This CPE does not support Layer 3 call procedures."

H.4 Global statement of conformance

Global statement: The implementation specified in this PICS meets all the mandatory requirements of the referenced standards:

Yes/No

NOTE – Answering "No" to this question indicates non-conformance to this Recommendation. Non-supported mandatory capabilities are to be listed in the PICS below, with an explanation for the abnormal status of the implementation.

The client will have fully complied with the requirements for a statement of conformance by completing the statement contained in this section. However, the client may find it helpful to continue to complete the detailed tabulations in the sections which follow.

H.5 Protocol Capabilities (PC)

Index	Protocol feature	Status	Reference	Support
PC 1.1	Is the CPE of the non-automatic TEI assignment category?	O.1	3.3.4.2	Yes:___ No:___ X:___
PC 1.2	Is the CPE of the automatic TEI assignment category?	O.1	3.3.4.2	Yes:___ No:___ X:___
PC 1.3	Does the CPE support point-to-point procedures?	M	Annex A	Yes:___ No:___ X:___
PC 2.1	If the CPE supports point-to-point configuration, does the CPE support the broadcast data link?	O	5.2.1, Annex A	Yes:___ No:___ X:___
PC 4	Does the CPE support data link monitor function?	O	5.10	Yes:___ No:___ X:___
PC 5	Does the CPE support reject retransmission procedure?	O	3.6.7, 5.8.1, APPX I	Yes:___ No:___ X:___
PC 6.1	Does the DLE support automatic negotiation of data link layer parameters?	O.2	APPX IV	Yes:___ No:___ X:___
PC 6.2	Does the DLE support internal parameter initialization?	O.2	5.4	Yes:___ No:___ X:___
PC 7	Does the CPE permit concurrent LAPB data link connection within the D-channel?	O	2.3	Yes:___ No:___ X:___
PC 7.1	Does the CPE support multi-selective reject?	O	Annex E	Yes:___ No:___ X:___
Service Access Point Identifier (SAPI)				
PC 8	If the CPE supports Layer 3 call control procedures, is SAPI = 0 supported?	M	3.3.3	Yes:___ No:___ X:___
PC 9	If the CPE supports X.25 Layer 3 packet procedures on D-channel, is SAPI = 16 supported?	M	3.3.3	Yes:___ No:___ X:___
PC 10	If the CPE supports point-to-multipoint, is SAPI = 63 supported?	M	3.3.3	Yes:___ No:___ X:___
PC 10.1	If the CPE supports teleaction communication on D-channel, is SAPI = 12 supported?	M	3.3.3	Yes:___ No:___ X:___
PC 10.2	If the CPE supports point-to-point, is SAPI = 0 supported?	M	5.2.1, 3.3.3, Annex A	Yes:___ No:___ X:___
PC 11.1	Does the implementation support the association of a given TEI with all SAPs which the CPE supports?	O	3.3.4, 5.3.1, (3.4.3/Q.920)	Yes:___ No:___ X:___
PC 11.2	If the CPE is an X.31 type of packet mode terminal equipment, is a given TEI for point-to-point data link connection (<127) associated with all SAPs which the CPE supports?	M	3.3.4, 5.3.1 (3.4.3/Q.920)	Yes:___ No:___ X:___
PC11.3	If the CPE supports point-to-point, does the CPE support the association of TEI = 0 with SAPI = 0?	M	Annex A	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
PC 12	Does the implementation support modulus 128 for frames numbering?	M	3.5.2.1, 5.5.1	Yes:___ No:___ X:___
Peer-to-peer procedures				
PC 13.1	Unacknowledged Information Transfer If the CPE supports point-to-point, does the CPE support UI-command?	O	5.2.1, Annex A	Yes:___ No:___ X:___
PC 14	If the CPE supports UI transfer, is the P/F bit set to 0?	M	5.1.1	Yes:___ No:___ X:___
TEI Management				
PC 15.1	If the CPE supports point-to-point, does the CPE transmit management entity messages in UI frames with DLCI = (63,127)	O	5.3.1	Yes:___ No:___ X:___
TEI Assignment Procedures				
PC 16.1	Does the CPE initiate TEI assignment upon power-up?	O.3	5.3.1	Yes:___ No:___ X:___
PC 16.2	Does the CPE initiate TEI assignment at the time an incoming or an outgoing call is handled, if there is no TEI assigned?	O.3	5.3.1	Yes:___ No:___ X:___
PC 17	If the CPE is of the non-automatic category, does the CPE side management entity assign a TEI value?	M	5.3.2	Yes:___ No:___ X:___
PC 18	If the CPE is of the automatic category: Does the CPE side management entity initiate TEI assignment?	M	5.3.2	Yes:___ No:___ X:___
PC 19	Is the Ri randomly generated?	M	5.3.2	Yes:___ No:___ X:___
PC 20	Is the Ai value in an Identity Request message always equal to 127?	M	5.3.2	Yes:___ No:___ X:___
PC 21	Does the CPE retransmit an Identity Request message upon timer T202 expiry?	M	5.3.2.1	Yes:___ No:___ X:___
PC 22	Does the CPE use a new value of Ri in the above instance (PC 21)?	M	5.3.2.1	Yes:___ No:___ X:___
TEI Check Response/Removal/Identity Verify				
PC 23.1	If the CPE supports TEI management (see Annex A): Does the CPE send a single Identity Check Response message, if the Ai value in the received Identity Check Request message is equal to 127?	O.4	5.3.3.2	Yes:___ No:___ X:___
PC 23.2	Does the CPE send an individual Identity Check Response message, for each TEI which is assigned to it, if the Ai value in the received Identity Check Request message is equal to 127?	O.4	5.3.3.2	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
PC 23.3	Does the CPE send any combination of (multiple) "single" and "individual" Identity Check Response messages in order to report all the TEIs assigned to it, if the Ai value in the received Identity Check Request message is equal to 127?	O.4	5.3.3.2	Yes:___ No:___ X:___
PC 24	Does the CPE support transmitting one Identity Check Response message in response to an Identity Check Request message with Ai < 127, if the TEI value being checked is in use?	M	5.3.3.2	Yes:___ No:___ X:___
PC 25	Does the DLE enter the TEI Unassigned state, upon removal of an automatic TEI?	M	5.3	Yes:___ No:___ X:___
PC 26	Does the CPE send an Identity Request message upon removal of an automatic TEI? If the CPE supports TEI management (see Annex A), and if an Identity Request message is outstanding:	M	5.3.4	Yes:___ No:___ X:___
PC 27.1	Does the CPE remove the TEI from the DLE on receipt of an Identity Assigned message containing a TEI value which is already in use?	O.5	5.3.2 5.3.4.2	Yes:___ No:___ X:___
PC 27.2	Does the CPE initiate a TEI identity verify procedure on receipt of an Identity Assigned message containing a TEI value which is already in use?	O.5	5.3.2	Yes:___ No:___ X:___
PC 28.1	If the CPE is of the non-automatic TEI assignment category, does the CPE: i) remove the TEI from the data link layer entity; ii) discard it within the Layer management; and iii) provide an indication of the removal to the user of the equipment, if one of the conditions for TEI removal applies?	O.17	5.3.4, 5.3.4.2	Yes:___ No:___ X:___
PC 28.2	If the CPE is of the non-automatic TEI assignment category, does the CPE: i) remove the TEI from the data link layer entity; ii) retain it within the Layer management; and iii) notify the user of the equipment the need for some corrective action, if one of the conditions for TEI removal applies?	O.17	5.3.4, 5.3.4.2	Yes:___ No:___ X:___
PC 29.1	If the CPE supports TEI management (see Annex A), and if the CPE checks the TEI of all Identity Assign messages: Does the CPE remove TEI from the DLE on receipt of an Identity Assigned message containing a TEI value which is already in use?	O.6	5.3.2, 5.3.4.2	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
PC 29.2	Does the CPE initiate a TEI identity verify procedure on receipt of an Identity Assigned message containing a TEI value which is already in use?	O.6	5.3.2	Yes:___ No:___ X:___
PC 30	If the CPE supports TEI management (see Annex A), and if the CPE initiates a TEI Identity Verify procedure, does the Ai contain the own TEI which has been assigned by ASP (automatic TEI) or entered (non-automatic TEI), respectively? If the CPE initiates the TEI identity verify procedure:	M	5.3.5.2	Yes:___ No:___ X:___
PC 31	Does the CPE remove the TEI from the DLE, if no Identity Check Request message with an Ai = 127 or an Ai value equal to Ai value in the Identity Verify Request message has been received when timer T202 (again) expired after retransmission of the Identity Verify Request message upon expiry of timer T202?	M	5.3.5.3	Yes:___ No:___ X:___
Establishment and Release of Multiple Frame Operation				
PC 32	Does the CPE support multiple frame operation?	M	5.5	Yes:___ No:___ X:___
PC 33.1	Does the DLE initiate multi-frame establishment a) immediately after TEI assignment?	O.7	5.5	Yes:___ No:___ X:___
PC 33.2	b) when there is an incoming or an outgoing call?	O.7	5.5	Yes:___ No:___ X:___
PC 34.1	c) Does the DLE remain in TEI Assigned state when the multiple frame operation is released?	O.8	5.5.3	Yes:___ No:___ X:___
PC 34.2	d) Does the DLE initiate immediate re-establishment when the multiple frame operation is released?	O.8	5.5.3	Yes:___ No:___ X:___
PC 35.1	Does the CPE re-establish the multi-frame (MF) operation: a) on receiving SABME command while in the MF mode of operation?	M	5.7.1	Yes:___ No:___ X:___
PC 35.2	b) if N200 retransmission failures occur while in the Timer recovery condition?	M	5.7.1	Yes:___ No:___ X:___
PC 35.3	c) on receiving a supervisory or unnumbered frame with incorrect length?	M	3.6.11 5.8.5	Yes:___ No:___ X:___
PC 35.4	d) on receiving an invalid sequence number N(R)?	M	3.6.11 5.8.5	Yes:___ No:___ X:___
PC 35.5	e) on receiving a frame with an information field exceeding N201 (maximum number of octets)?	M	3.6.11 5.8.5	Yes:___ No:___ X:___
PC 35.6	f) on receiving a FRMR response?	M	5.8.6	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
PC 35.7	g) on receiving an unsolicited DM (F = 0) response while in the MF mode of operation?	M	5.7.1	Yes:___ No:___ X:___
PC 35.7	h) on receiving an unsolicited DM (F = 1) response while in the Timer recovery condition?	M	5.7.1	Yes:___ No:___ X:___
Error Conditions				
PC 36	Does the CPE transmit a REJ frame in the event of a N(S) sequence error?	M	5.8.1	Yes:___ No:___ X:___
PC 37.1	Does the CPE issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the TEI Assigned state?	O.10	5.8.7, APPX II	Yes:___ No:___ X:___
PC 37.2	Does the CPE issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the TEI Assigned state?	O.10	5.8.7, APPX II	Yes:___ No:___ X:___
PC 38.1	Does the CPE issue an MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the Awaiting establishment state?	O.11	5.8.7 APPX II	Yes:___ No:___ X:___
PC 38.2	Does the CPE issue an MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the Awaiting establishment state?	O.11	5.8.7, APPX II	Yes:___ No:___ X:___
PC 39.1	Does the CPE issue an MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the Awaiting release state?	O.12	5.8.7, APPX II	Yes:___ No:___ X:___
PC 39.2	Does the CPE issue an MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the Awaiting release state?	O.12	5.8.7, APPX II	Yes:___ No:___ X:___
PC 40.1	Does the CPE issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the Multiple Frame Established state?	O.13	5.8.7, APPX II	Yes:___ No:___ X:___
PC 40.2	Does the CPE issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the Multiple Frame Established state?	O.13	5.8.7, APPX II	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
PC 41.1	Does the CPE issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the Timer Recovery state?	O.14	5.8.7, APPX II	Yes:___ No:___ X:___
PC 41.2	Does the CPE issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the Timer Recovery state?	O.14	5.8.7, APPX II	Yes:___ No:___ X:___
PC 42.1	Does the CPE issue an MDL-ERROR.indication (G) and initiate TEI check procedure, after N200 unsuccessful retransmissions of SABME in the Awaiting establishment state?	O.15	APPX II	Yes:___ No:___ X:___
PC 42.2	Does the CPE issue an MDL-ERROR.indication (G) and remove the TEI value, after N200 unsuccessful retransmissions of SABME in the Awaiting establishment state?	O.15	APPX II	Yes:___ No:___ X:___
PC 43.1	Does the CPE issue an MDL-ERROR.indication (H) and initiate TEI check procedure, after N200 unsuccessful retransmissions of DISC in the Awaiting release state?	O.16	APPX II	Yes:___ No:___ X:___
PC 43.2	Does the CPE issue an MDL-ERROR.indication (H) and remove the TEI value, after N200 unsuccessful retransmissions of DISC in the Awaiting release state?	O.16	APPX II	Yes:___ No:___ X:___
Point-to-point procedures				
PC 45	If the CPE supports point-to-point configuration, does the CPE support only one TEI?	M	Annex A	Yes:___ No: X:___
PC 46	If the CPE supports point-to-point configuration, does the CPE not support peer-to-peer management procedures?	M	Annex A	Yes:___ No:___X:___
PC 47	If the CPE supports point-to-point configuration, does the CPE support TEI = 0?	M	Annex A	Yes:___ No:___X:___
PC 48	If the CPE supports point-to-point configuration, does the CPE use acknowledged information transfer service for peer-to-peer communication?	M	Annex A	Yes:___ No:___X:___
Multi-Selective Reject				
PC 49	Does the implementation maintain a poll sequence number?	M	E.3.5.2.7	Yes:___ No:___ X:___
PC 50	Does the implementation increment the poll sequence number after a frame with the P bit set to 1 is sent?	M	E.3.5.2.7	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
PC 51	Does the N(R) subfield of the control field of the SREJ response contain the sequence number of the oldest missing I frame?	M	E.3.6.7.1	Yes:___ No:___ X:___
PC 52	Does the information field of the SREJ response contain the sequence numbers of the remaining missing I frames?	M	E.3.6.7.1	Yes:___ No:___ X:___
PC 53	Is the identity of the missing I frames indicated by one octet for every I frame?	O.14	E.3.6.7.1	Yes:___ No:___ X:___
PC 54	Is the identity of the missing I frames indicated by one octet for every stand-alone I frame plus a span list for every sequence of two or more contiguously numbered I frames?	O.14	E.3.6.7.1	Yes:___ No:___ X:___
O.1 = Support of at least one of these items is required. O.2 = Support of at least one of these items is required. O.3 = Support of at least one of these items is required. O.4 = Support of one, and only one, of these items is required. O.5 = Support of one, and only one, of these items is required. O.6 = Support of one, and only one, of these items is required. O.7 = Support of at least one of these items is required. O.8 = Support of at least one of these items is required. O.10 = Support of one, and only one, of these items is required. O.11 = Support of one, and only one, of these items is required. O.12 = Support of one, and only one, of these items is required. O.13 = Support of one, and only one, of these items is required. O.14 = Support of one, and only one, of these items is required. O.15 = Support of one, and only one, of these items is required. O.16 = Support of one, and only one, of these items is required. O.17 = Support of one, and only one, of these items is required.				

H.6 Frames – Protocol Data Units (FR)

Index	Protocol feature	Status	Reference	Support
Frame Format				
FR 1	Format A	M	2.1	Yes:___ No:___ X:___
FR 2	Format B	M	2.1	Yes:___ No:___ X:___
Flag Sequence				
FR 3	Opening flag	M	2.2	Yes:___ No:___ X:___
FR 4	Closing flag	M	2.2	Yes:___ No:___ X:___
Address Field				
FR 5	Two octets	M	2.3	Yes:___ No:___ X:___
FR 6	If the DLE permits concurrent LAPB data link connection within the D-channel, is the one octet address field recognized?	M	2.3	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
Control Field				
FR 7	Unacknowledged operation Single octet	M	2.4	Yes:___ No:___ X:___
FR 8	Multiple frame operation Two octets	M	2.4	Yes:___ No:___ X:___
FR 9	Single octet (unnumbered frame)	M	2.4	Yes:___ No:___ X:___
Order of Bit Transmission				
FR 10	Ascending numerical order	M	2.8.2	Yes:___ No:___ X:___
Field Mapping Convention				
FR 11	Lowest bit number = Lowest order value	M	2.8.3	Yes:___ No:___ X:___
FR 12.1	Do all transmitted frames contain the following fields? – Flag	M	2.2	Yes:___ No:___ X:___
FR 12.2	– Address	M	2.3	Yes:___ No:___ X:___
FR 12.3	– Control	M	2.4	Yes:___ No:___ X:___
FR 12.4	– FCS	M	2.7	Yes:___ No:___ X:___
FR 13	Is the CPE capable of accepting the closing flag as the opening flag of the next frame?	M	2.2	Yes:___ No:___ X:___
FR 14	Does the CPE generate a single flag as above?	O	2.2	Yes:___ No:___ X:___
FR 15	Does the CPE ignore one flag, or two or more consecutive flags that do not delimit frames?	M	2.2	Yes:___ No:___ X:___
FR 16	Are all invalid frames discarded and no action taken?	M	2.9	Yes:___ No:___ X:___
FR 17	Are seven or more contiguous 1 bits interpreted as an abort and the associated frames ignored?	M	2.10	Yes:___ No:___ X:___
FR 18	If the CPE supports the automatic negotiation of data link layer parameters, does the CPE support XID frames?	M	APPX IV	Yes:___ No:___ X:___

H.7 System Parameters (SP)

Index	System parameters	Status	Reference	Support/Range
SP 1	If the DLE supports multiple frame operation: Retransmission time (T200)	M	5.9.1	Yes:___ No:___ Value:___
SP 2	Maximum number of retransmissions (N200) Maximum number of octets in information field (N201)	M	5.9.2	Yes:___ No:___ Value:___
SP 3	For SAP supporting signalling	M	5.9.3	Yes:___ No:___ Value:___
SP 4	For SAP supporting packet on the D-channel Maximum number of outstanding I frames (k)	M	5.9.3	Yes:___ No:___ Value:___
SP 5	For SAP supporting primary access signalling	M	5.9.5	Yes:___ No:___ Value:___
SP 6	For SAP supporting primary access packet on the D-channel	M	5.9.5	Yes:___ No:___ Value:___

Index	System parameters	Status	Reference	Support/Range
SP 7	If the CPE is of the automatic TEI assignment category: Maximum number of transmissions of TEI Identity Request message (N202)	M	5.9.4	Yes:___ No:___ Value:___
SP 8	Minimum time between the transmission of TEI Identity Request message (T202)	M	5.9.7	Yes:___ No:___ Value:___
SP 9	If the CPE supports the data link monitor function: Maximum time allowed without frames being exchanged (T203)	M	5.9.8	Yes:___ No:___ Value:___
SP 10	If the CPE supports the automatic negotiation of data link parameters: Retransmission time of XID frame (TM20)	M	IV.2	Yes:___ No:___ Value:___
SP 11	Maximum number of retransmissions of XID frame (NM20)	M	IV.2	Yes:___ No:___ Value:___

ANNEX I⁶

Protocol Implementation Conformance Statement (PICS) to Recommendation Q.921 for Primary Rate (Network-side)

I.1 General

The supplier of a protocol implementation claiming to conform to this Recommendation shall complete the following Protocol Implementation Conformance Statement (PICS) proforma and accompany it by the information necessary to identify fully both the supplier and the implementation. The PICS proforma applies to the basic rate user-side interface.

The PICS is a document specifying the capabilities and options which have been implemented, and any features which have been omitted, so that the implementation can be tested for conformance against relevant requirements, and against those requirements only.

This PICS has several uses, the most important are the static conformance review and test case selection in order to identify which conformance tests are applicable to this product.

The PICS proforma is a document, in the form of a questionnaire, normally designed by the protocol specifier or conformance test suite specifier which, when completed for an implementation or system, becomes the PICS.

I.2 Abbreviations and special symbols

APPX Appendix

ASP Assignment Source Point

DLCI Data Link Connection Identifier, DLCI = (SAPI, TEI)

⁶ Copyright release for PICS proforma

Users of this Recommendation may freely reproduce the PICS proforma in this Annex so that it can be used for its intended purpose and may further publish the completed PICS.

DLE	Data Link Entity
FR	Prefix for the Index number of the Frames group
IUT	Implementation Under Test
M	Mandatory
N/A	Not Applicable
O	Optional
O.<n>	Optional, but, if chosen, support is required for either at least one or only one of the options in the group labelled by the same numeral <n>
P	Prohibited
PC	Prefix for the Index number of the Protocol Capabilities group
PICS	Protocol Implementation Conformance Statement
<r>	receive (frame)
<s>	send (frame)
SAPI	Service Access Point Identifier
SP	Prefix for the Index number of System Parameter group
TEI	Terminal End-point Identifier

I.3 Instructions for completing for PICS Proforma

The main part of the PICS proforma is a fixed-format questionnaire, divided into three sections. Answers to the questionnaire are to be provided in the rightmost column, either by simply marking an answer to indicate a restricted choice (such as Yes or No), or by entering a value or a set or range of values.

A supplier may also provide additional information categorized as either Exceptional Information or Supplementary Information (other than PIXIT). When present, each kind of additional information is to be provided as items labelled X.<i> or S.<i> respectively for cross-reference purposes, where <i> is any unambiguous identification for the item. An exception item should contain the appropriate rationale. The Supplementary Information is not mandatory and the PICS is complete without such information. The presence of optional supplementary or exceptional information should not affect test execution, and will in no way affect static conformance verification.

NOTE – Where an implementation is capable of being configured in more than one way, a single PICS may be able to describe all such configurations. However, the supplier has the choice of providing more than one PICS, each covering some subset of the implementation's configuration capabilities, in case this makes for easier or clearer presentation of the information.

In the case in which an IUT does not implement a condition listed, such as in PC 8, where a CPE may not support Layer 3 call procedures, the Support column of the PICS proforma table should be completed as: " Yes: No: X: X2". The entry of the exceptional information would read: "X2 This CPE does not support Layer 3 call procedures."

I.4 Global statement of conformance

Global statement: The implementation specified in this PICS meets all the mandatory requirements of the referenced standards:

Yes/No

NOTE – Answering "No" to this question indicates non-conformance to this Recommendation. Non-supported mandatory capabilities are to be listed in the PICS below, with an explanation for the abnormal status of the implementation.

The client will have fully complied with the requirements for a statement of conformance by completing the statement contained in this section. However, the client may find it helpful to continue to complete the detailed tabulations in the sections which follow.

I.5 Protocol Capabilities (PC)

Index	Protocol feature	Status	Reference	Support
PC 1.1	Does the implementation accept non-automatic TEI assignment?	M	3.3.4.2	Yes: __ No: __ X: __
PC 1.2	Does the implementation support automatic TEI assignment?	M	3.3.4.2	Yes: __ No: __ X: __
PC 1.3	Does the implementation support point-to-point procedures?	M	Annex A	Yes: __ No: __ X: __
PC 2.1	If the implementation supports point-to-point configuration, does the implementation support the broadcast data link?	O	5.2.1, Annex A	Yes: __ No: __ X: __
PC 3	Does the implementation support the TEI Identity verify procedure?	O	5.3.5	Yes: __ No: __ X: __
PC 4	Does the implementation support data link monitor function?	O	5.10	Yes: __ No: __ X: __
PC 5	Does the implementation support reject retransmission procedure?	O	3.6.7 5.8.1, APPX I	Yes: __ No: __ X: __
PC 6.1	Does the implementation support automatic negotiation of data link layer parameters?	O.2	APPX IV	Yes: __ No: __ X: __
PC 6.2	Does the implementation support internal parameter initialization?	O.2	5.4	Yes: __ No: __ X: __
PC 7	Does the implementation permit concurrent LAPB data link connection within the D-channel?	O	2.3	Yes: __ No: __ X: __
PC 7.1	Does the implementation support multi-selective reject?	O	Annex E	Yes: __ No: __ X: __
Service Access Point Identifier (SAPI)				
PC 8	Does the implementation support layer 3 call control procedures (SAPI = 0)?	M	3.3.3	Yes: __ No: __ X: __
PC 9	Does the implementation support X.25 Layer 3 packet procedures on D-channel (SAPI = 16)?	M	3.3.3	Yes: __ No: __ X: __
PC 10	Does the implementation support layer 2 management procedures on D-channel (SAPI = 63)?	M	3.3.3	Yes: __ No: __ X: __
PC 10.1	Does the implementation support teleaction communication on D-channel (SAPI = 12)?	M	3.3.3	Yes: __ No: __ X: __
PC 10.2	If the implementation supports point-to-point, is SAPI = 0 supported?	M	5.2.1, 3.3.3, Annex A	Yes: __ No: __ X: __
PC 11	Does the implementation give priority to SAPI = 0 information?	M	5.2/Q.920	Yes: __ No: __ X: __

Index	Protocol feature	Status	Reference	Support
PC 11.1	Does the implementation support the association of a given TEI with all SAPs which the implementation supports?	O	3.3.4, 5.3.1, (3.4.3/Q.920)	Yes:___ No:___ X:___
PC 11.2	If the implementation is an X.31 type of packet mode terminal equipment, is a given TEI for point-to-point data link connection (<127) associated with all SAPs which the implementation supports?	M	3.3.4, 5.3.1, (3.4.3/Q.920)	Yes:___ No:___ X:___
PC11.3	If the implementation supports point-to-point, does the implementation support the association of TEI = 0 with SAPI = 0?	M	Annex A	Yes:___ No:___ X:___
PC 12	Does the implementation support modulus 128 for frames numbering?	M	3.5.2.1, 5.5.1	Yes:___ No:___ X:___
Peer-to-peer procedures				
PC 13.1	Unacknowledged Information Transfer If the implementation supports point-to-point, does the implementation support UI-command?	O	5.2.1, Annex A	Yes:___ No:___ X:___
PC 14	If the implementation supports UI transfer, is the P/F bit set to 0?	M	5.1.1	Yes:___ No:___ X:___
TEI Management				
PC 17	Does the ASP transmit management entity messages in UI frames with SAPI = 63 and TEI = 127?	M	5.3.1	Yes:___ No:___ X:___
PC 17.1	If the implementation supports point-to-point, does the ASP transmit management entity messages in UI frames with SAPI = 63 and TEI = 127?	O	5.3.1	Yes:___ No:___ X:___
PC 18	Does the ASP allocate, select and assign TEI values?	M	5.3.2	Yes:___ No:___ X:___
PC 19.1	Does the ASP support a map of the full range of automatic TEI values?	O.3	5.3.2	Yes:___ No:___ X:___
PC 19.2	Does the ASP support an updated list of all automatic TEI values available for Assignment, or a smaller subset?	O.3	5.3.2	Yes:___ No:___ X:___
TEI Assignment Procedures				
PC 20	Does the ASP ignore Identity Request messages containing identical Ri values?	M	5.3.2	Yes:___ No:___ X:___
PC 21	Does the ASP ignore Identity Request messages with Ai = 0 to 63?	M	5.3.2	Yes:___ No:___ X:___
PC 22	Does the ASP deny Identity Request messages with Ai = 64 to 126?	M	5.3.2	Yes:___ No:___ X:___
PC 23	Does the ASP initiate TEI check procedure if available TEI values are exhausted?	M	5.3.2	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
TEI Check Procedures				
PC 24	Does the ASP transmit an Identity Check Request message containing either the specific TEI value to be checked or the value 127 when all TEI values are to be checked?	M	5.3.3.2	Yes:___ No:___ X:___
PC 25	When the TEI check procedure is used to test whether a TEI value is in use, does the ASP retransmit an Identity Check Request message containing either the specific TEI value to once if no answer is received?	M	5.3.3.2	Yes:___ No:___ X:___
PC 26	Does the ASP accept a multiple Identity Check Response message in response to an Identity Check Request message with Ai = 127 ?	M	5.3.3.2	Yes:___ No:___ X:___
PC 27	Does the ASP assume that the TEI value under check is free if no response is received from the user?	M	5.3.3.2	Yes:___ No:___ X:___
PC 28	Does the ASP assume that the TEI value being checked is in use on receipt of one Identity Check Response message?	M	5.3.3.2	Yes:___ No:___ X:___
PC 29	Does the ASP assume duplicate TEI assignment on receipt of more than one Identity Check Response message containing the same TEI value?	M	5.3.3.2	Yes:___ No:___ X:___
TEI Removal/Identity Verify procedures				
PC 30	Does the ASP remove a non-automatic TEI value when duplicate TEI assignment has occurred?	M	5.3.4.2	Yes:___ No:___ X:___
PC 31	Does the ASP remove an automatic TEI value when either it is no longer in use or duplicate TEI assignment has occurred?	M	5.3.4.2	Yes:___ No:___ X:___
PC 32	Does the ASP transmit twice in succession an Identity Remove message containing either the specific TEI value to be removed or Ai = 127 when all TEI values are to be removed?	M	5.3.4	Yes:___ No:___ X:___
PC 33	Does the ASP respond with an Identity Check Request message, if the TEI Identity verify procedure is implemented and if an Identity Verify message is received from the user?	M	5.3.5	Yes:___ No:___ X:___
Establishment and Release of Multiple Frame Operation				
PC 34	Does the implementation support multiple frame operation?	M	5.5	Yes:___ No:___ X:___
PC 35.1	Does the implementation re-establish the multi-frame (MF) operation: a) on receiving SABME command while in the MF mode of operation?	M	5.7.1	Yes:___ No:___ X:___
PC 35.2	b) if N200 retransmission failures occur while in the Timer recovery condition?	M	5.7.1	Yes:___ No:___ X:___
PC 35.3	c) on receiving a supervisory or unnumbered frame with incorrect length?	M	3.6.11, 5.8.5	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
PC 35.4	d) on receiving an invalid sequence number N(R)?	M	3.6.11, 5.8.5	Yes:___ No:___ X:___
PC 35.5	e) on receiving a frame with an information field exceeding N201 (maximum number of octets)?	M	3.6.11, 5.8.5	Yes:___ No:___ X:___
PC 35.6	f) on receiving a FRMR response?	M	5.8.6	Yes:___ No:___ X:___
PC 35.7	g) on receiving an unsolicited DM (F = 0) response while in the MF mode of operation?	M	5.7.1	Yes:___ No:___ X:___
PC 35.8	h) on receiving an unsolicited DM (F = 1) response while in the Timer recovery condition?	M	5.7.1	Yes:___ No:___ X:___
Error Conditions				
PC 36	Does the implementation transmit a REJ frame in the event of a N(S) sequence error?	M	5.8.1	Yes:___ No:___ X:___
PC 37.1	Does the implementation issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the TEI Assigned state?	O.4	APPX II	Yes:___ No:___ X:___
PC 37.2	Does the implementation issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the TEI Assigned state?	O.4	APPX II	Yes:___ No:___ X:___
PC 38.1	Does the implementation issue an MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the Awaiting establishment state?	O.5	APPX II	Yes:___ No:___ X:___
PC 38.2	Does the implementation issue an MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the Awaiting establishment state?	O.5	APPX II	Yes:___ No:___ X:___
PC 39.1	Does the implementation issue an MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the Awaiting release state?	O.6	APPX II	Yes:___ No:___ X:___
PC 39.2	Does the implementation issue an MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the Awaiting release state?	O.6	APPX II	Yes:___ No:___ X:___
PC 40.1	Does the implementation issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the Multiple Frame Established state?	O.7	APPX II	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
PC 40.2	Does the implementation issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the Multiple Frame Established state?	O.7	APPX II	Yes:___ No:___ X:___
PC 41.1	Does the implementation issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and initiate TEI check procedure on the receipt of an unsolicited UA response in the Timer Recovery state?	O.8	APPX II	Yes:___ No:___ X:___
PC 41.2	Does the implementation issue an MDL-ERROR.indication (C) or MDL-ERROR.indication (D) and remove the TEI value on the receipt of an unsolicited UA response in the Timer Recovery state?	O.8	APPX II	Yes:___ No:___ X:___
PC 42.1	Does the implementation issue an MDL-ERROR.indication (G) and initiate TEI check procedure, after N200 unsuccessful retransmissions of SABME in the Awaiting establishment state?	O.9	APPX II	Yes:___ No:___ X:___
PC 43.1	Does the implementation issue an MDL-ERROR.indication (H) and initiate TEI check procedure, after N200 unsuccessful retransmissions of DISC in the Awaiting release state?	O.10	APPX II	Yes:___ No:___ X:___
Other network management actions				
PC 44.1	Does the implementation log the event on error code A?	O	APPX II	Yes:___ No:___ X:___
PC 44.2	Does the implementation log the event on error code B?	O	APPX II	Yes:___ No:___ X:___
PC 44.3	Does the implementation log the event on error code E?	O	APPX II	Yes:___ No:___ X:___
PC 44.4	Does the implementation log the event on error code F?	O	APPX II	Yes:___ No:___ X:___
PC 44.5	Does the implementation log the event on error code I?	O	APPX II	Yes:___ No:___ X:___
PC 44.6	Does the implementation log the event on error code J?	O	APPX II	Yes:___ No:___ X:___
PC 44.7	Does the implementation log the event on error code K?	O	APPX II	Yes:___ No:___ X:___
PC 44.8	Does the implementation log the event on error code L?	O	APPX II	Yes:___ No:___ X:___
PC 44.9	Does the implementation log the event on error code N?	O	APPX II	Yes:___ No:___ X:___
PC 44.10	Does the implementation log the event on error code O?	O	APPX II	Yes:___ No:___ X:___

Index	Protocol feature	Status	Reference	Support
Point-to-point procedures				
PC 45	If the implementation supports point-to-point configuration, does the implementation support only one TEI?	M	Annex A	Yes: __ No: X: __
PC 46	If the implementation supports point-to-point configuration, does the implementation not support peer-to-peer management procedures?	M	Annex A	Yes: __ No: __ X: __
PC 47	If the implementation supports point-to-point configuration, does the implementation support TEI = 0?	M	Annex A	Yes: __ No: __ X: __
PC 48	If the CPE supports point-to-point configuration, does the CPE use acknowledged information transfer service for peer-to-peer communication?	M	Annex A	Yes: __ No: __ X: __
Multi-Selective Reject				
PC 49	Does the implementation maintain a poll sequence number?	M	E.3.5.2.7	Yes: __ No: __ X: __
PC 50	Does the implementation increment the poll sequence number after a frame with the P bit set to 1 is sent?	M	E.3.5.2.7	Yes: __ No: __ X: __
PC 51	Does the N(R) subfield of the control field of the SREJ response contain the sequence number of the oldest missing I frame?	M	E.3.6.7.1	Yes: __ No: __ X: __
PC 52	Does the information field of the SREJ response contain the sequence numbers of the remaining missing I frames?	M	E.3.6.7.1	Yes: __ No: __ X: __
PC 53	Is the identity of the missing I frames indicated by one octet for every I frame?	O.18	E.3.6.7.1	Yes: __ No: __ X: __
PC 54	Is the identity of the missing I frames indicated by one octet for every standalone I frame plus a span list for every sequence of two or more contiguously numbered I frames?	O.18	E.3.6.7.1	Yes: __ No: __ X: __
O.2 = Support of at least one of these items is required. O.3 = Support of at least one of these items is required. O.4 = Support of one, and only one, of these items is required. O.5 = Support of one, and only one, of these items is required. O.6 = Support of one, and only one, of these items is required. O.7 = Support of one, and only one, of these items is required. O.8 = Support of one, and only one, of these items is required. O.9 = This action is preferred. O.10 = This action is preferred. O.18 = Support of one, and only one, of these items is required.				

I.6 Frames – Protocol Data Units (FR)

Index	Protocol feature	Status	Reference	Support
Frame Format				
FR 1	Format A	M	2.1	Yes: __ No: __ X: __
FR 2	Format B	M	2.1	Yes: __ No: __ X: __
Flag Sequence				
FR 3	Opening flag	M	2.2	Yes: __ No: __ X: __
FR 4	Closing flag	M	2.2	Yes: __ No: __ X: __
Address Field				
FR 5	Two octets	M	2.3	Yes: __ No: __ X: __
FR 6	If the DLE permits concurrent LAPB data link connection within the D-channel, is the one octet address field recognized?	M	2.3	Yes: __ No: __ X: __
Control Field				
FR 7	Unacknowledged operation	M	2.4	Yes: __ No: __ X: __
	Single octet			
FR 8	Multiple frame operation	M	2.4	Yes: __ No: __ X: __
	Two octets			
FR 9	Single octet (unnumbered frame)	M	2.4	Yes: __ No: __ X: __
Order of Bit Transmission				
FR 10	Ascending numerical order	M	2.8.2	Yes: __ No: __ X: __
Field Mapping Convention				
FR 11	Lowest bit number = Lowest order value	M	2.8.3	Yes: __ No: __ X: __
	Do all transmitted frames contain the following fields?			
FR 12.1	– Flag	M	2.2	Yes: __ No: __ X: __
FR 12.2	– Address	M	2.3	Yes: __ No: __ X: __
FR 12.3	– Control	M	2.4	Yes: __ No: __ X: __
FR 12.4	– FCS	M	2.7	Yes: __ No: __ X: __
FR 13	Is the implementation capable of accepting the closing flag as the opening flag of the next frame?	M	2.2	Yes: __ No: __ X: __
FR 14	Does the implementation generate a single flag as above?	O	2.2	Yes: __ No: __ X: __
FR 15	Does the implementation ignore one flag, or two or more consecutive flags that do not delimit frames?	M	2.2	Yes: __ No: __ X: __
FR 16	Are all invalid frames discarded and no action taken?	M	2.9	Yes: __ No: __ X: __
FR 17	Are seven or more contiguous 1 bits interpreted as an abort and the associated frames ignored?	M	2.10	Yes: __ No: __ X: __
FR 18	Does the implementation discard frame types associated with an application (see Table 5/Q.921) not implemented?	M	3.6.1	Yes: __ No: __ X: __

Index	Protocol feature	Status	Reference	Support
FR 19	If the implementation supports the automatic negotiation of data link layer parameters, does the implementation support XID frames?	M	3.6.12, APPX IV	Yes: __ No: __ X: __
FR 20	Does the implementation discriminate invalid frames and frames with an information field exceeding N201 value?	M	5.8.5	Yes: __ No: __ X: __
FR 21	Does the implementation discard unbounded frames?	M	5.8.5	Yes: __ No: __ X: __

I.7 System Parameters (SP)

Index	System parameters	Status	Reference	Support/Range
SP 1	If the DLE supports multiple frame operation Retransmission time (T200)	M	5.9.1	Yes: __ No: __ Value: __
SP 2	Maximum number of retransmissions (N200) Maximum number of octets in information field (N201)	M	5.9.2	Yes: __ No: __ Value: __
SP 3	For SAP supporting signalling	M	5.9.3	Yes: __ No: __ Value: __
SP 4	For SAP supporting packet on the D-channel Maximum number of outstanding I frames (k)	M	5.9.3	Yes: __ No: __ Value: __
SP 5	For SAP supporting primary access signalling	M	5.9.5	Yes: __ No: __ Value: __
SP 6	For SAP supporting primary access packet on the D-channel	M	5.9.5	Yes: __ No: __ Value: __
SP 7	Maximum time between retransmission of TEI Identity Check Request messages (T201)			
SP 8	If the implementation supports the data link monitor function: Maximum time allowed without frames being exchanged (T203)	M	5.9.8	Yes: __ No: __ Value: __
SP 9	If the implementation supports the automatic negotiation of data link parameters: Retransmission time of XID frame (TM20)	M	IV.2	Yes: __ No: __ Value: __
SP 10	Maximum number of retransmissions of XID frame (NM20)	M	IV.2	Yes: __ No: __ Value: __

APPENDIX I

Retransmission of REJ response frames

I.1 Introduction

This Appendix describes an optional procedure which may be used to provide a reject retransmission procedure.

I.2 Procedure

This optional reject retransmission procedure can supplement the Q.921 LAPD protocol by defining a new variable for multiple frame operation (see 3.5.2), and by modifying the N(S) sequence error exception condition reporting and recovery (see 5.8.1).

I.2.1 Recovery state variable V(M)

Each point-to-point data link entity may have an associated V(M) when using I frame commands and supervisory frame commands/responses. V(M) denotes the sequence number of the last frame received which caused an N(S) sequence error condition. V(M) can take on the value 0 to 127 and may be used to determine if another REJ response frame should be sent on receipt of an N(S) sequence error while in the REJ exception condition.

I.2.2 N(S) sequence error supplementary procedure

The first three paragraphs of subclause 5.8.1, N(S) sequence error, apply. The remainder of the subclause is as follows:

The REJ frame is used by a receiving data link layer entity to initiate an exception recovery (retransmission) following the detection of an N(S) sequence error. The receiving data link entity shall set V(M) to the N(S) sequence number which caused the N(S) sequence error condition.

Only one REJ exception condition for a given direction of information transfer shall be established at a time [that is, all REJ frames must have the same N(R) value until the REJ reception is cleared].

A data link layer entity receiving an REJ command or response shall initiate sequential transmission (retransmission) of I frames starting with the I frame indicated by the N(R) contained in the REJ frame.

An REJ exception is cleared when the requested I frame is received or when SABME, or DISC is received.

If an N(S) sequence error exception occurs when the receiving data link layer entity is in the REJ exception condition, then it shall check the N(S) of the received frame to see if the data link layer entity which received the REJ frame has retransmitted in response to the REJ frame [i.e. is N(S) within the range $V(R) + 1 \leq N(S) \leq V(M)$]. If the N(S) of the received frame is within the above range, then it shall send another REJ response frame, issue an MDL-ERROR indication primitive to the connection management entity, and it shall set V(M) equal to N(S). The transmitting side will not need to wait for timer T200 to expire before it can retransmit the lost frame.

If an N(S) sequence error occurs when the receiving data link layer entity is in the REJ exception condition, and it cannot be determined if the data link layer entity which received the REJ frame has retransmitted in response to that frame [i.e. if $N(S) > V(M)$], then it shall set V(M) equal to the N(S) of the received frame.

APPENDIX II

Occurrence of MDL-ERROR indication within the basic states and actions to be taken by the management entity

II.1 Introduction

Table II.1 gives the error situations in which the MDL-ERROR indication primitive will be generated. This primitive notifies the data link layer's connection management entity of the occurred error situation. The associated error parameter contains the error code that describes the unique error conditions. Table II.1 also identifies the associated connection management actions to be taken from the network and the user side, based on the types of error conditions reported.

This Appendix does not incorporate the retransmission of REJ response frames described in Appendix I.

II.2 Layout of Table II.1

The "Error code" column gives the identification value of each error situation to be included as a parameter with the MDL-ERROR indication primitive.

The column entitled "Error condition" together with the "Affected states" describes unique protocol error events and the basic state of the data link layer entity at the point that the MDL-ERROR indication primitive is generated.

For a given error condition, the column entitled "Network management action" describes the preferred action to be taken by the network management entity.

The column entitled "User management action" describes the preferred action to be taken by the user side management entity on a given error condition.

II.3 Preferred management actions

The various preferred layer management actions on an error situation may be described as one of the following:

a) *Error log*

This suggests that the network side connection management entity has the preferred action of logging the event into an error counter. The length and the operation of the counter mechanisms for the error situations is implementation dependent.

b) *TEI check*

This means that the network side layer management entity invokes the TEI check procedure.

c) *TEI verify*

This means that the user side layer management entity may optionally invoke a TEI verify request procedure that asks the network side layer management entity to issue a TEI check procedure.

d) *TEI remove*

This means that the user side layer management entity may directly remove its TEI value from service.

In most of the described error situations, there is either no action to be taken on the user side layer management or the action to be taken is implementation dependent, as Table II.1 shows. "Implementation dependent" means that it is optional whether the user side layer management has incorporated any form of error counter to log (store) the reported event. If action is taken, the layer management has to take into account that the data link layer will have initiated a recovery procedure.

Table II.1/Q.921 – Management Entity Actions for MDL-Error-Indications

Error type	Error code	Error condition	Affected states (Note 1)	Network management action	User management action
Receipt of unsolicited response	A	Supervisory (F = 1)	7	Error log	Implementation dependent
	B	DM (F = 1)	7, 8	Error log	Implementation dependent
	C	UA (F = 1)	4, 7, 8	TEI removal procedure or TEI check procedure; then, if TEI:	TEI identity verify procedure, if implemented, or remove TEI
	D	UA (F = 0)	4, 5, 6, 7,8	<ul style="list-style-type: none"> – free, remove TEI; – single, no action; – multiple, TEI removal procedure. 	
	E	Receipt of DM response (F = 0)	7, 8	Error log	Implementation dependent
Peer initiated re-establishment	F	SABME	7, 8	Error log	Implementation dependent
Unsuccessful retransmission (N200 times)	G	SABME	5	TEI check procedure; then, if TEI:	TEI identity verify procedure, if implemented, or remove TEI
	H	DIS	6	<ul style="list-style-type: none"> – free, remove TEI; – single, error log; – multiple, TEI removal procedure. 	
	I	Status enquiry	8	Error log	Implementation dependent

Table II.1/Q.921 – Management Entity Actions for MDL-Error-Indications (concluded)

Error type	Error code	Error condition	Affected states (Note 1)	Network management action	User management action
Other	J	N(R) error	7, 8	Error log	Implementation dependent
	K	Receipt of FRMR response	7, 8	Error log	Implementation dependent
	L	Receipt of undefined frame	4, 5, 6, 7, 8	Error log	Implementation dependent
	M (Note 2)	Receipt of I field not permitted	4, 5, 6, 7, 8	Error log	Implementation dependent
	N	Receipt of frame with wrong size	4, 5, 6, 7, 8	Error log	Implementation dependent
	O	N201 error	4, 5, 6, 7, 8	Error log	Implementation dependent
NOTE 1 – For the description of the affected states, see Annex B.					
NOTE 2 – According to 5.8.5, this error code will never be generated.					

APPENDIX III

Optional basic access deactivation procedures

III.1 Introduction

This Appendix provides one example of a deactivation procedure which optionally may be used by the network side system management to control deactivation of the access. Figure III.1 provides a conceptual model of the interactions which are required for this deactivation procedure.

III.2 Description of the Conceptual Model

The monitor function uses layer 2 activity as the basis for establishing whether deactivation of the access can take place. The signal INFORMATION is used to report the layer 2 activity in the following manner:

- a) INFORMATION (FREE) indicates that there is no data link connection in the multiple-frame mode of operation;
- b) INFORMATION (IN USE) indicates that there is at least one data link connection in the mode-setting or multiple-frame mode of operation; and
- c) INFORMATION (UNIT DATA) indicates that a UI frame is about to be transmitted, or has just been received.

Within the data link layer entity the DL-ESTABLISH request/indication primitives and DL-RELEASE indication/confirm primitives mark the duration of the multiple-frame mode of operation, and the MDL/DL-UNIT DATA request/indication primitives mark the transmission and reception of UI frames.

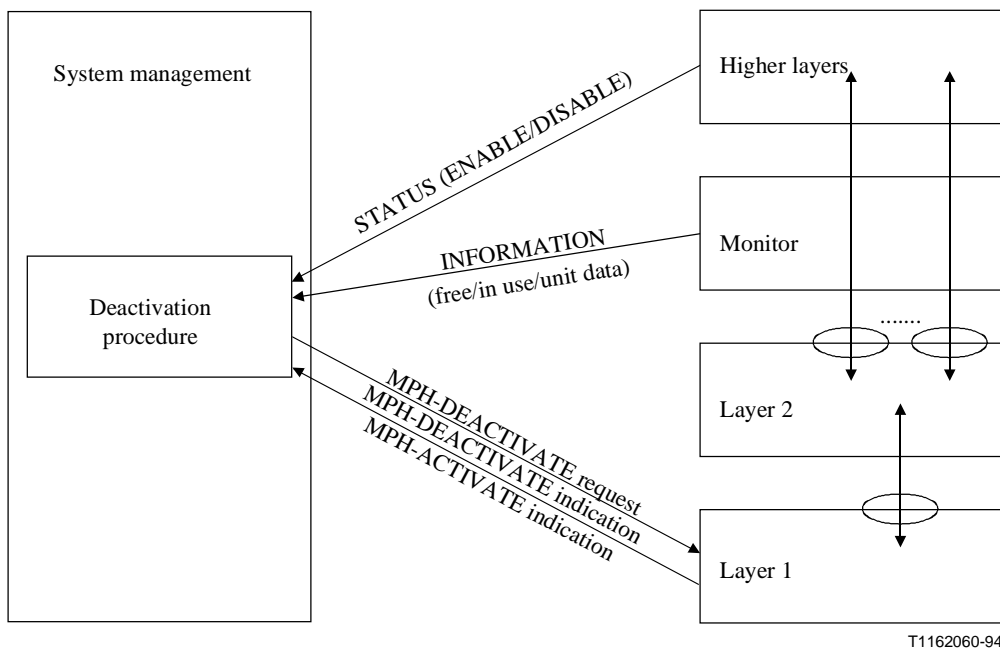


Figure III.1/Q.921 – Conceptual model of the interactions for an example deactivation procedure

A signal Status is used to represent the ability of higher layers to enable or disable the deactivation procedures:

- STATUS (ENABLE) deactivation procedures enabled; and
- STATUS (DISABLE) deactivation procedures disabled.

The MPH-DEACTIVATE request, MPH-DEACTIVATE indication and MPH-ACTIVATE indication primitives are used as described in clause 4. The definition and usage of these primitives are also described in Recommendation I.430 [5] which specifies layer 1.

Since, in Recommendation I.430 [5] the usage of the MPH-DEACTIVATE indication primitive is an implementation option, two cases of deactivation are described below.

Subclause III.3 provides a description of the deactivation procedure when the MPH-DEACTIVATE indication primitive is delivered to the system management entity.

Subclause III.4 provides a description of the deactivation procedure when the MPH-DEACTIVATE indication primitive is not delivered to the system management entity.

NOTE – These procedures require that all layer 3 entities making use of the acknowledged information transfer service must release the data link connection at an appropriate point after the completion of the information transfer.

III.3 Deactivation procedure with MPH-DEACTIVATE indication

This deactivation procedure makes use of the MPH-DEACTIVATE indication primitive to provide an option of layer 1 implementation.

Figure III.2 provides a state transition diagram of the deactivation procedure with the MPH-DEACTIVATE indication primitive.

This deactivation procedure can be represented by six states:

- State 1: Information transfer not available and free;
(*No information transfer and free*)
- State 2: Information transfer available and free;
(*Information transfer and free*)
- State 3: Information transfer available and in use;
(*Information transfer and in use*)
- State 4: Information transfer not available and in use;
(*No information transfer and in use*)
- State 5: Information transfer interrupted and free;
(*Information interrupted and free*)
- State 6: Information transfer interrupted and in use;
(*Information interrupted and in use*)

These six states are described as follows:

- a) State 1 represents the state where the access is assumed to be deactivated and no data link connections are in a mode setting or multiple-frame mode of operation.
- b) State 2 represents the state where the access is activated and no data link connection is in a mode setting or multiple-frame mode of operation. Timer TM01 is running, and upon its expiry, if deactivation is enabled, then an MPH-DEACTIVATE request primitive may be issued to layer 1. The access is then assumed to be deactivated.
- c) State 3 represents the state where the access is activated and at least one data link connection is in a mode setting or multiple-frame mode of operation.
- d) State 4 represents the state where the access is regarded as being in a transient state (neither deactivated nor activated) and at least one data link connection is in a mode setting or multiple-frame mode of operation. [This state can be entered, for example, due to the arrival of an INFORMATION (IN USE) signal before an MPH-ACTIVATE indication primitive.]
- e) State 5 represents the state where the access is regarded as being in a transient state (neither deactivated nor activated) and no data link connection is in a mode setting or multiple-frame mode of operation. Timer TM01 is running and upon its expiry, if deactivation is enabled, then an MPH-DEACTIVATE request primitive will be issued to layer 1. The access is assumed to be deactivated.
- f) State 6 represents the state where the access is regarded as being in the transient state (neither deactivated nor activated) and at least one data link connection is in a mode setting or multiple frame mode of operation.

Timer TM01 is started whenever state 2 is entered:

- i) on receipt of an MPH-ACTIVATE indication primitive in state 1; and
- ii) on receipt of an INFORMATION (FREE) signal in state 3.

Timer TM01 is started whenever state 5 is entered:

- on receipt of an INFORMATION (FREE) signal in state 6.

Timer TM01 is restarted in states 2 and 5 when:

- TM01 expires while deactivation is disabled by the receipt of a STATUS (DISABLE) signal; and
- an INFORMATION (UNIT DATA) signal is received in order to allow sufficient time for current and further unacknowledged information transfer.

Timer TM01 has a value of ten seconds at the network side.

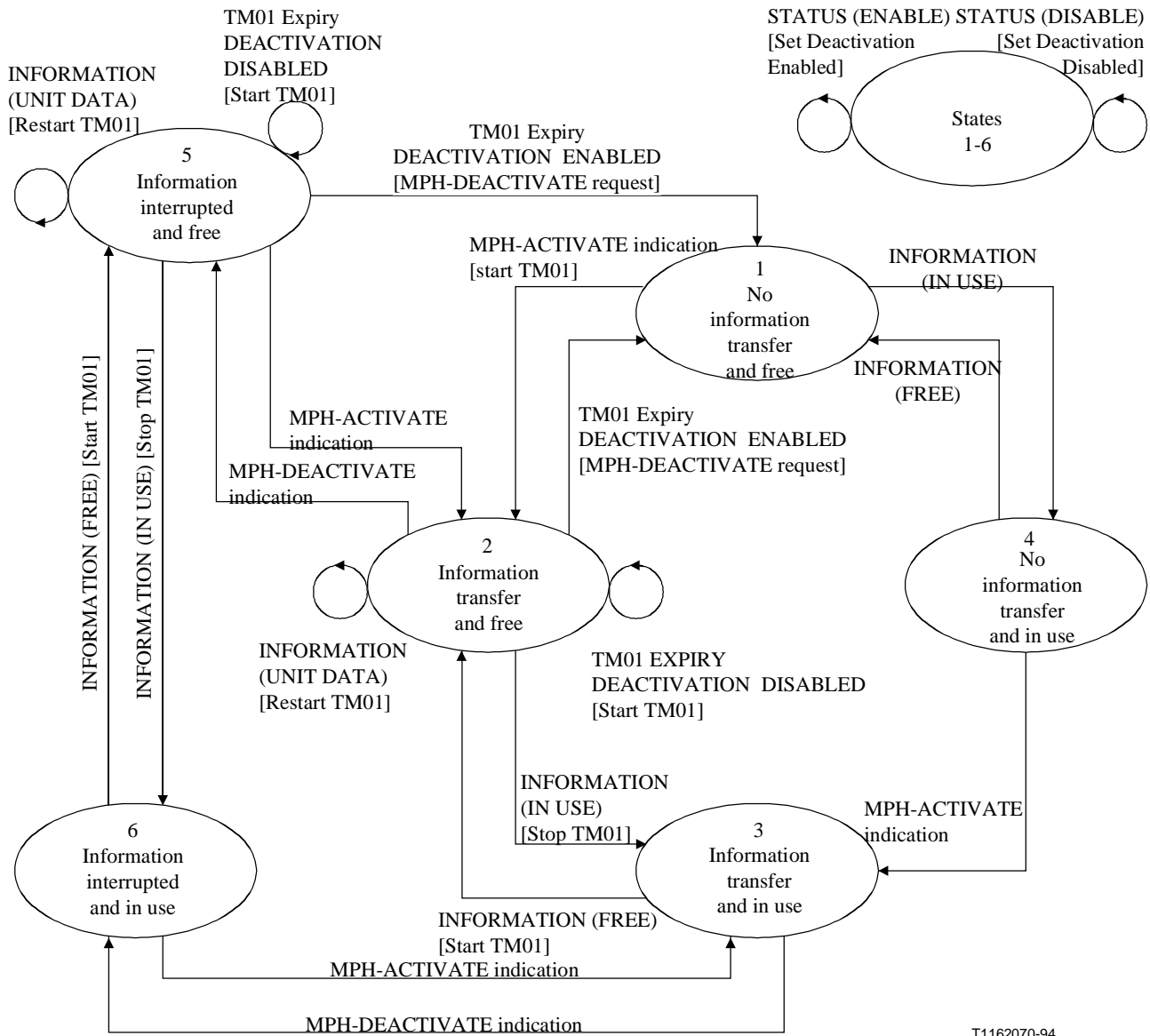


Figure III.2/Q.921 – State transition diagram of a deactivation procedure with MPH-DEACTIVATE indication

III.4 Deactivation procedure without MPH-DEACTIVATE indication

This deactivation procedure does not make use of the MPH-DEACTIVATE indication primitive to provide an option of layer 1 implementation. Thus this procedure can be represented by only four states, i.e. state 1, state 2, state 3, and state 4. States 5 and 6 have disappeared.

Figure III.3 provides a state transition diagram of this deactivation procedure without the MPH-DEACTIVATE indication primitive.

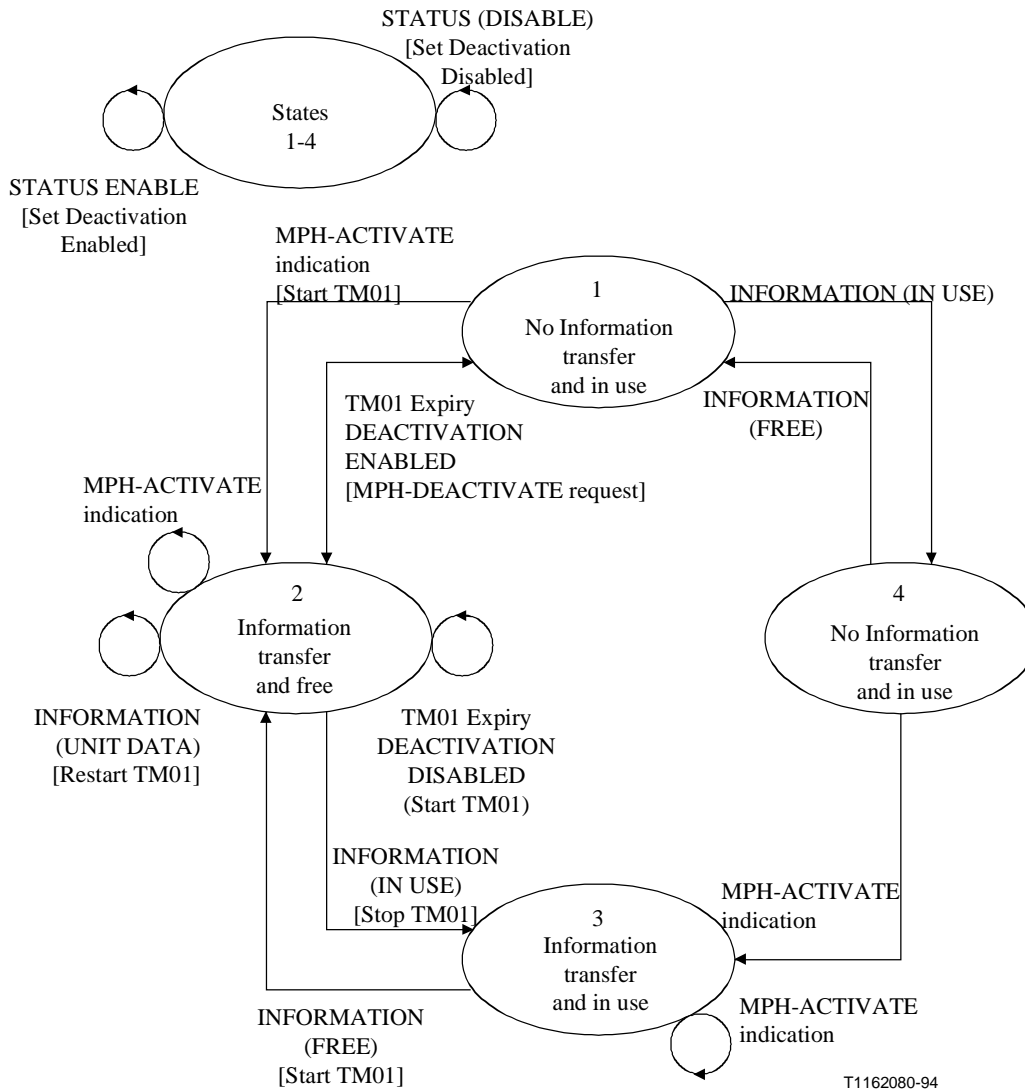


Figure III.3/Q.921 – State transition diagram of a deactivation procedure without MPH-DEACTIVATE indication

APPENDIX IV

Automatic negotiation of data link layer parameters

IV.1 General

The initialization of data link layer parameters is defined in 5.4. This Appendix defines a procedure suitable to negotiate these parameters with a peer entity.

Typically, after the assignment of a TEI value to the management entity, the data link connection management entity is notified by its layer management entity that parameter initialization is required.

The data link connection management entity will then invoke the peer-to-peer negotiation procedure.

IV.2 Automatic negotiation of data link layer parameter values

For each data link layer an exchange of certain data link layer parameters may take place between the peer data link connection management entities before entering the *TEI-assigned* state. This exchange may be initiated after acquiring a TEI, that is, after:

- receipt of a DL-ESTABLISH request or a DL-UNIT DATA request primitive following a power-up condition associated with non-automatic TEI user equipment;
- receipt of the Identity assigned response message for automatic TEI assignment user equipment. This message contains the TEI received by the layer management entity.

All messages used for automatic negotiation of data link parameters are carried in the information field of XID frames with a TEI value set to the value acquired as indicated above and with a SAPI set to a value identical to that associated with the TEI on the data link entity whose parameters are being negotiated. Once a TEI value has been assigned to a terminal which supports multiple data link access points (e.g. SAPI = 0 assigned to call control procedures and SAPI = 16 assigned to packet mode communications), this terminal may negotiate link layer parameters for each SAPI in use.

The data link connection management entity, following assignment of a TEI from the layer management entity, shall issue an XID command with the P bit set to 0 and the I field coded as shown in Figure IV.1, and start the connection management timer TM20.

The I field of the XID command frame shall reflect the parameters desired for future communications across this data link layer connection.

The peer data link connection management entity, upon receipt of this XID command frame, shall transmit an XID response with the F bit set to 0 containing the list of parameter values that the peer can support.

If the data link connection management entity receives the above XID response prior to expiry of timer TM20, it shall stop the timer, and shall notify the layer management entity of a successful parameter exchange. However, if timer TM20 expires before receiving the XID response, the data link connection management entity shall retransmit the XID command, increment the retransmission counter and restart timer TM20. This retransmission process is repeated if timer TM20 expires again. Should the retransmission counter equal NM20, or an XID response frame with a zero length I field be received, the data link connection management entity shall issue an indication to the layer management entity and initialize the parameters to the default values. The layer management entity may log this condition and then issue the MDL-ASSIGN request primitive to the data link layer.

The timer TM20 is set to 2.5 seconds and NM20 is set to 3.

Octet	8	7	6	5	4	3	2	1	
5	1	0	0	0	0	0	1	0	Format Identifier (FI)
6	1	0	0	0	0	0	0	0	Group Identifier (GI)
7	0	0	0	0	0	0	0	0	Group Length (GL)
8	0	0	0	0	1	1	1	0	Group Length (GL)
9	0	0	0	0	0	1	0	1	Parameter Identifier (PI) = Frame Size (Transmit)
10	0	0	0	0	0	0	1	0	Parameter length (PL) = 2
11	2^{15}							2^8	Parameter Value (PV) = N201 Value of Transmitter
12	2^7							2^0	PV = N201 Value of Transmitter
13	0	0	0	0	0	1	1	0	PI = Frame Size (Receive)
14	0	0	0	0	0	0	1	0	PL = 2
15	2^{15}							2^8	PV = N201 Value of Receiver
16	2^7							2^0	PV = N201 Value of Receiver
17	0	0	0	0	0	1	1	1	PI = Window Size (Transmit)
18	0	0	0	0	0	0	0	1	PL = 1
19	0	2^6						2^0	PV = k Value
20	0	0	0	0	1	0	0	1	PI = Retransmission Timer (T200)
21	0	0	0	0	0	0	0	1	PL = 1
22	2^7							2^0	PV = T200 Value ^{a)}

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^{a)} Increments of 0.1 seconds; maximum range 25.5 seconds.

Figure IV.1/Q.921 – XID I field encoding for parameter negotiation

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